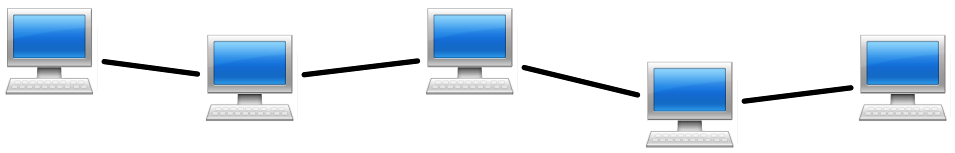
- **Weeks 1–2: informal introduction**
  - network = path 
- **Week 3: graph theory**
- **Weeks 4–7: models of computing**
  - what can be computed (efficiently)?
- **Weeks 8–11: lower bounds**
  - what cannot be computed (efficiently)?
- **Week 12: recap**

Exam

- **Problems, model solutions, grading, feedback: available in Noppa**
- **Quick overview:**
  - max points: 24
  - point distribution: 14 · 15 · 18 · 18 · 18 · 24
  - 18/24 ≈ grade 4/5

Week 7

– Randomised algorithms

Deterministic algorithms

- **init_d(…):** state
- **send_d(…):** message vector
- **receive_d(…):** state

Randomised algorithms

- **init$_d$(…):** *probability distribution* over states
- **send$_d$(…):** message vector
- **receive$_d$(…):** *probability distribution* over states

Randomised algorithms

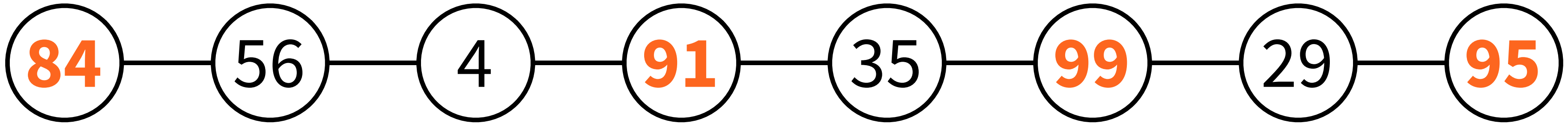
- You can always toss coins when you pick the new state

Randomised algorithms

- **Randomised algorithm in PN model**
- **Randomised algorithm in LOCAL model**
- **Randomised algorithm in CONGEST model**

