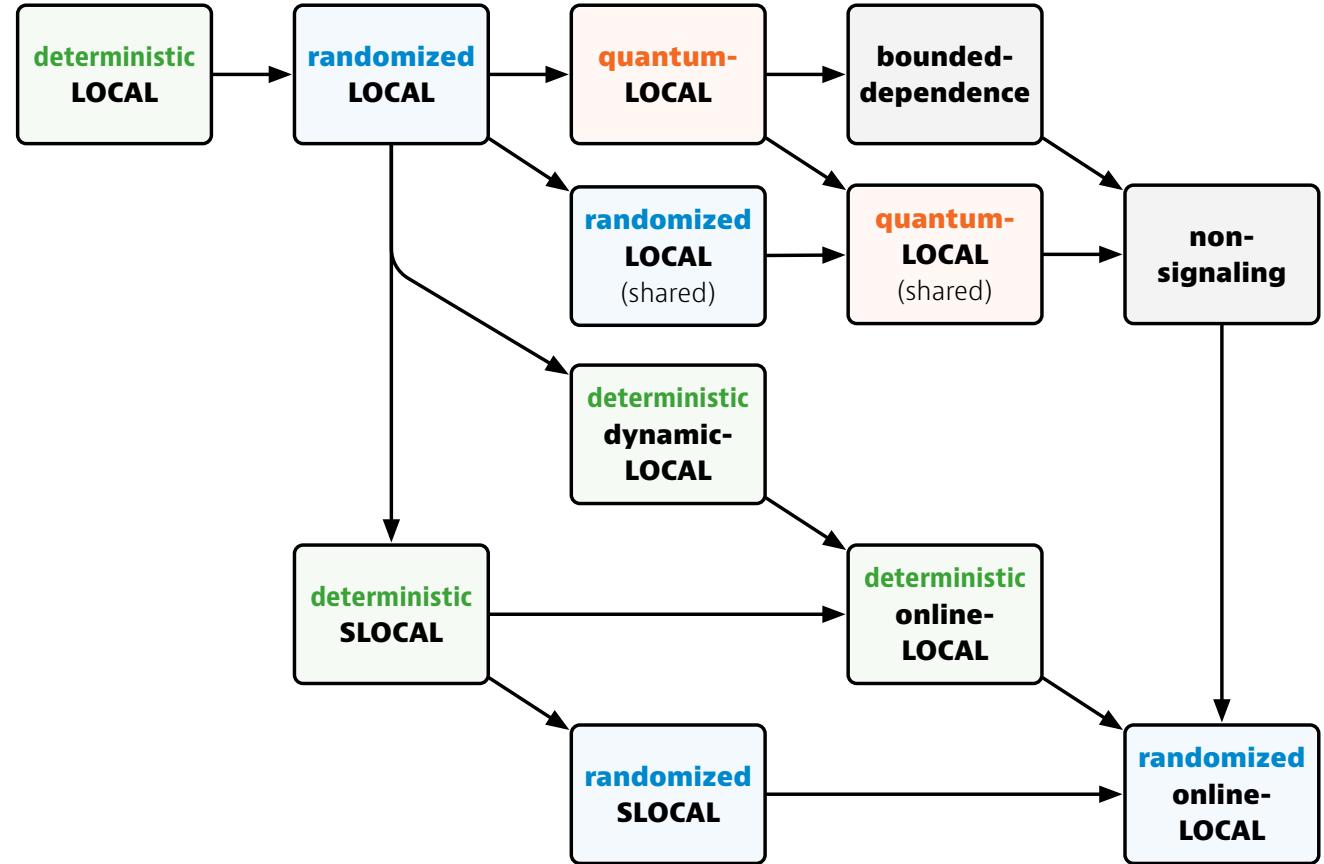


Locality in distributed, dynamic, online, and quantum settings

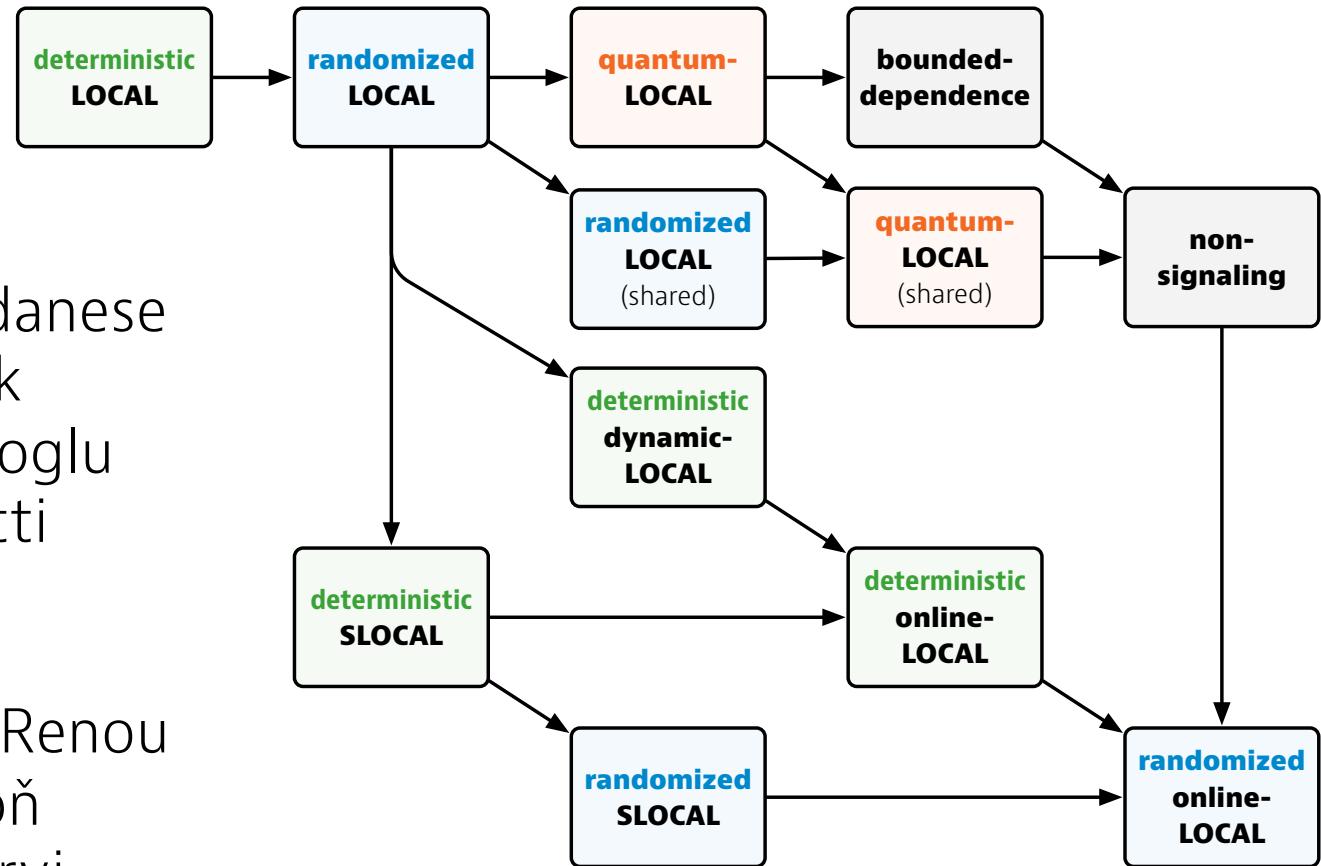


Jukka Suomela
Aalto University

Joint work with e.g.

Amirreza Akbari
Alkida Balliu
Sebastian Brandt
Xavier Coiteux-Roy
Francesco d'Amore
Anubhav Dhar
Massimo Equi
Navid Eslami
Rishikesh Gajjala
Mohsen Ghaffari
Fabian Kuhn
Eli Kujawa
François Le Gall
Henrik Lievonen

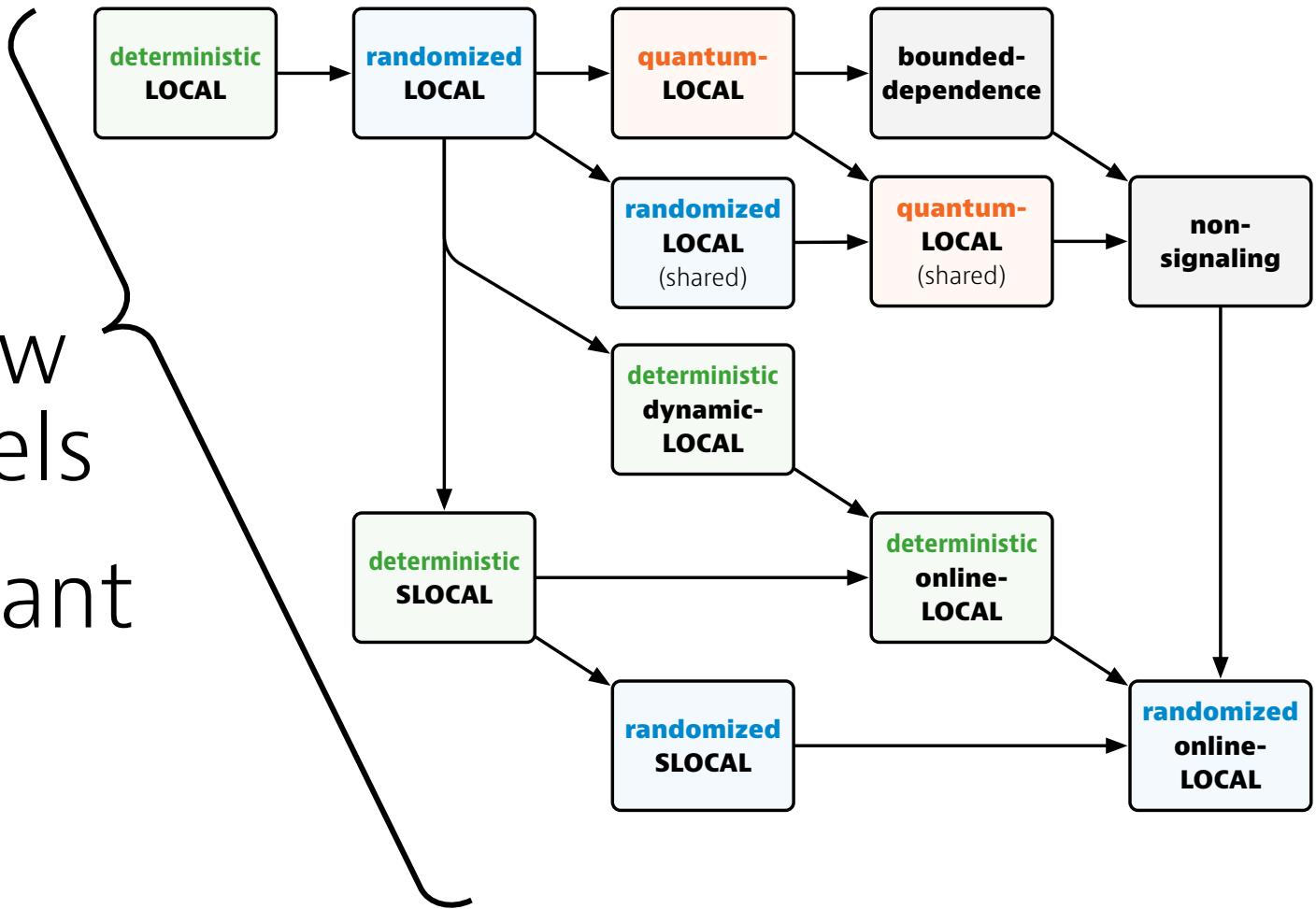
Augusto Modanese
Darya Melnyk
Mikail Muftuoglu
Dennis Olivetti
Shreyas Pai
Mikaël Rabie
Marc-Olivier Renou
Václav Rozhoň
Joonas Särkijärvi
Gustav Schmid
Jan Studený
Lucas Tendick
Jara Uitto
Isadora Veeren

