

Distributed Algorithms 2024



LOCAL model port-numbering model + unique identifiers

Nodes have distinct labels from {1, 2, ..., poly(*n*)}

CONGEST model LOCAL model + bandwidth limitation

Messages at most $O(\log n)$ bits

LOCAL · unbounded messages

• everything trivial to solve in O(diameter) rounds: gather full input and solve locally

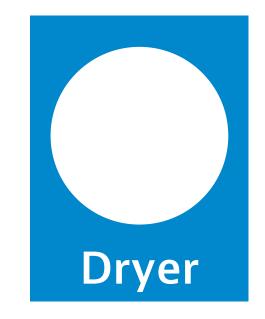
CONGEST · bounded messages

- gathering everything is way too expensive
- O(diameter) and O(n) is nontrivial

Designing efficient algorithms in **CONGEST** model

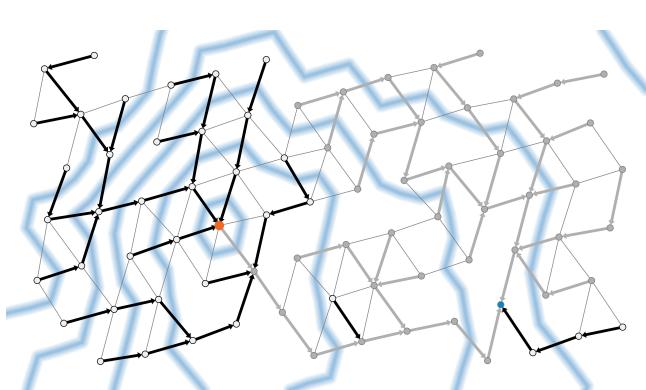
Pipelining





Pipelining

- Multiple operations in progress *simultaneously*
- Using *different resources*
- In APSP algorithm:
 - multiple waves
 - using different communication links



Pipelining

- Does not reduce the total number of messages
 - only removes idle periods between messages
- If all communication links are already sending useful data every round, no room for pipelining

What kind of problems cannot be solved fast in **CONGEST** model?

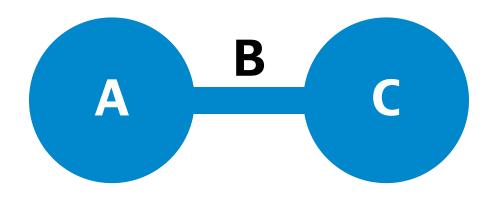
Typical hard problems

•A: complicated, lots of information

• B: bottleneck

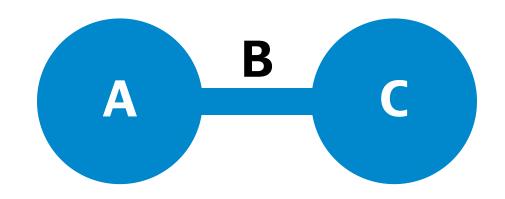
 can only send small number of bits per round from A to C

•C: need to know A



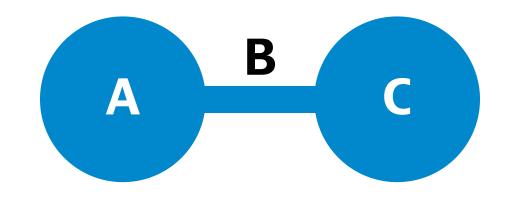
Proving hardness

- Counting argument
- Many possible inputs in A
- Few possible messages across bottleneck B



Proving hardness

- Counting argument
- Many possible inputs in A
- Few possible messages across bottleneck B
- Contradiction:
 - *different* inputs in A
 - same messages across B
 - same output in C





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Randomized algorithms



Deterministic algorithms in PN model init_d(...), send_d(...), receive_d(...)

- Deterministic algorithms in LOCAL model
 add unique identifiers
- Deterministic algorithms in CONGEST model
 - add bandwidth constraints

Randomized algorithms

- Randomized algorithms in PN model
 - init_d(...), receive_d(...): probability distribution
- Randomized algorithms in LOCAL model
 add unique identifiers
- Randomized algorithms in CONGEST model
 - add bandwidth constraints

Guarantees

Monte Carlo

- guaranteed running time
- probabilistic output quality

• Las Vegas

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Monte Carlo

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- probabilistic running time
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• "With high probability" (w.h.p.)

Role of randomness

- Sometimes randomness is the only way to design fast distributed algorithms
- Example: sinkless orientation
 - deterministic LOCAL: **O(log n)** is best possible
 - randomized LOCAL: O(log log n) w.h.p. is best possible

Role of randomness

- Sometimes randomness is just one of many ways to break symmetry
- Example:
 - **PN model** + randomness + knowledge of *n*: you can construct **unique identifiers** w.h.p.



Pretty simple idea:

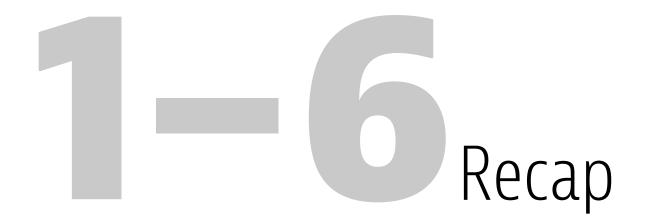
- nodes are *active* with probability 1/2
- only active nodes try to pick a random free color
- stop if successful

Simplest possible idea:

- everyone tries to pick
 a random free color
- stop if successful



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Models of computing

- PN
- **LOCAL** unique identifiers
- **CONGEST** bandwidth constraints
- Deterministic and **randomized** algorithms

Canonical problems

Vertex coloring

- coloring = schedule coloring breaks symmetry
- Used to solve many other problems
- Coming later: used to show that other problems are hard
- Demonstrates different algorithm design ideas

Algorithm ideas

- Conflict avoidance & coordination
- Process nodes by color classes
- Send proposals one by one
- Random subset of nodes is active
- Pipelining
- Algebraic techniques

New kinds of challenges

Unknown systems

algorithms that work in any network

Partial information

• making decisions based on local information

Parallelism

• many nodes act simultaneously