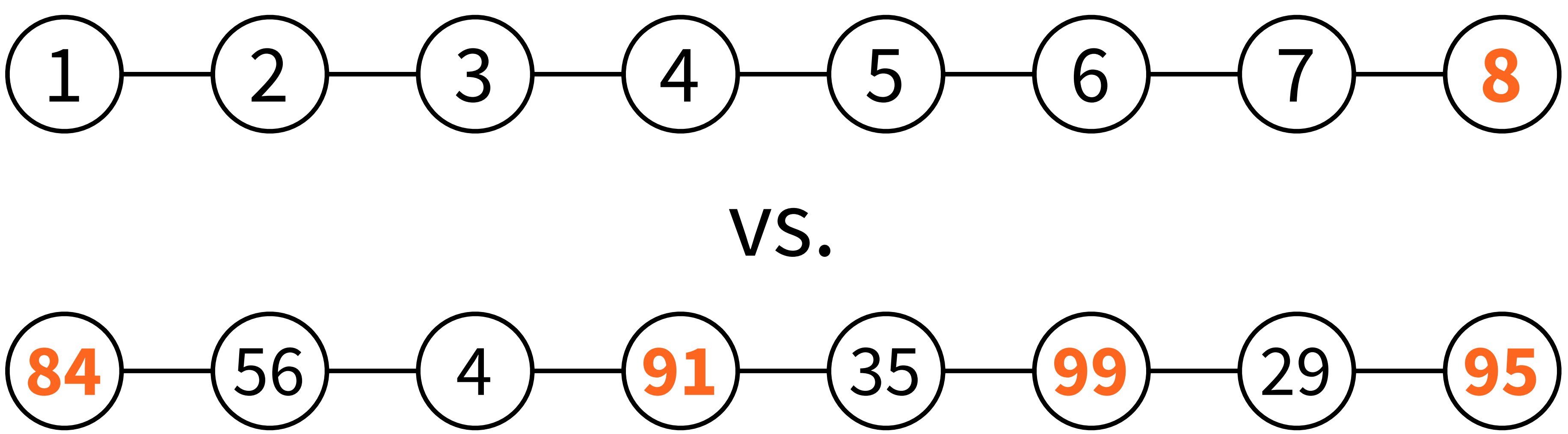
Uses of randomness

- **Break symmetry**
- **Similar to unique identifiers**



Uses of randomness

- **Better than unique identifiers: worst-case inputs unlikely?**

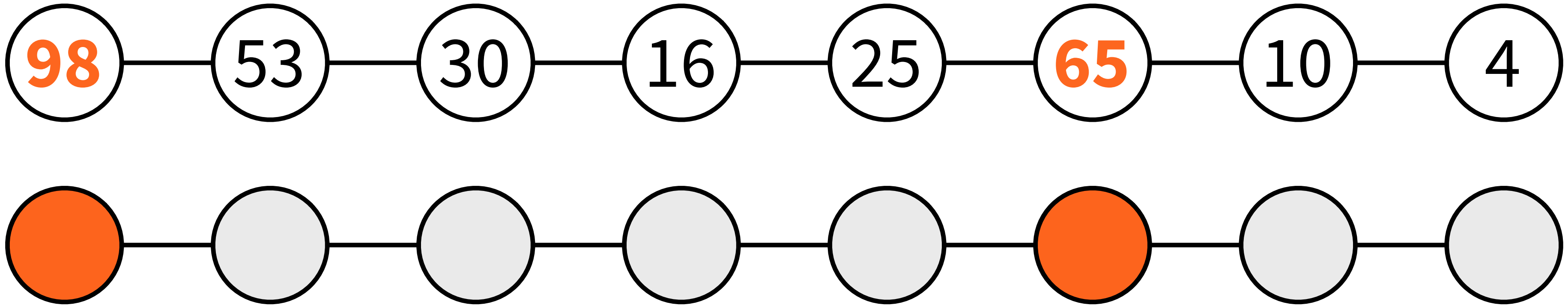


Guarantees

- ***Monte Carlo*: always fast**
  - running time deterministic
  - quality of output probabilistic
- ***Las Vegas*: always correct**
  - running time probabilistic
  - quality of output deterministic