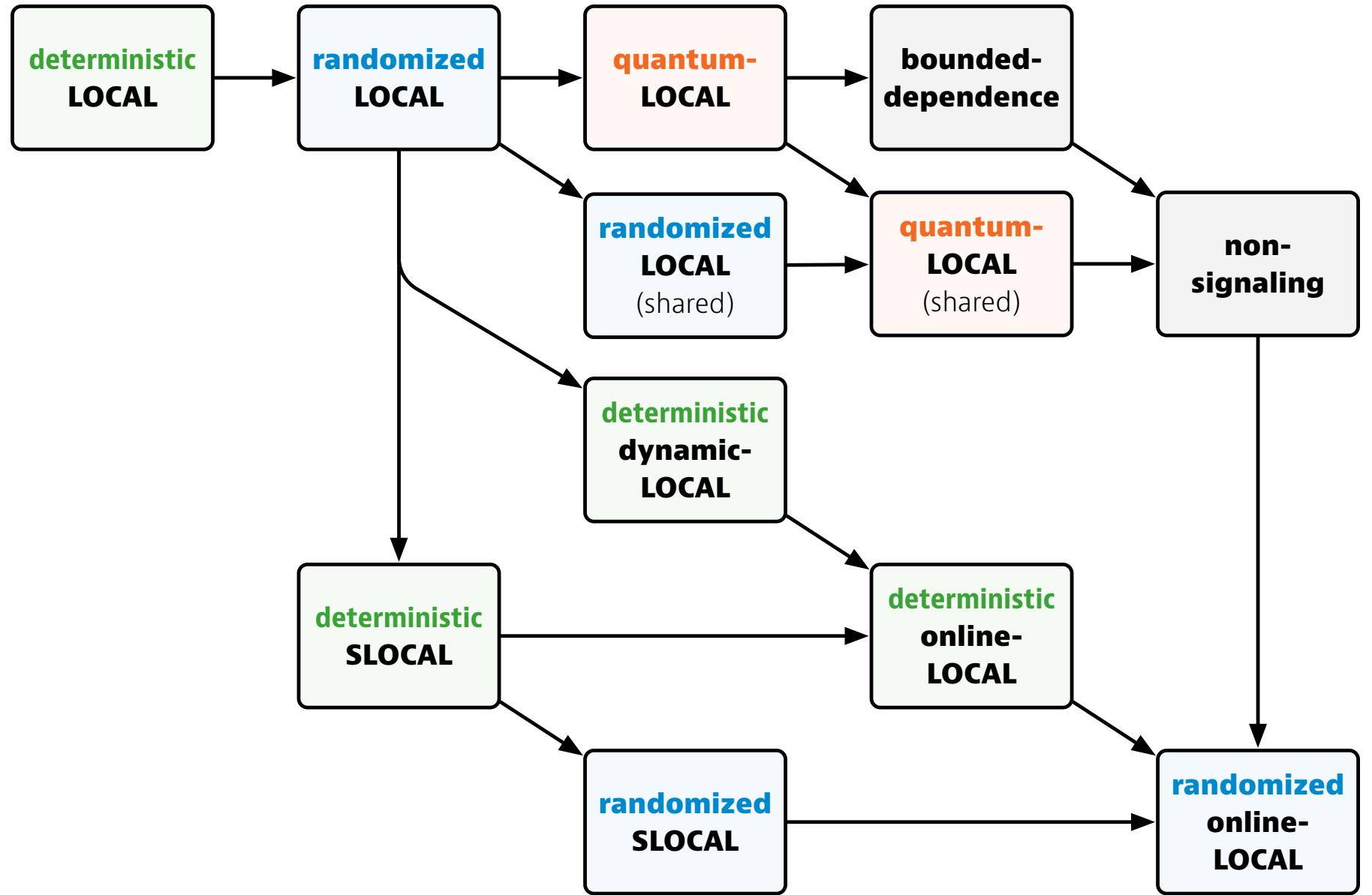


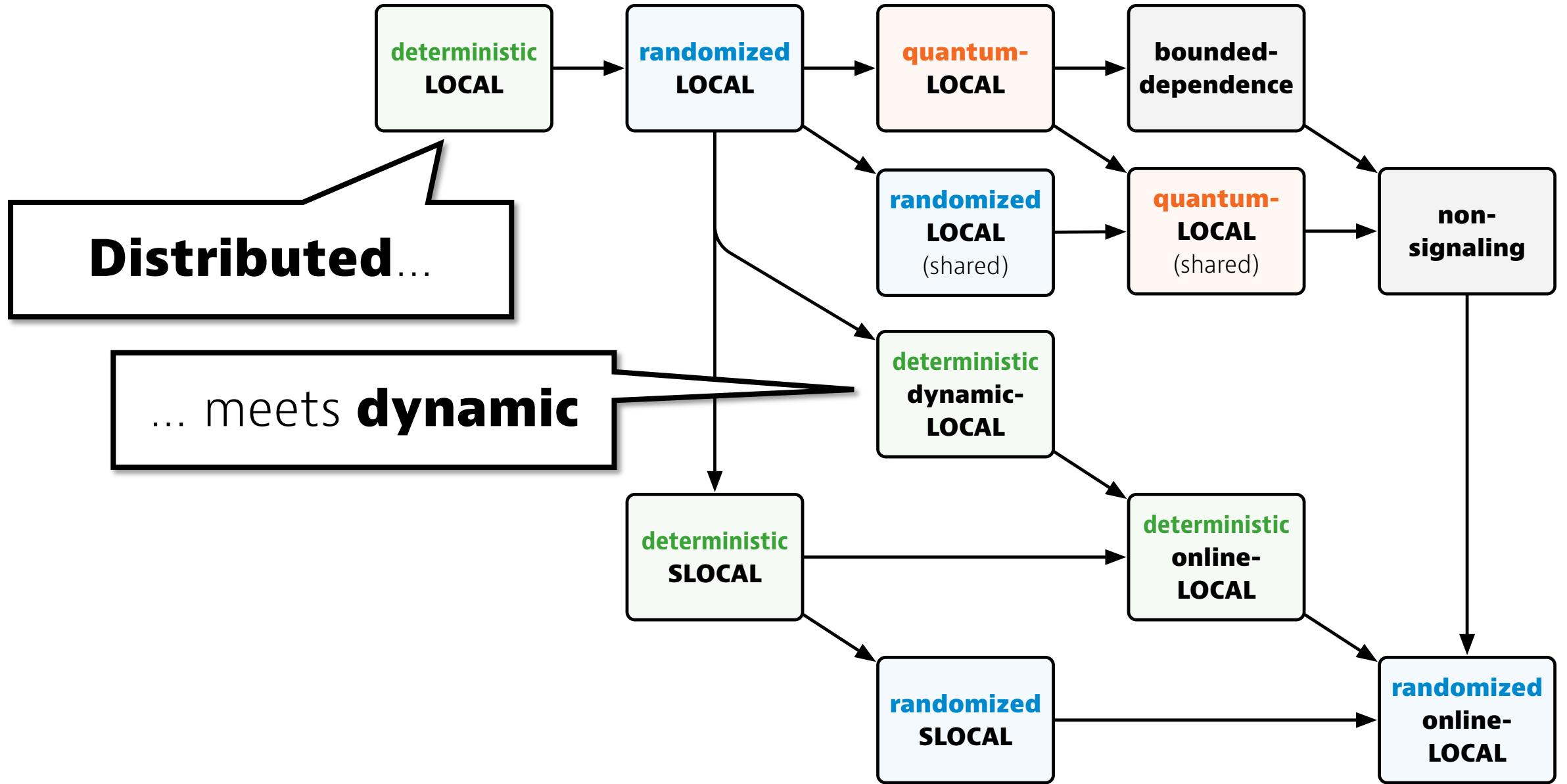
+ many others

Goals

1. Informal overview
of all these models
2. Why is this relevant
and interesting?
3. What is known,
what is open?

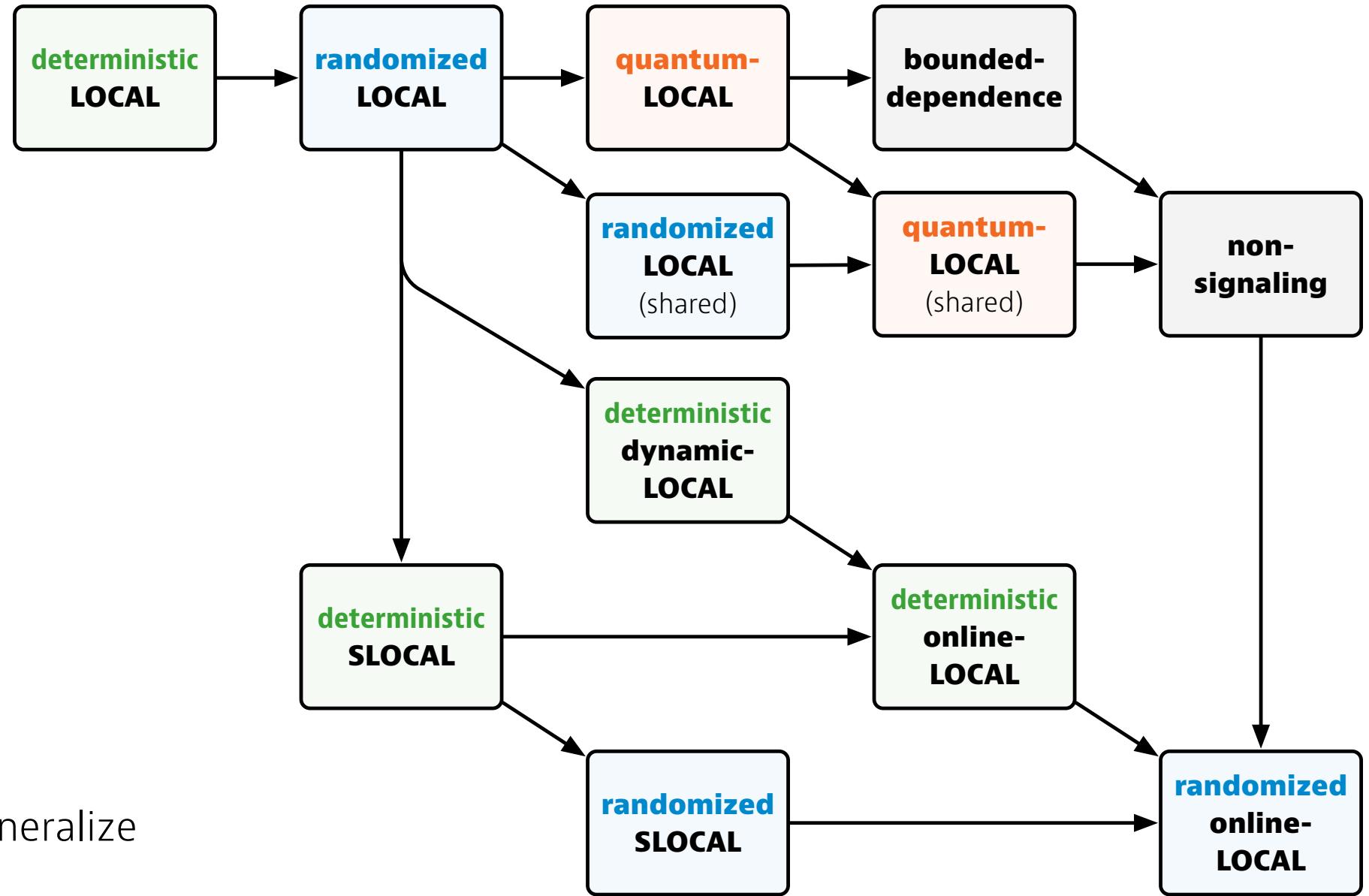