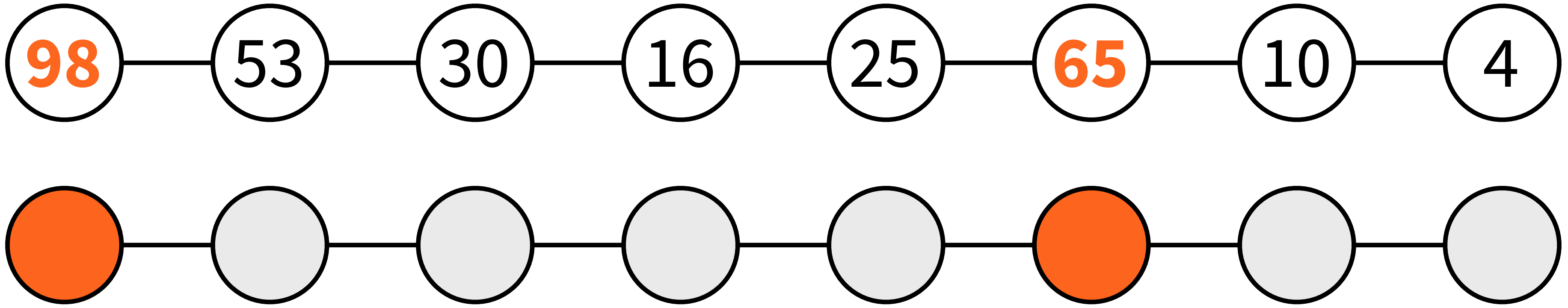
Monte Carlo

- Example: large independent set
- Pick random values, local maxima join



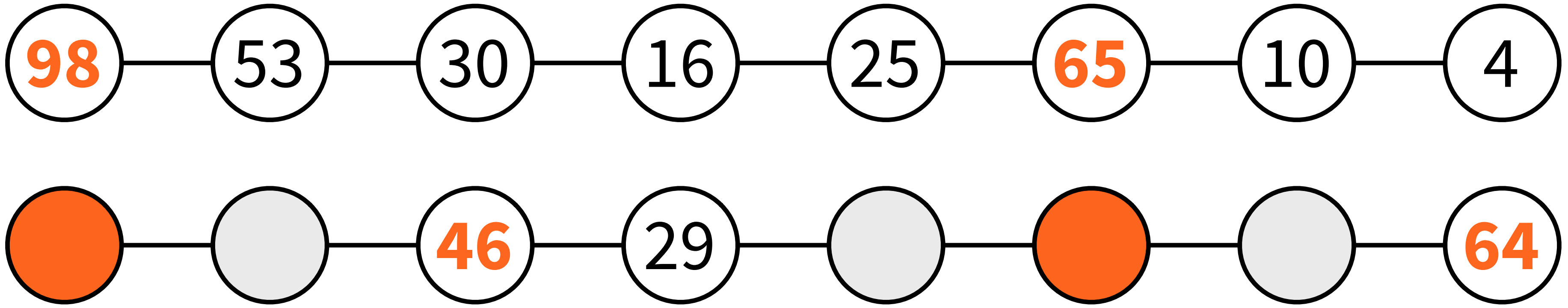
Monte Carlo

- Running time always $O(1)$
- Size of the set depends on random values



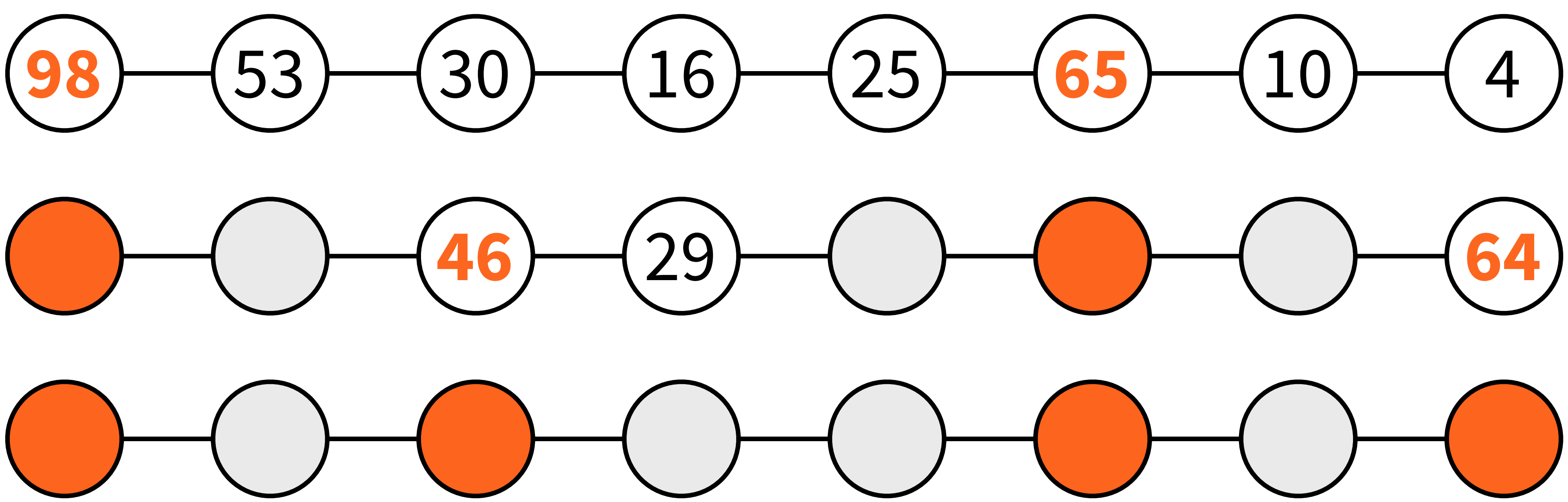
Las Vegas

- Pick random values, local maxima join, …



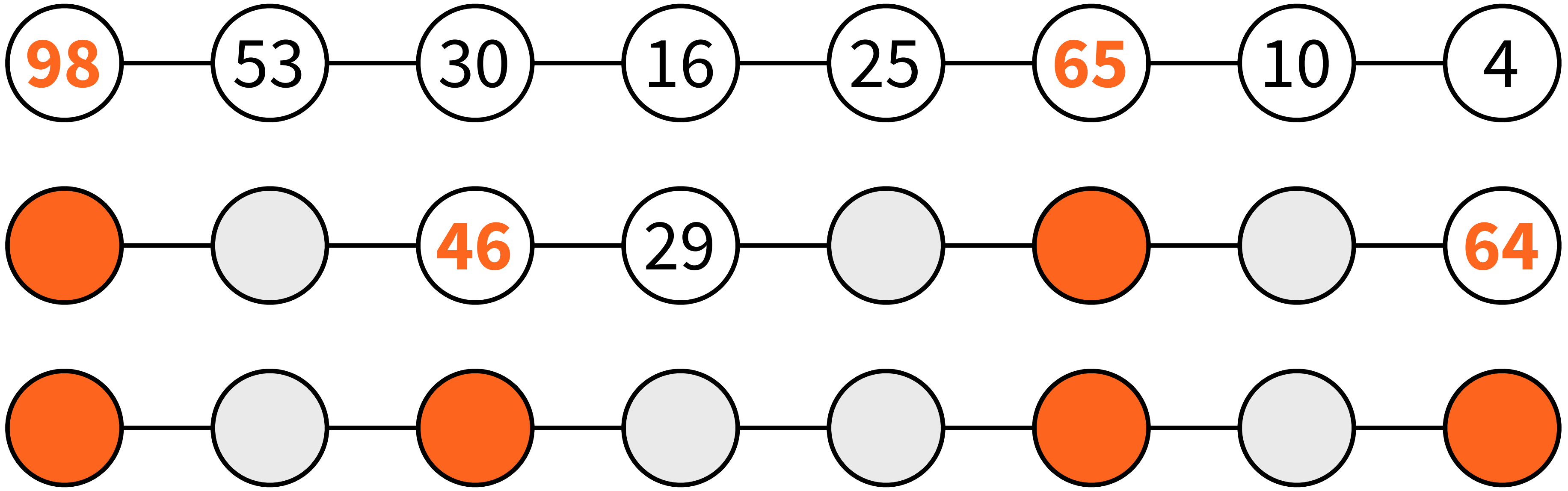
Las Vegas

- **Pick random values, local maxima join, repeat until maximal**



Las Vegas