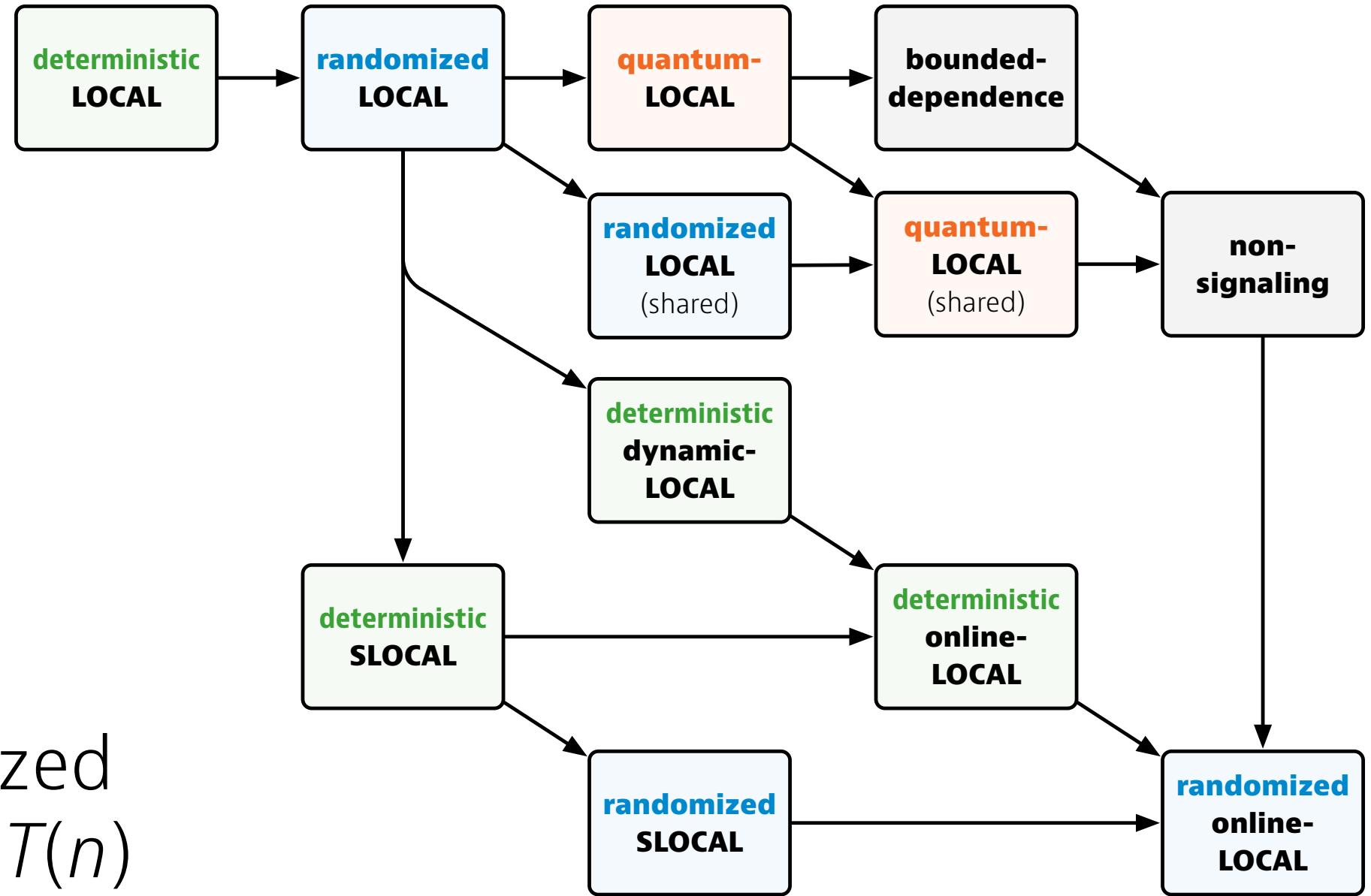


Context:
locally
checkable
labelings
(LCL)

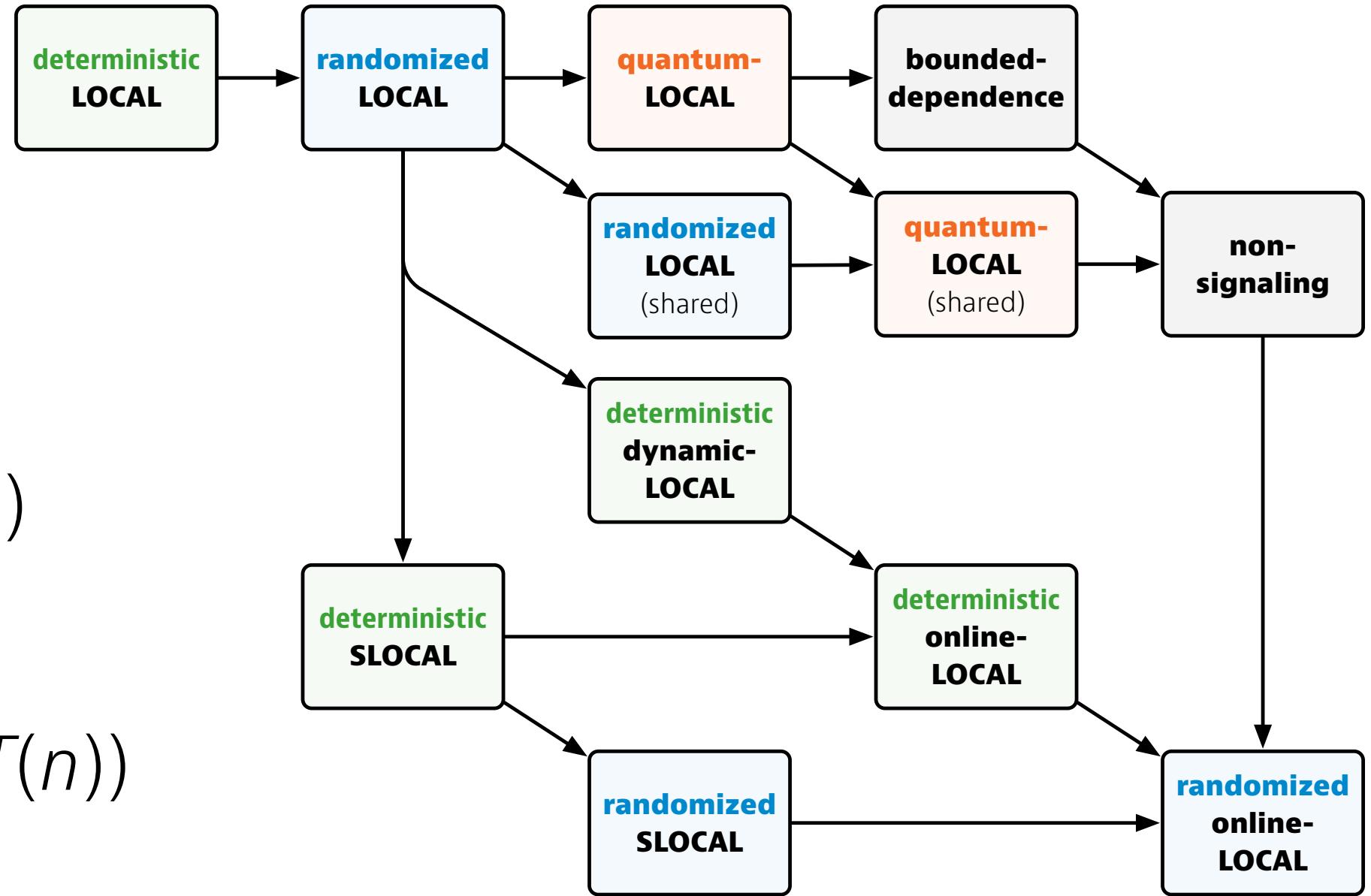
but many things generalize
far beyond LCLs



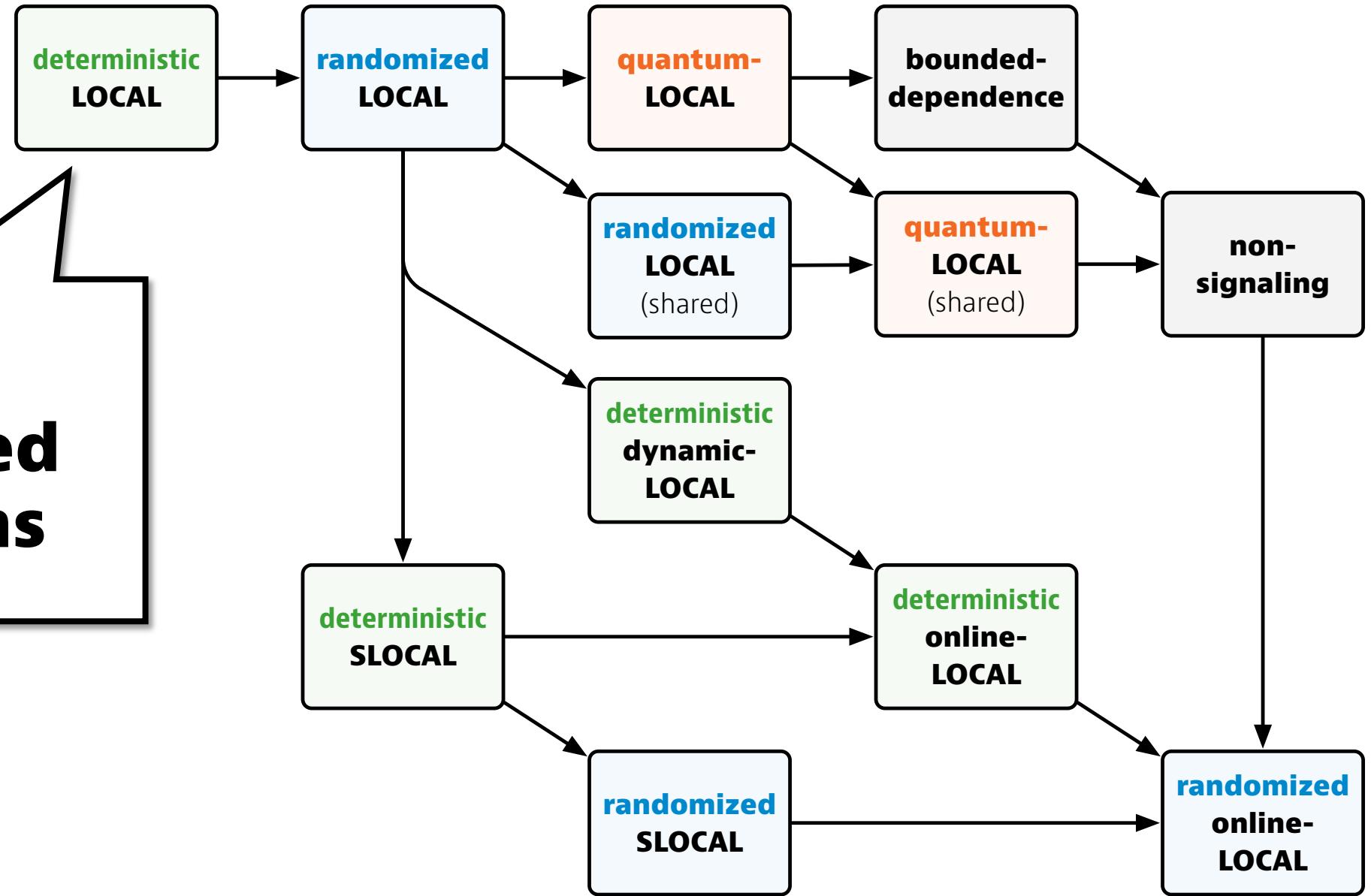
Each box:
model of computing
parameterized
by locality $T(n)$