- **Output is always maximal independent set, running time probabilistic**



With high probability

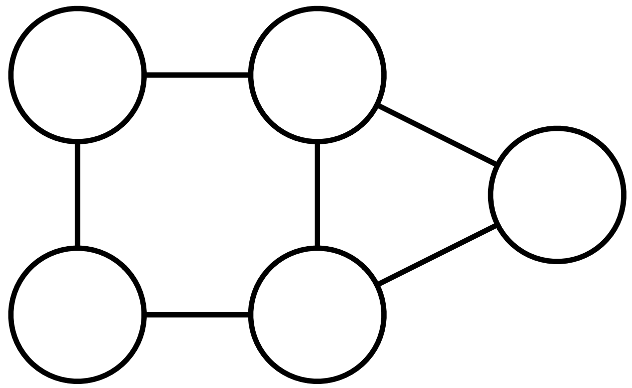
- **Success probability 1 – 1/n^c**
  - I can choose any constant c
- **"algorithm *A* stops in time *O*(log *n*) with high probability"**
- **"running time is *O*(log *n*) w.h.p."**

Example:
Graph colouring

- **Chapter 5: deterministic algorithm, $(\Delta + 1)$-colouring in $O(\Delta^2 + \log^* n)$ rounds**
- **Today: randomised algorithm, $(\Delta + 1)$-colouring in $O(\log n)$ rounds w.h.p.**

Algorithm idea

- Colour palette: $\{1, 2, \ldots, \Delta + 1\}$
- Pick a random colour
- Try again if conflicts...



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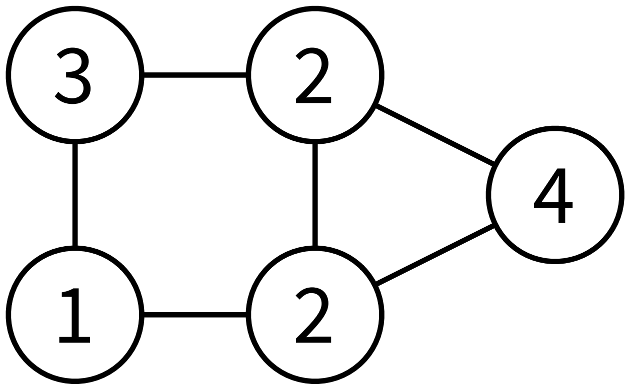

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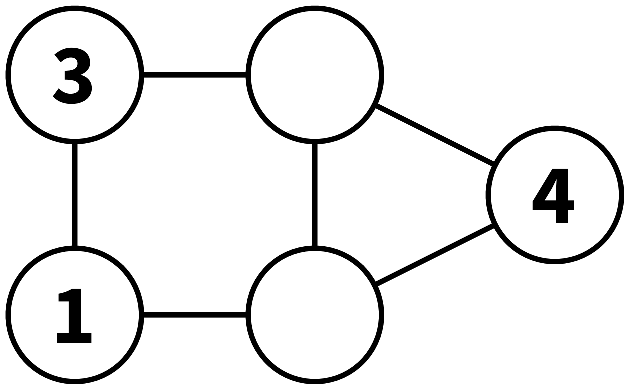

Algorithm idea

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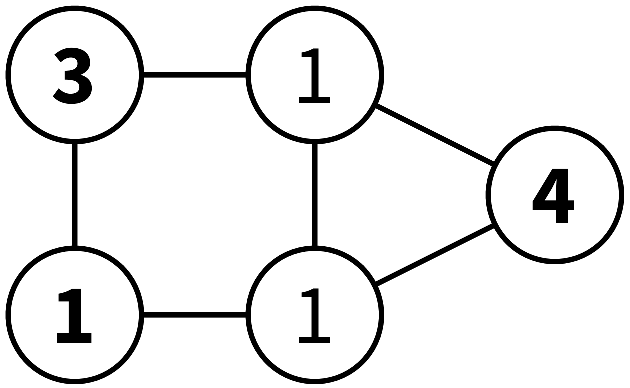

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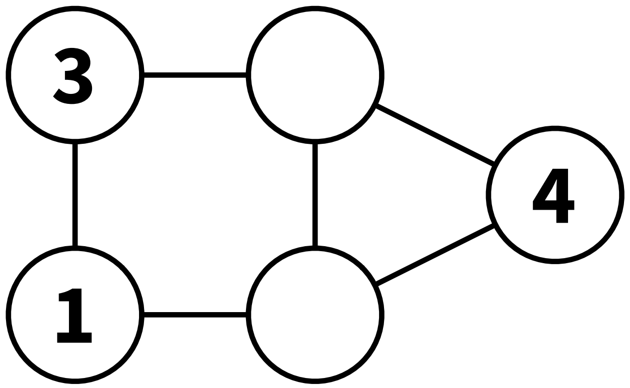

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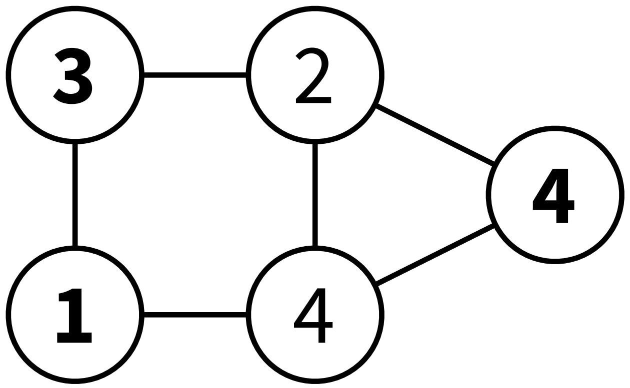

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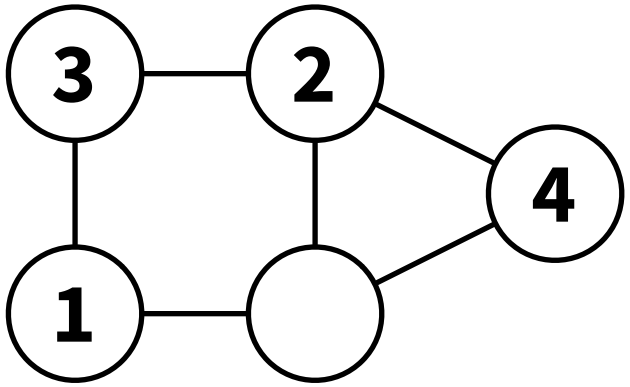

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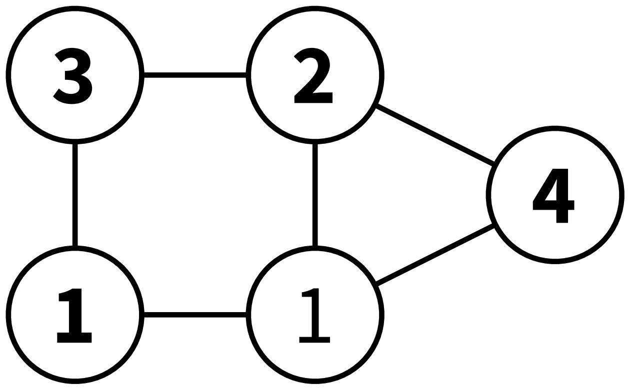

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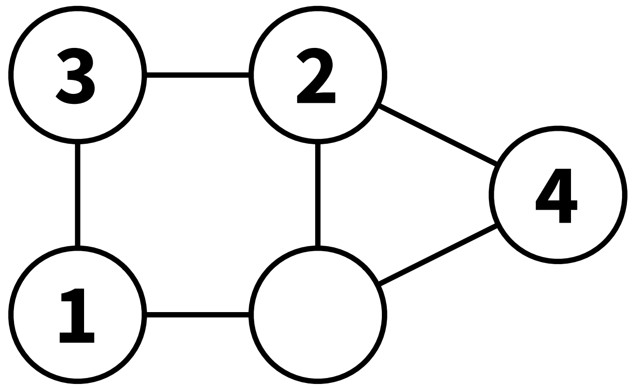