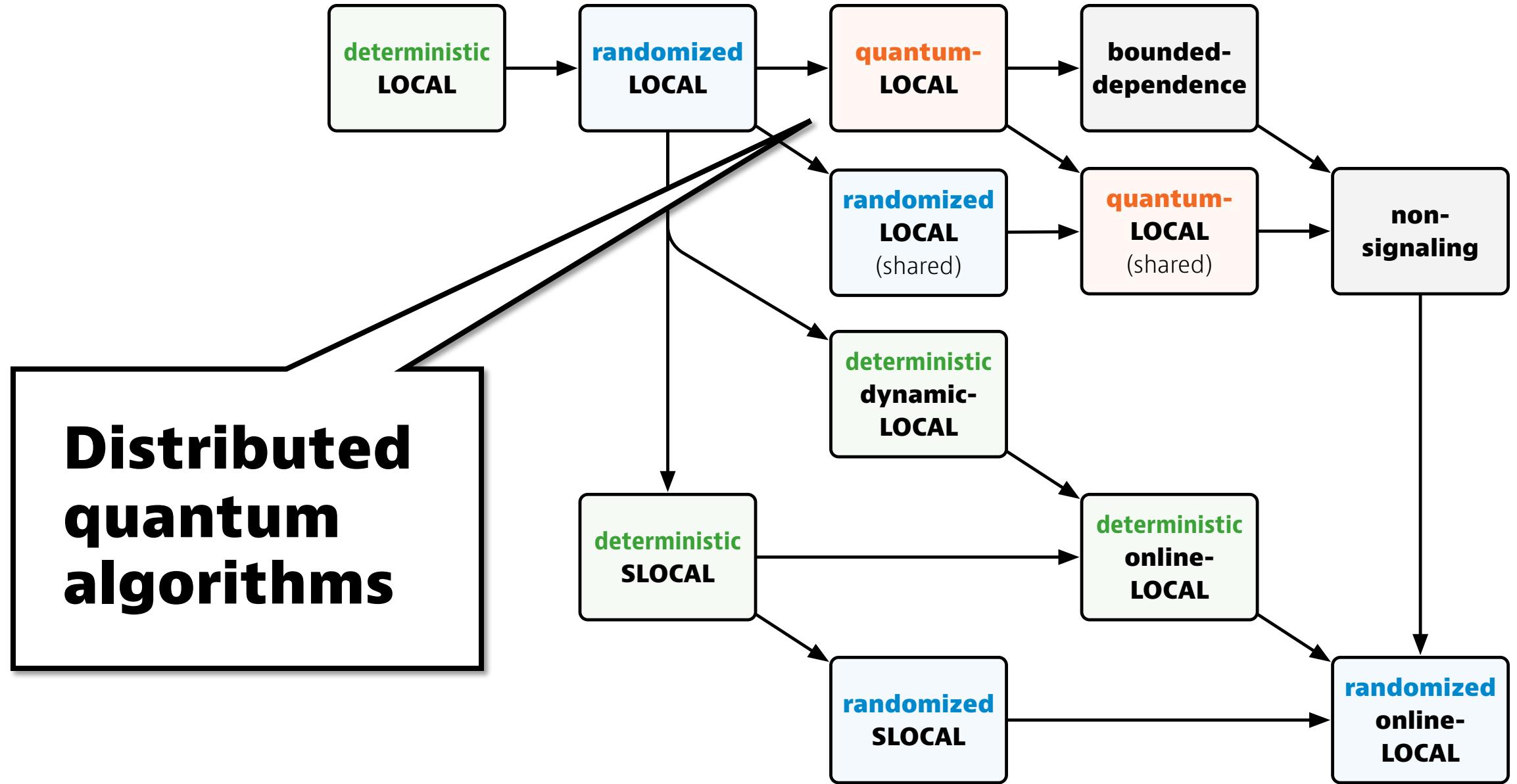


A → B:
locality $T(n)$
in model A
implies
locality $O(T(n))$
in model B

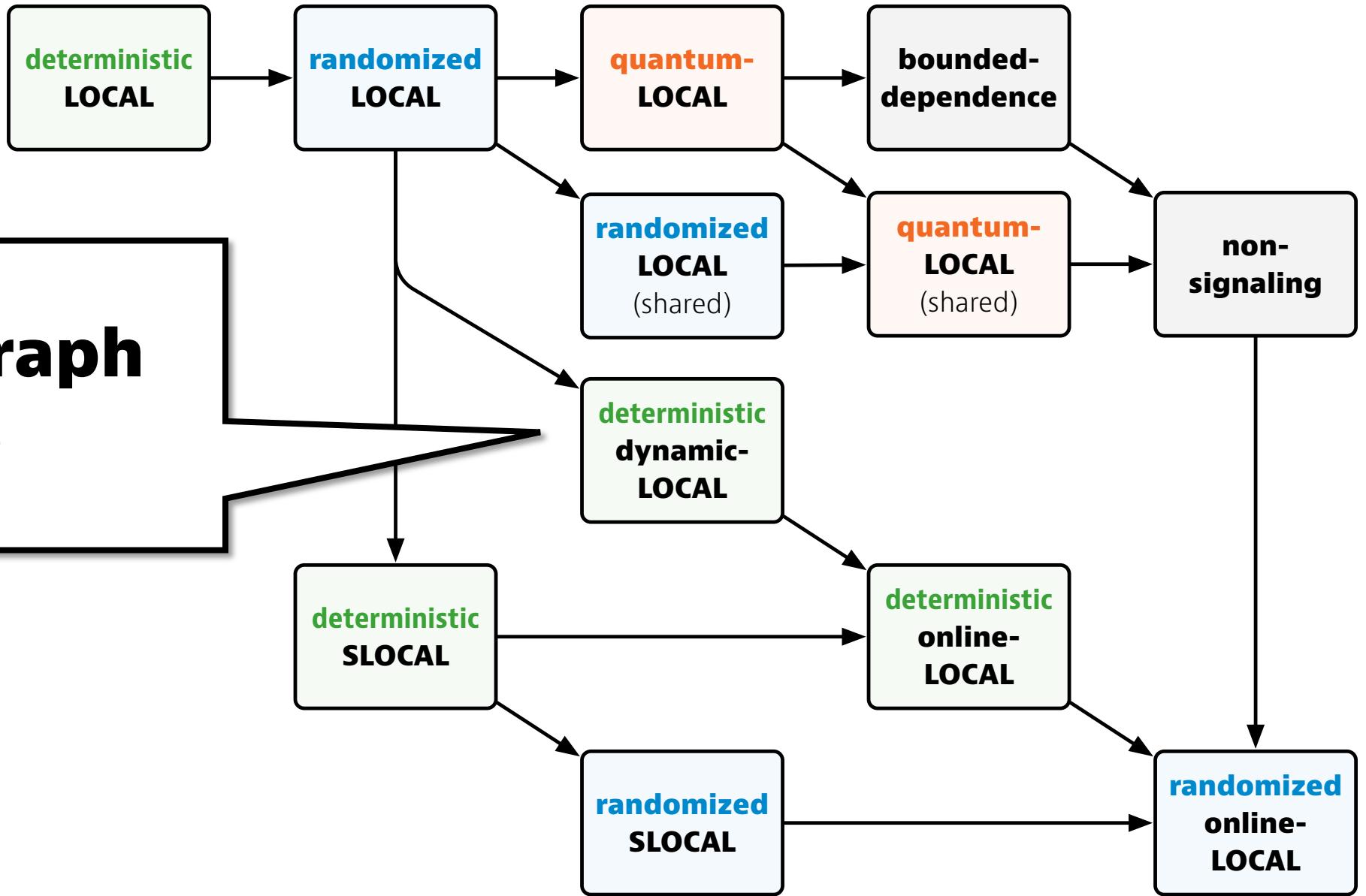


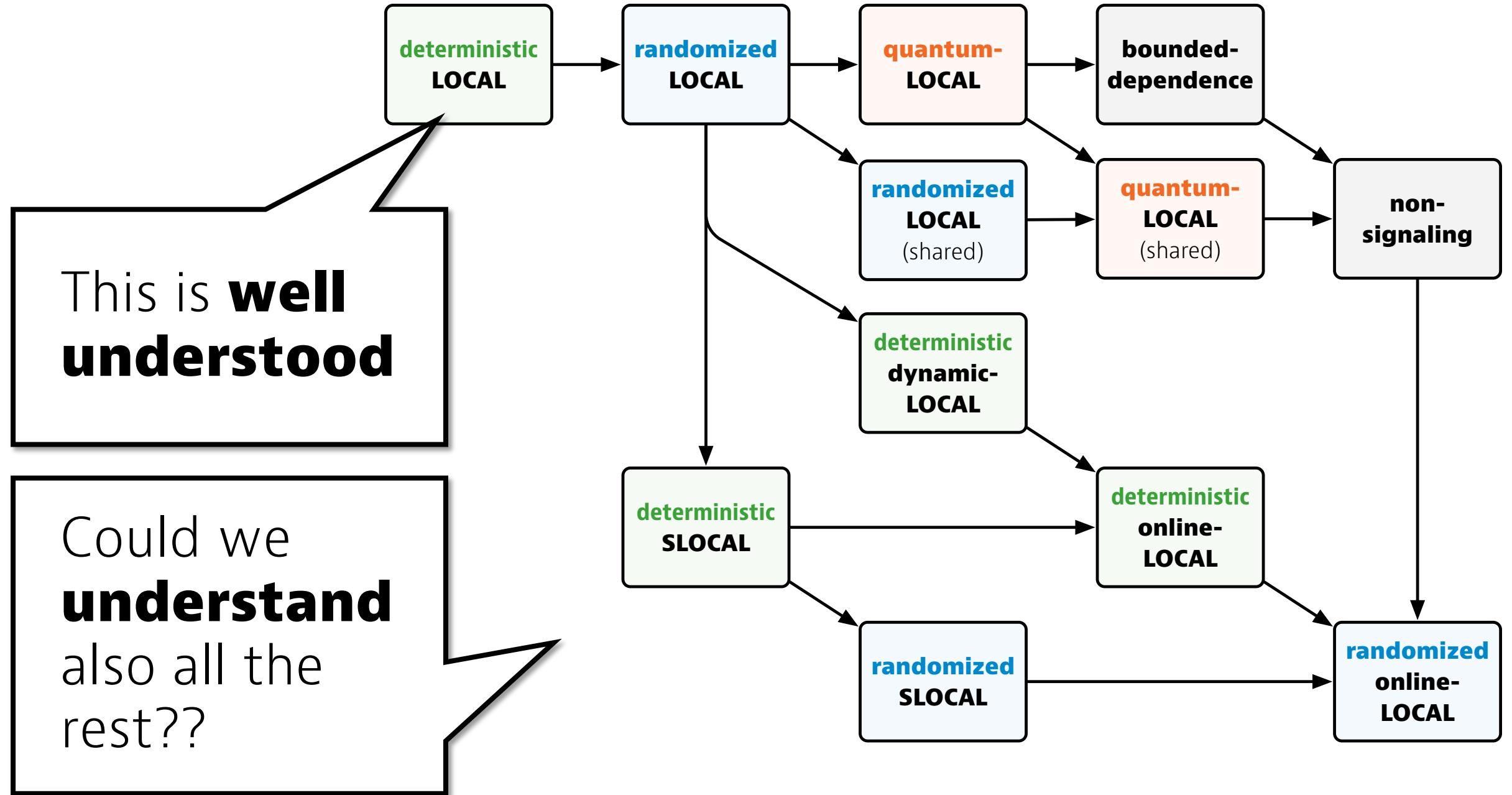
Classical distributed algorithms

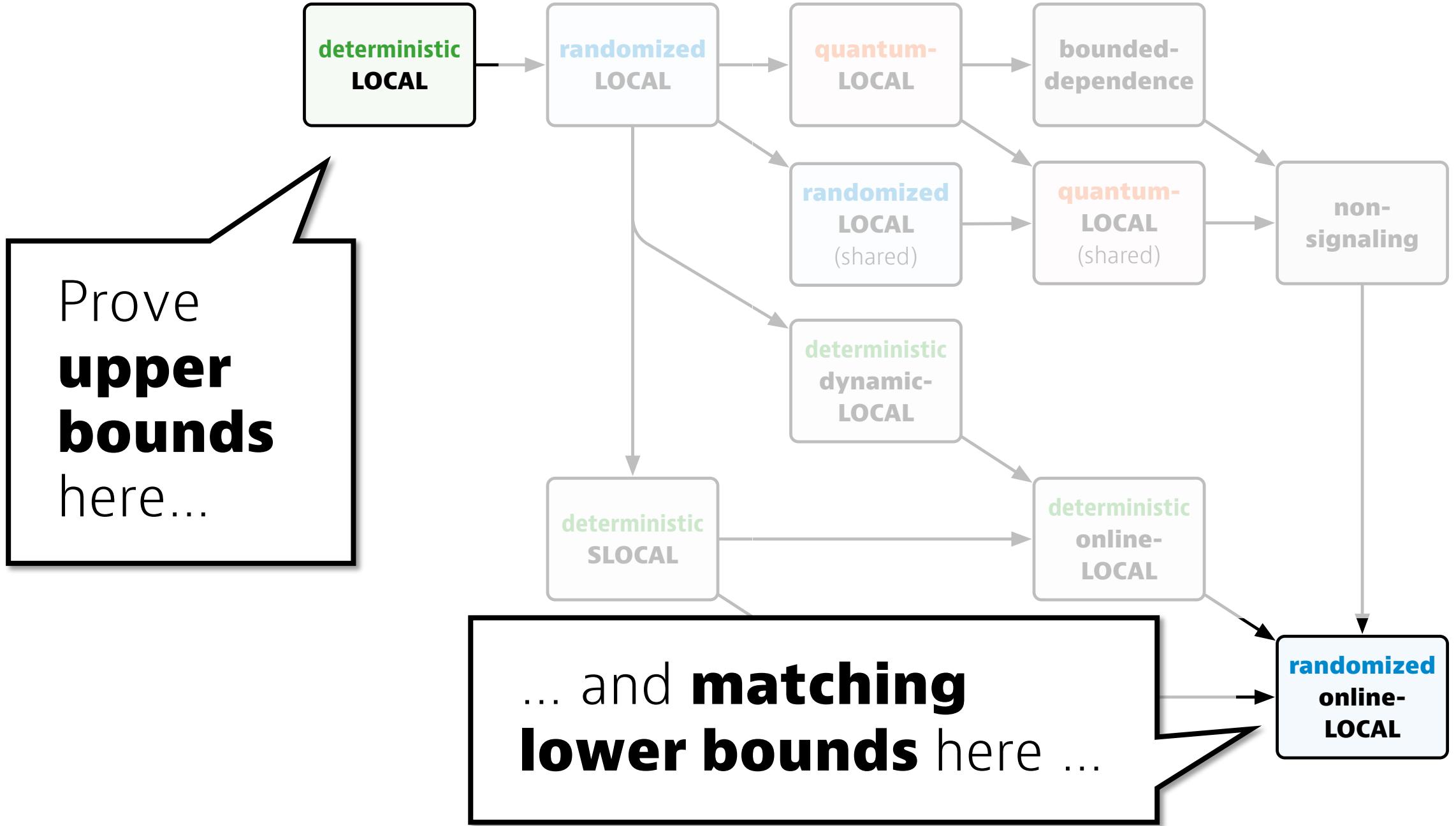