Algorithm idea

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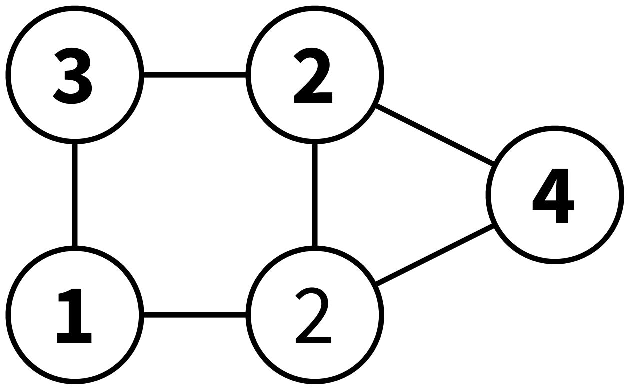

Algorithm idea

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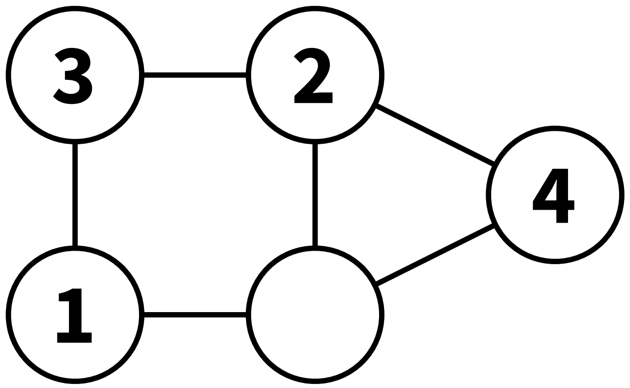
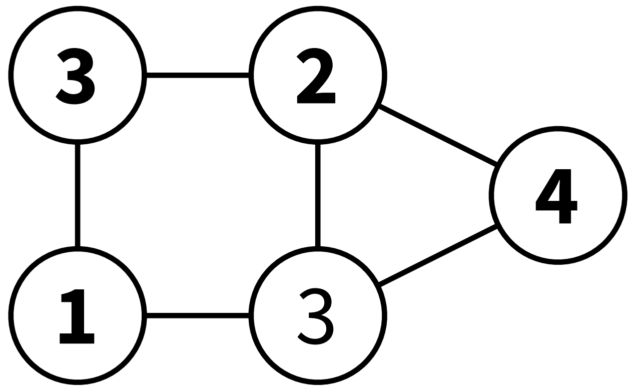


Algorithm idea

- Colour palette: $\{1, 2, \dots, \Delta + 1\}$
- Pick a random colour
- Try again if conflicts…



Algorithm idea

- Colour palette: $\{1, 2, \ldots, \Delta + 1\}$
- Pick a random colour
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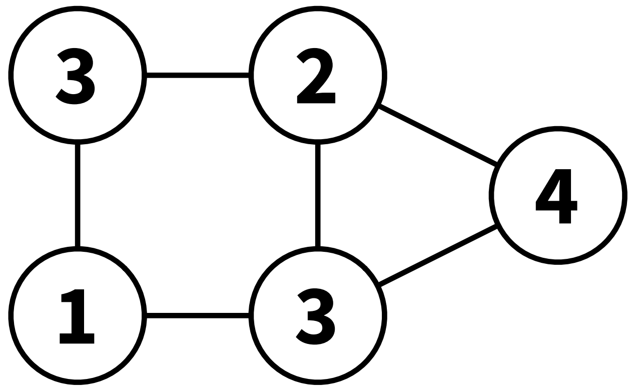


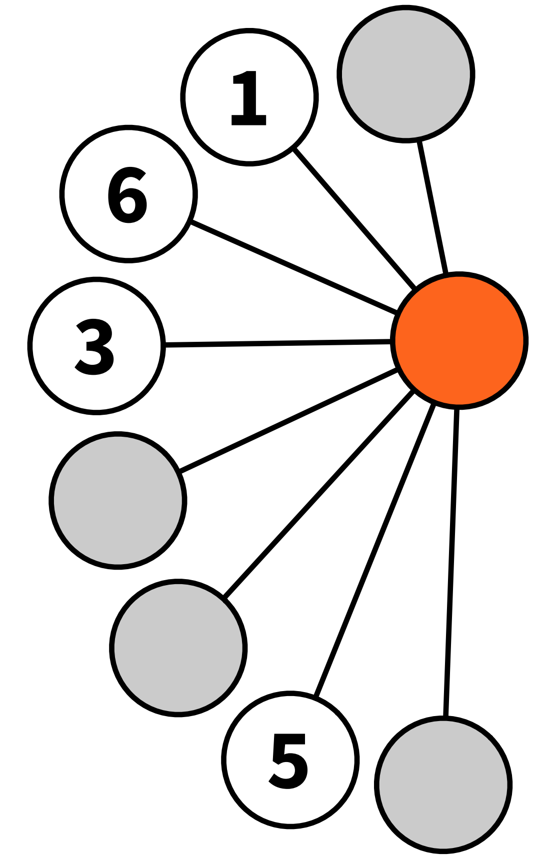
Algorithm idea 2

- **Colour palette: {1, 2, ..., Δ + 1}**
- **Pick a random *free* colour**
  - not used by any neighbour that has stopped
- **Try again if conflicts...**

Algorithm idea 3

- **Colour palette: $\{1, 2, \ldots, \Delta + 1\}$**
- **Active with probability 1/2**
- **If *active*, pick a random *free* colour**
  - not used by any neighbour that has stopped
- **Try again if conflicts…**

Algorithm idea 3



- **Active with probability 1/2**
- **Intuition:**
  - assume: d neighbours still running
  - roughly $d/2$ of them active
  - at least $d + 1$ colours in my palette
  - easy to pick a colour without conflicts (?)

- $s = 1$, $c \neq \bot$:
  - stopping state; **output c**
- $s = 1$, $c = \bot$:
  - probability 1/2: $c \leftarrow \bot$
  - probability 1/2: $c \leftarrow$ **random free colour**
  - $s \leftarrow 0$
- $s = 0$:
  - if conflicts: $c \leftarrow \bot$
  - $s \leftarrow 1$

Algorithm DBRand: **Randomised colouring**

- **Lemma 1:** A running node succeeds with probability 1/4

Algorithm DBRand: **Randomised colouring**

- **Lemma 1:** A running node succeeds with probability 1/4
- $\boldsymbol{T(n) = 2(c + 1) \log_{4/3} n = O(\log n)}$
- **Lemma 2:** For any v, node v stops in time $T(n)$ with probability $\mathbf{1 - 1/n^{c+1}}$

Algorithm DBRand:
Randomised colouring

- $T(n) = 2(c+1)\log_{4/3} n = O(\log n)$
- **Lemma 2:** For any *v*, node *v* stops in time $T(n)$ with probability $1 - 1/n^{c+1}$
- **Theorem 3:** All nodes stop in time $T(n)$ with probability $1 - 1/n^{c}$

Summary

- **Randomness may help**
- **Common idea: each node makes random trials until successful**
- **Typical running time: $O(\log n)$ w.h.p.**
  - proof technique: in each round, each node successful with some constant probability

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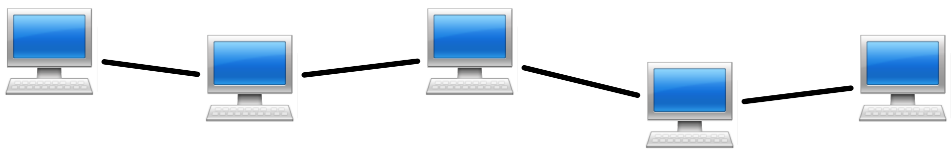
    - network = path 
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