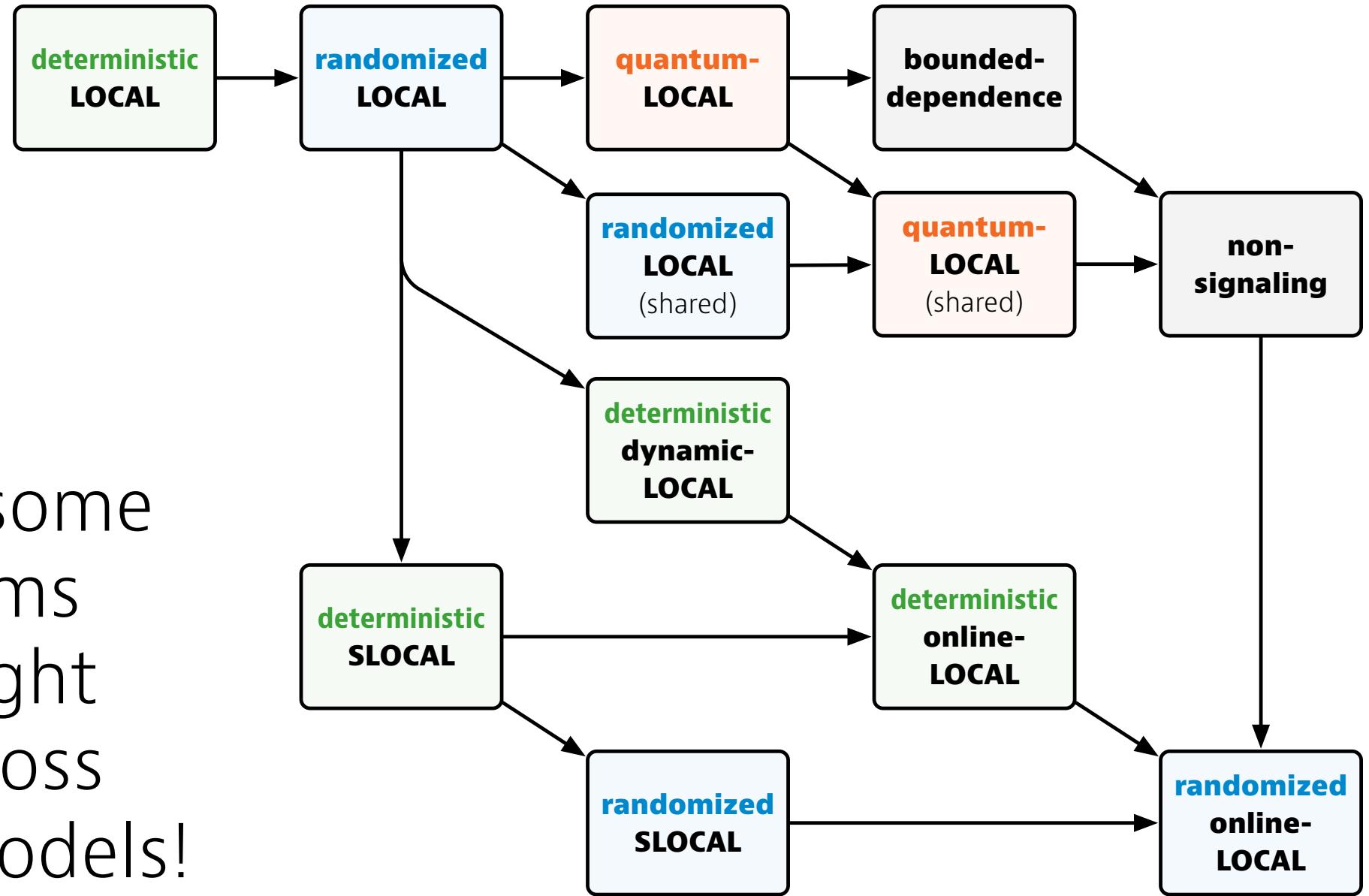
Dynamic graph algorithms

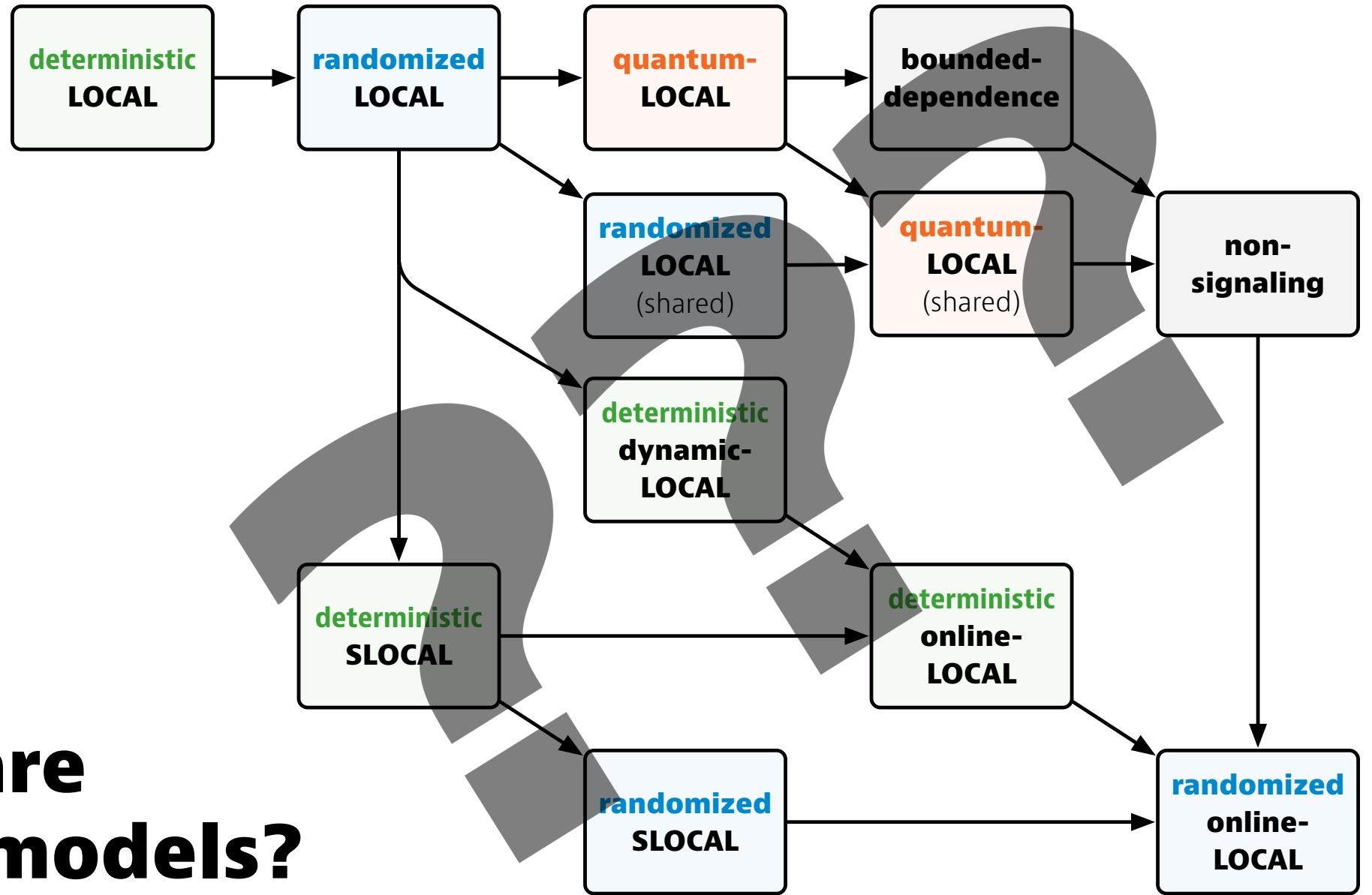




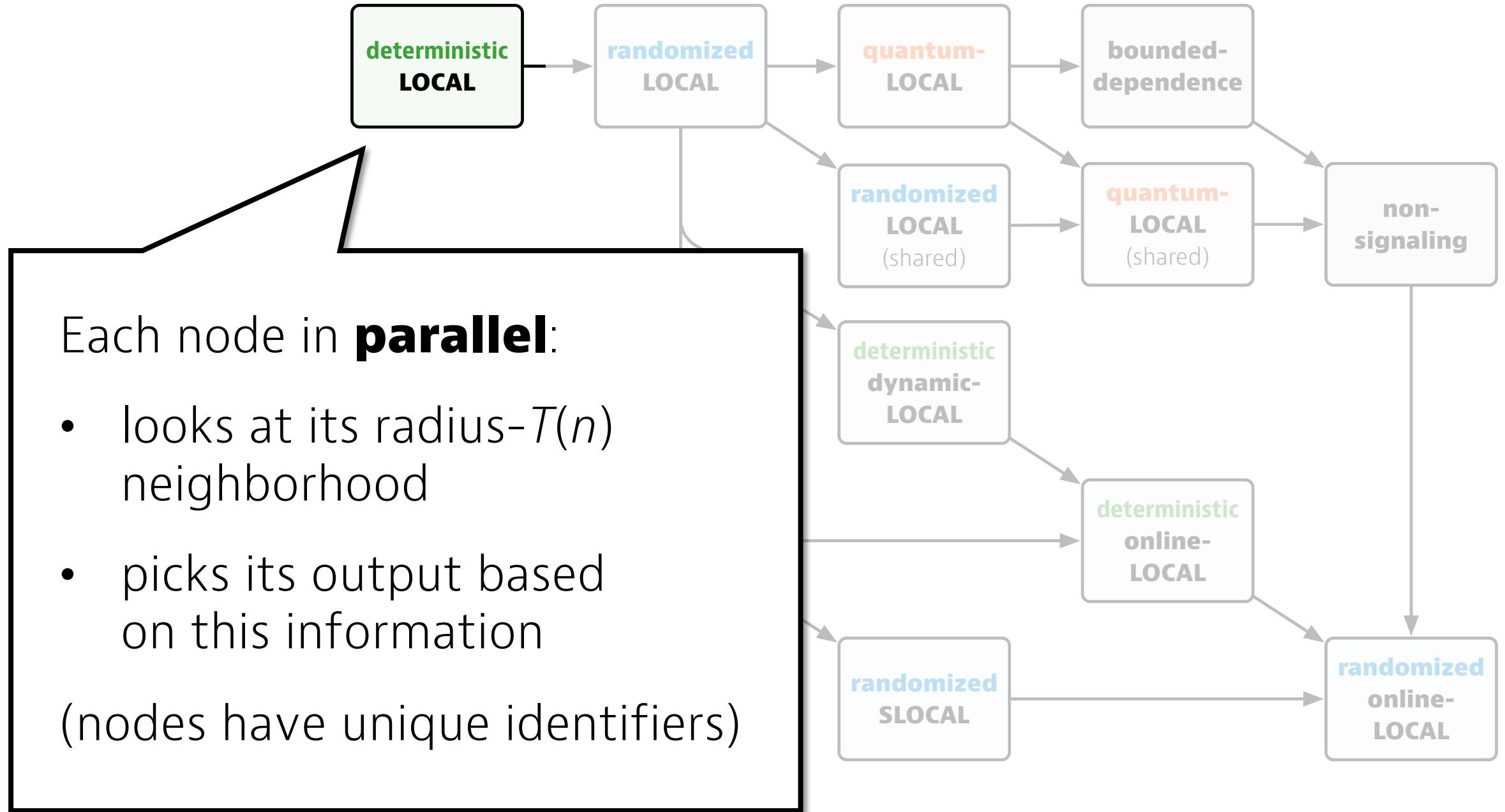


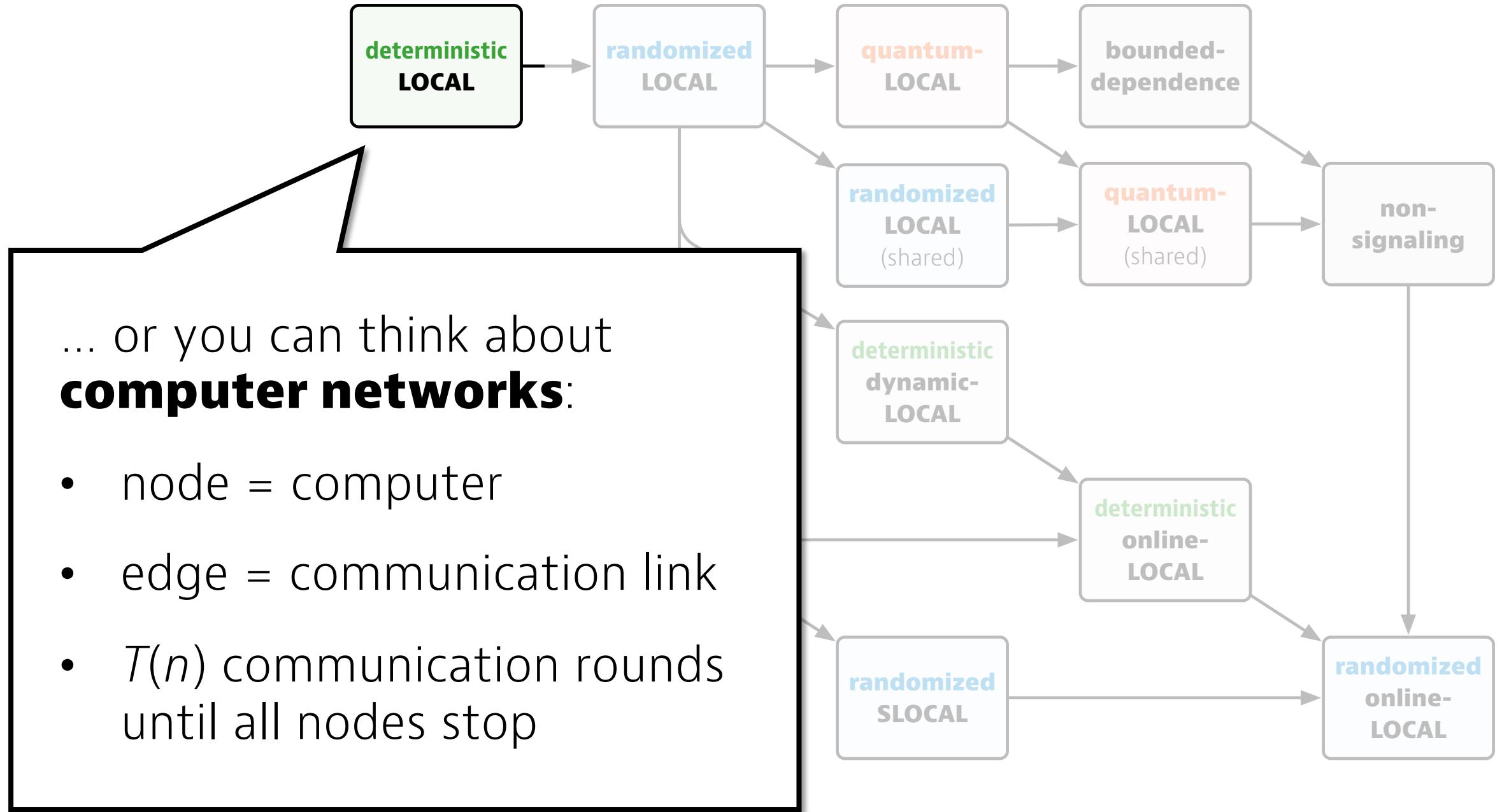
... and for some
LCL problems
we have tight
bounds across
all these models!

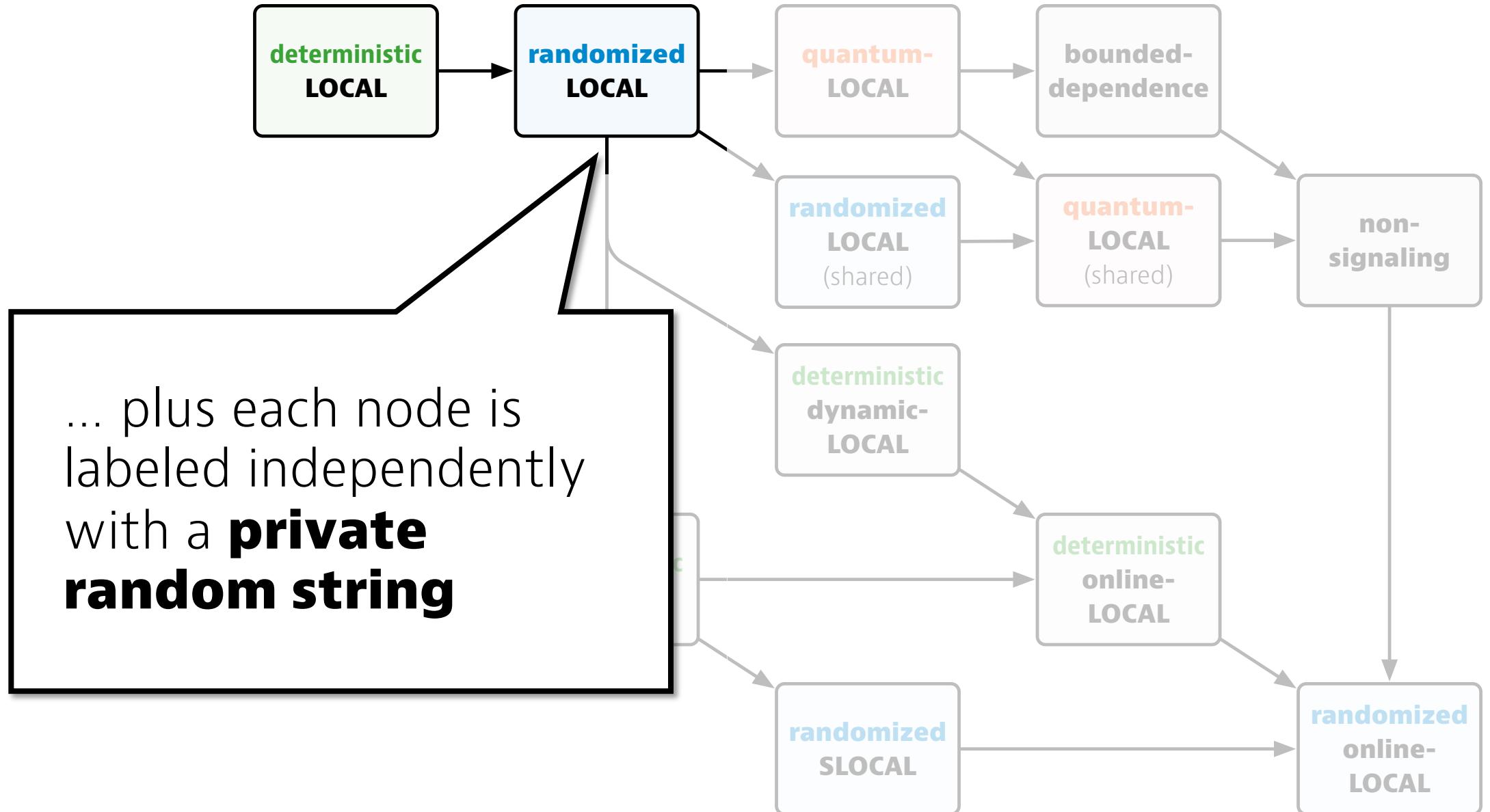


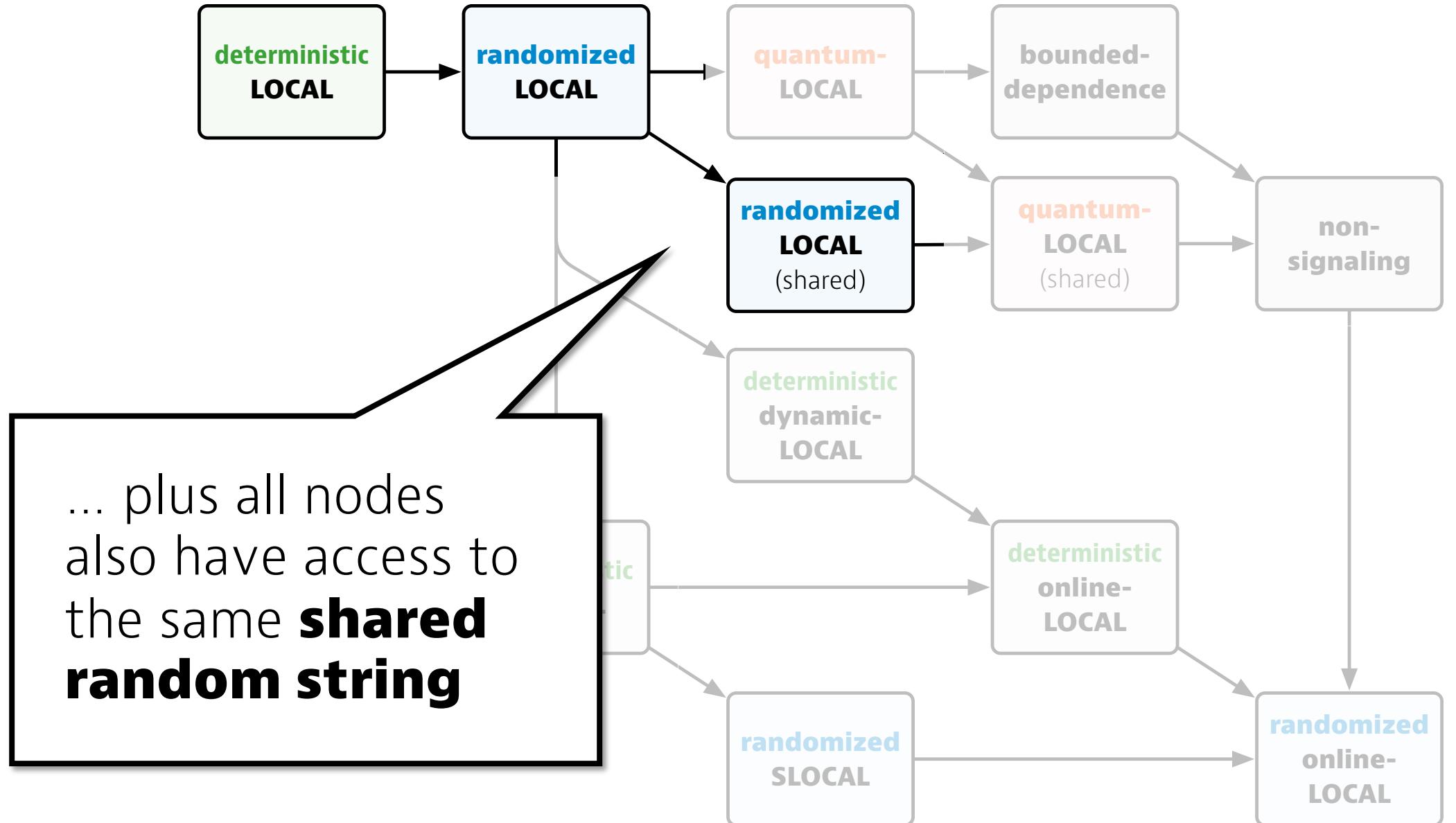


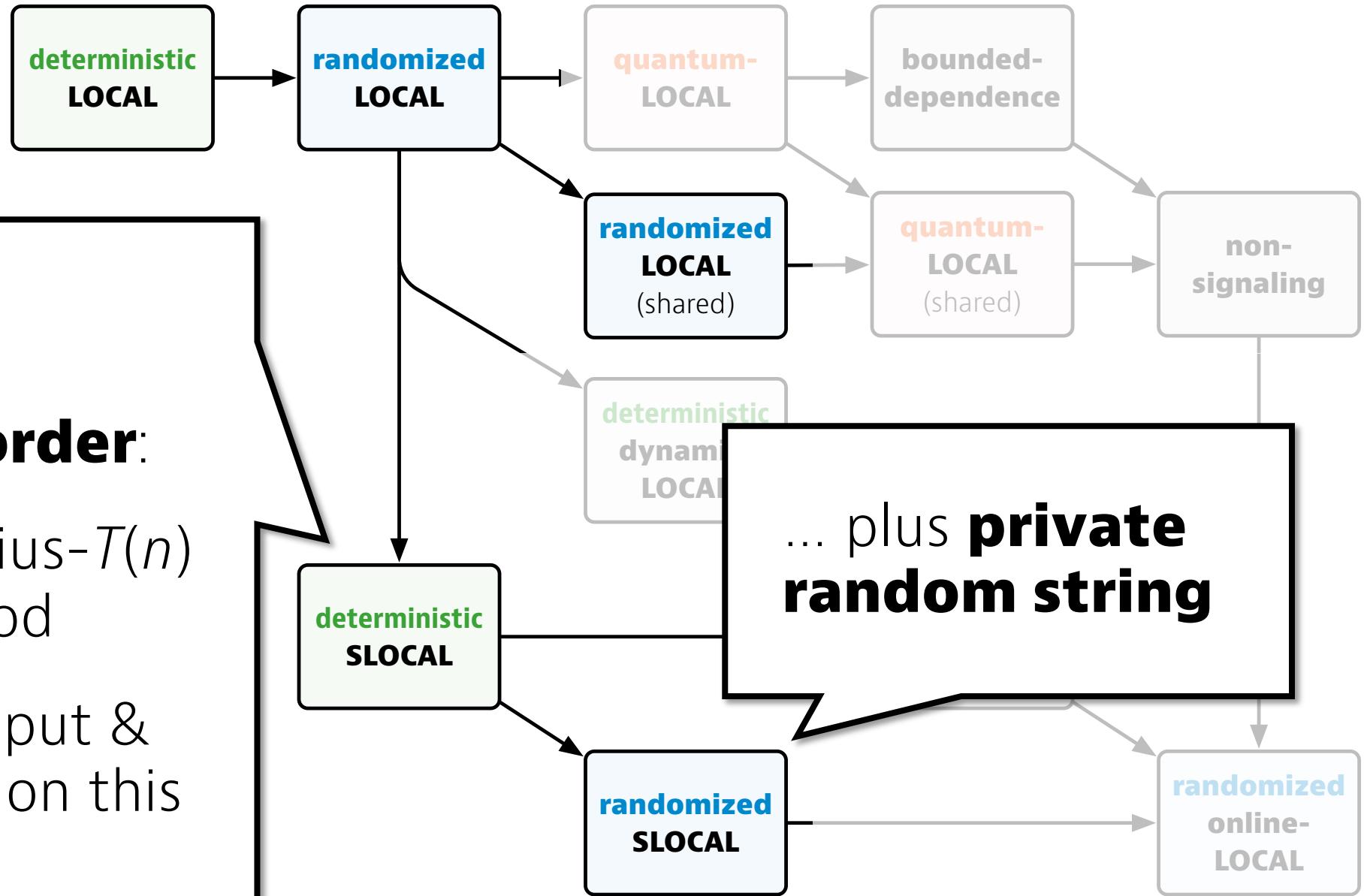
**So what are
all these models?**

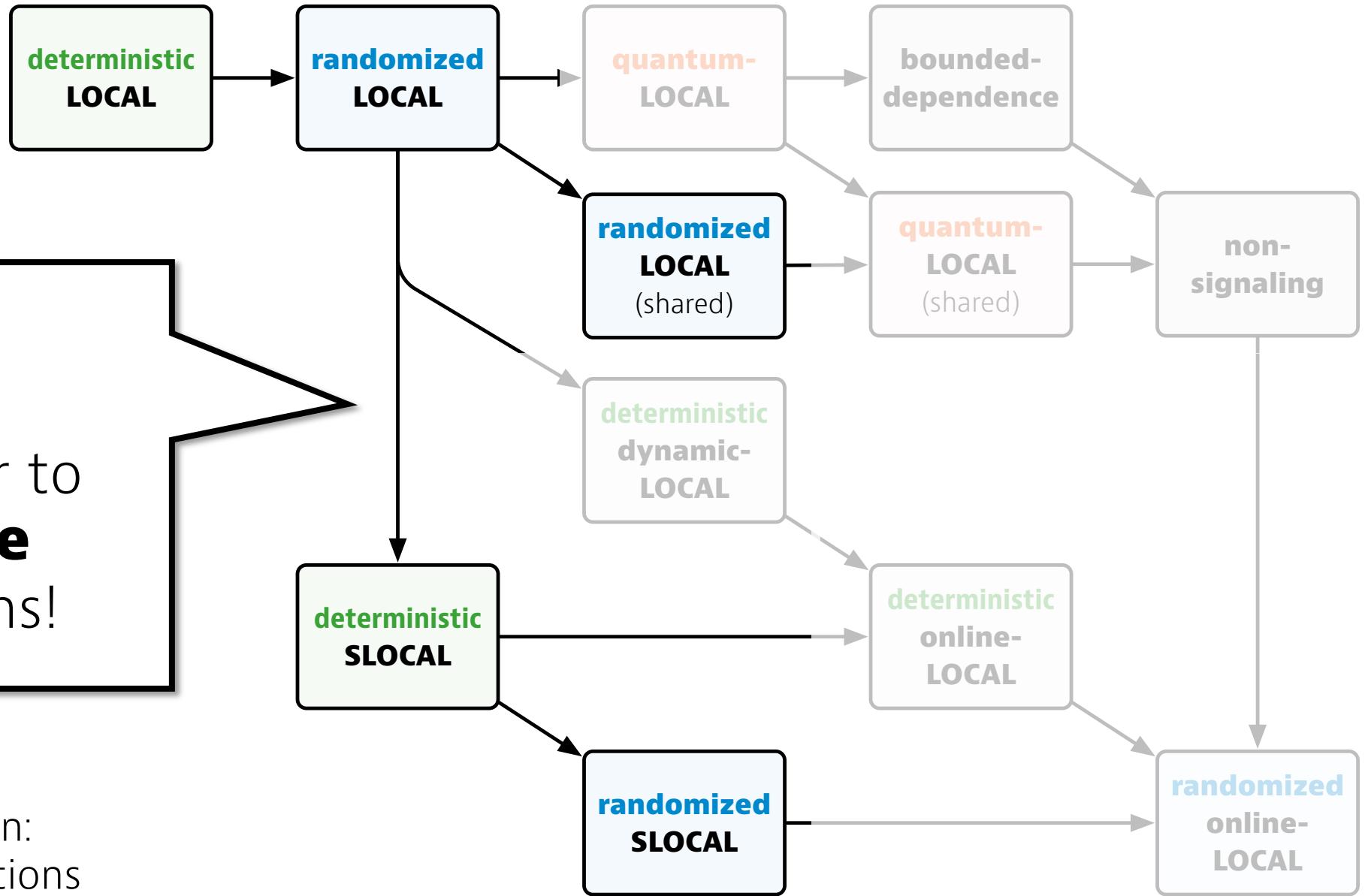










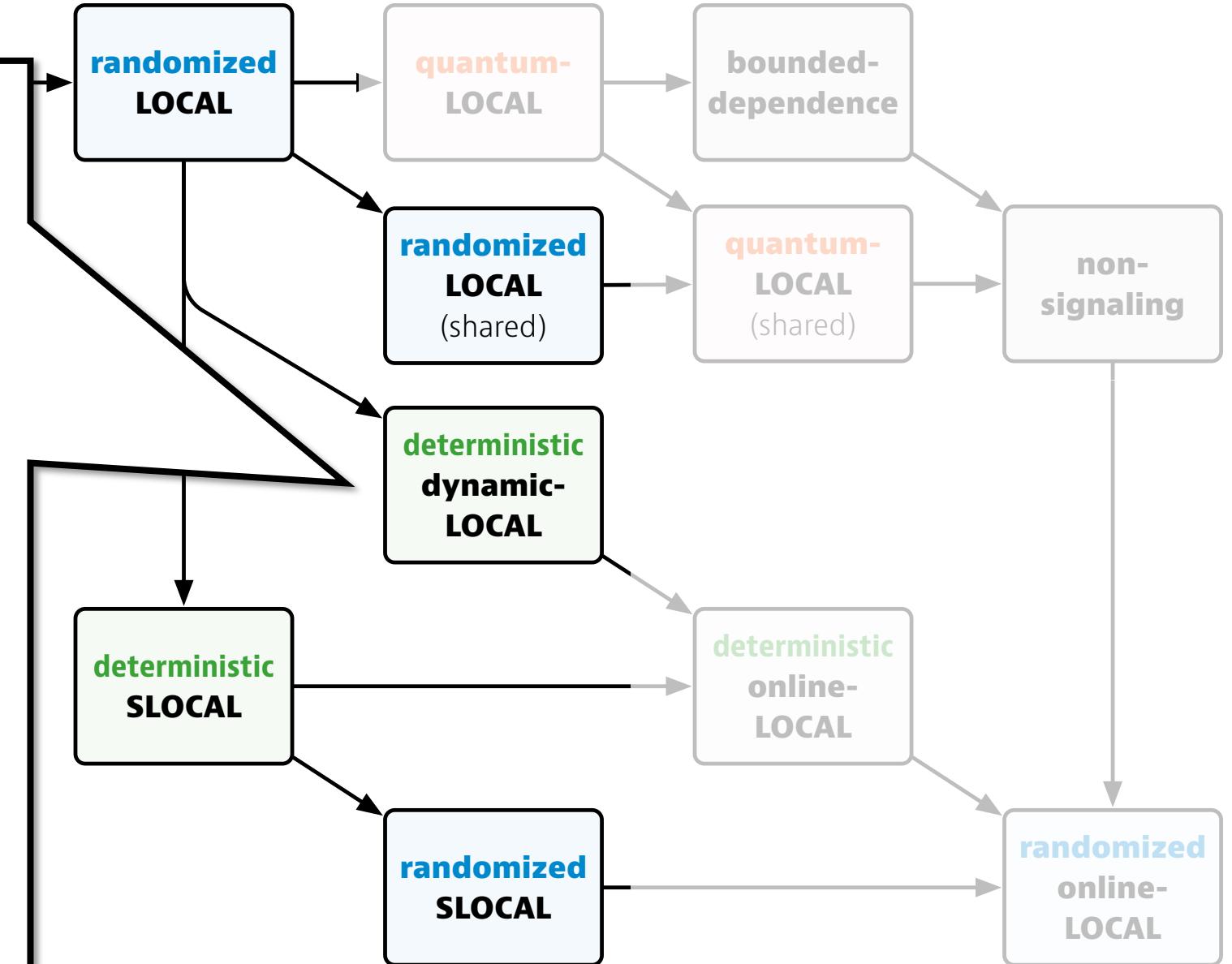


Ghaffari, Harris, Kuhn:
conditional expectations

Adversary adds nodes
and edges one by one

- We can **see everything**
- Have to maintain valid solution

We can **change** our output only within distance $T(n)$ from a point of change



Limited
what we
can **see**

deterministic
LOCAL

randomized
LOCAL

quantum-
LOCAL

bounded-
dependence

randomized
LOCAL
(shared)

quantum-
LOCAL
(shared)

non-
signaling

Limited
what we
can **change**

deterministic
dynamic-
LOCAL

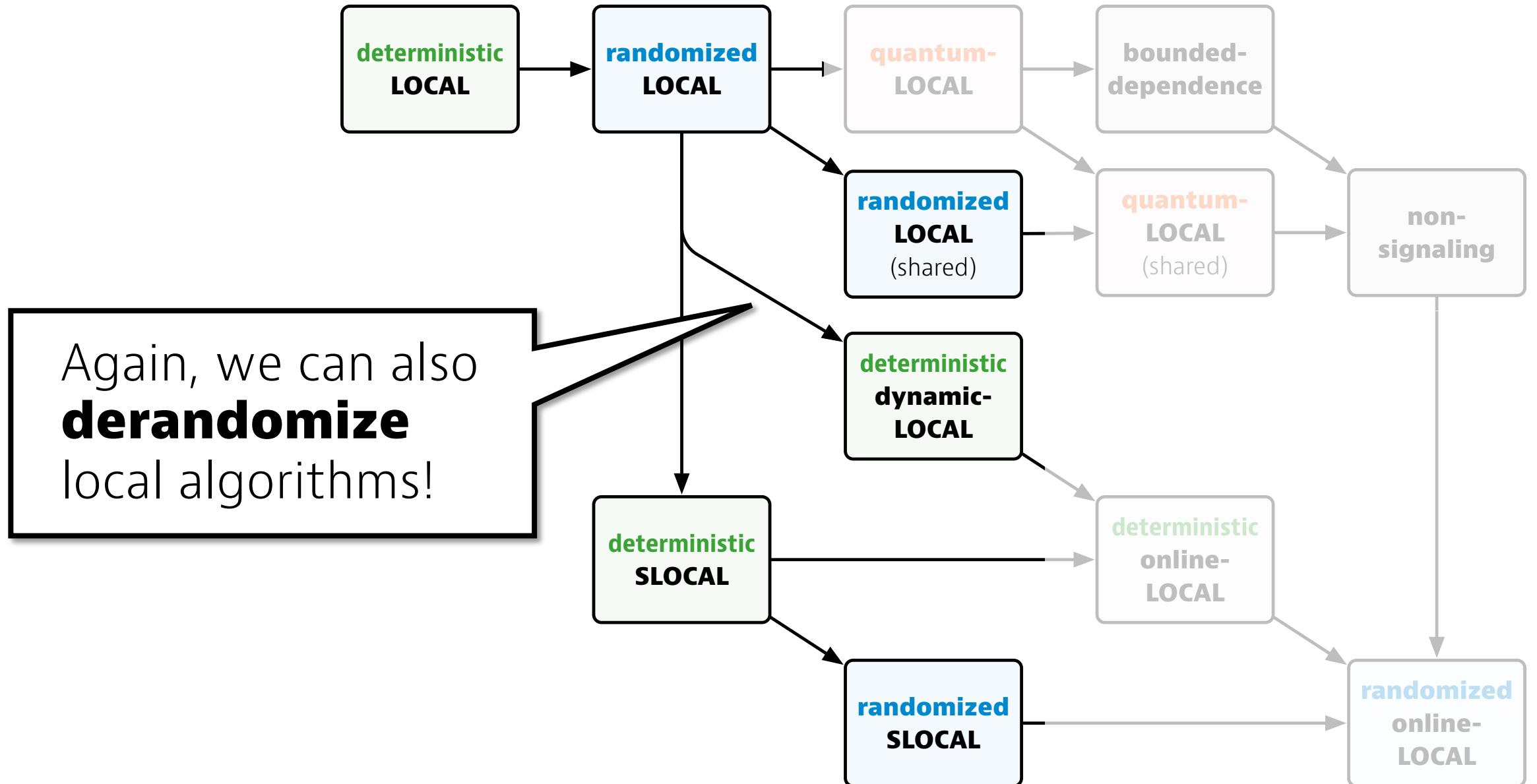
deterministic
SLOCAL

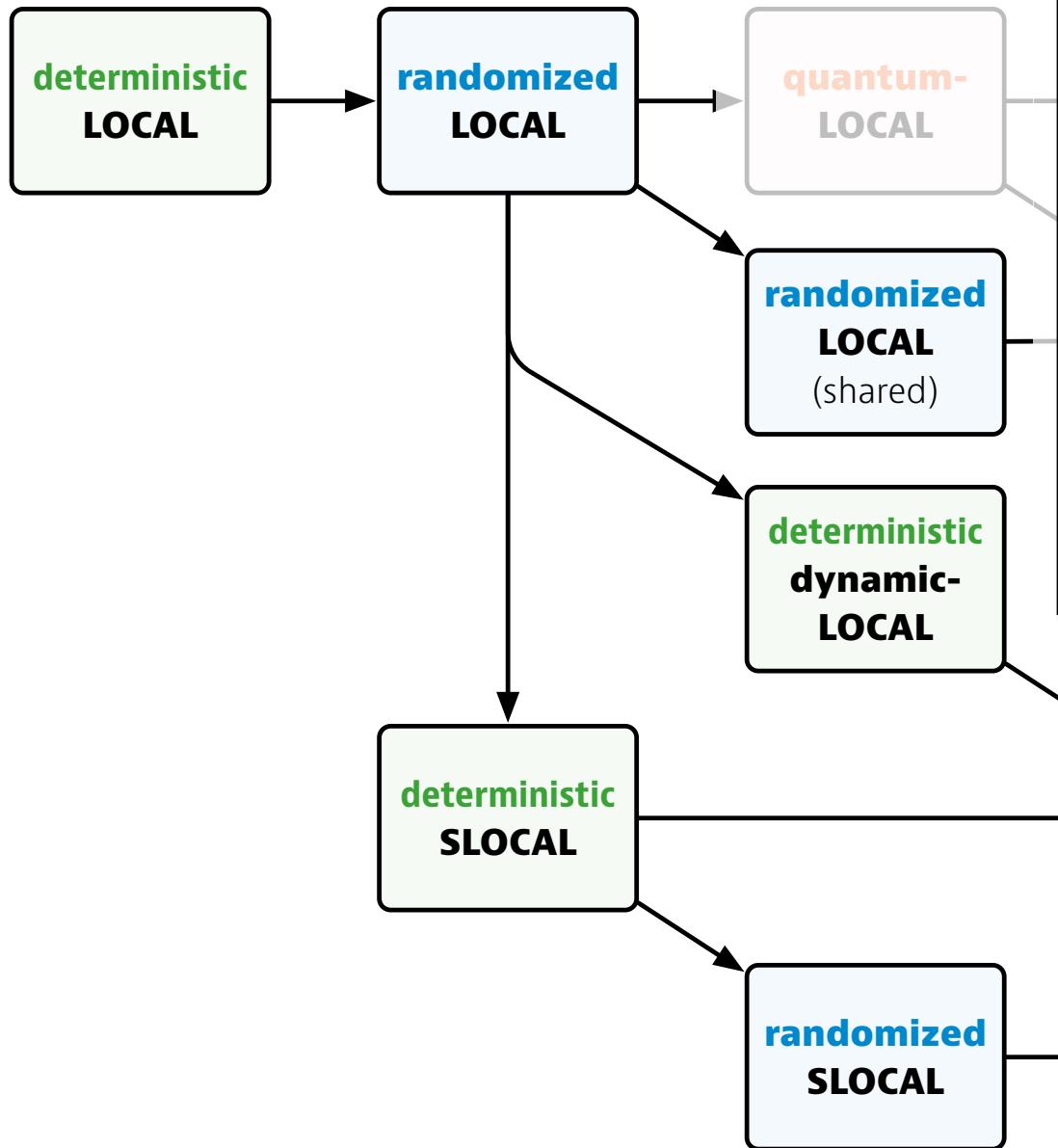
deterministic
online-
LOCAL

randomized
SLOCAL

randomized
online-
LOCAL

Why do we
have **LOCAL** →
dynamic-LOCAL?

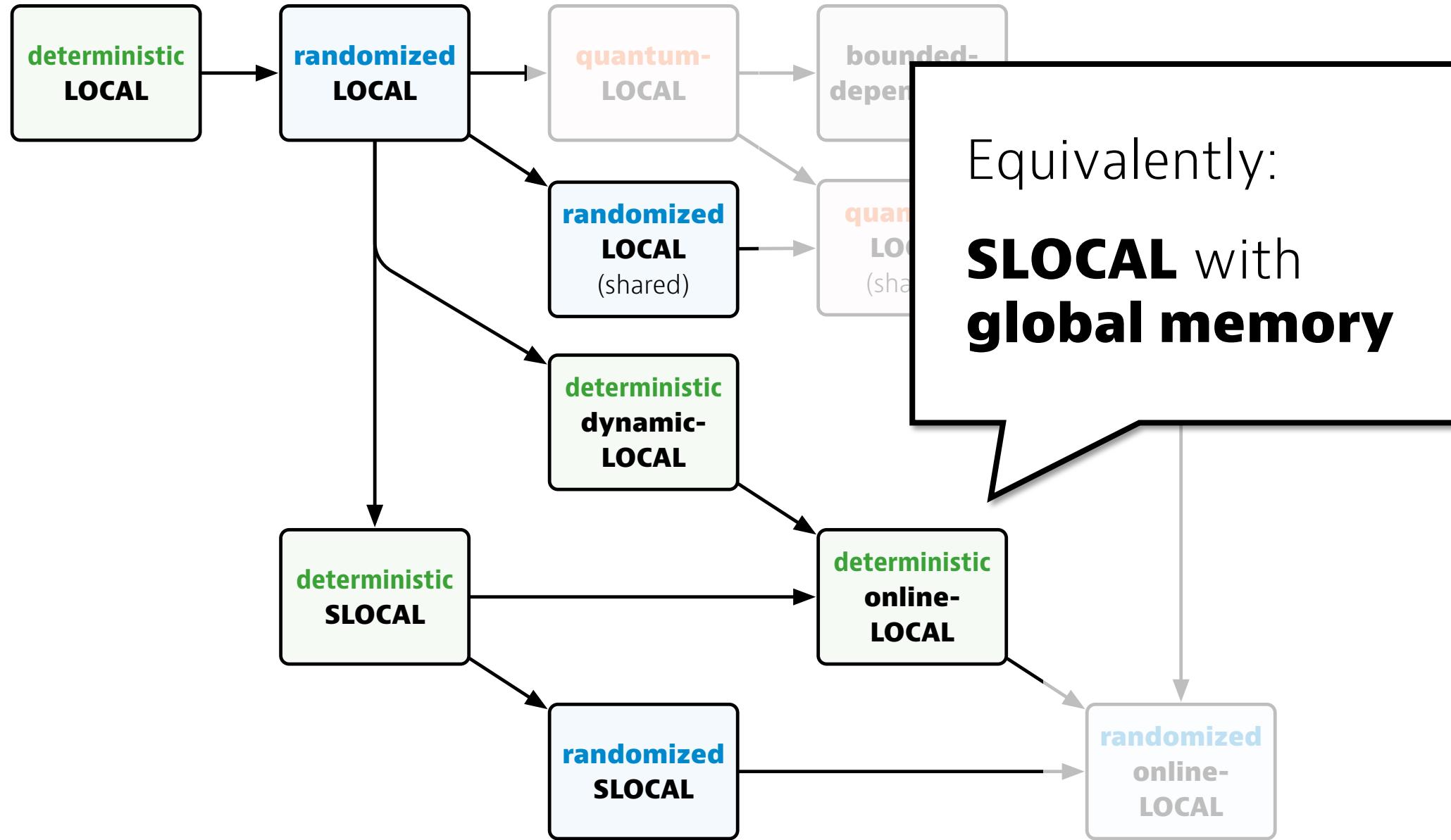


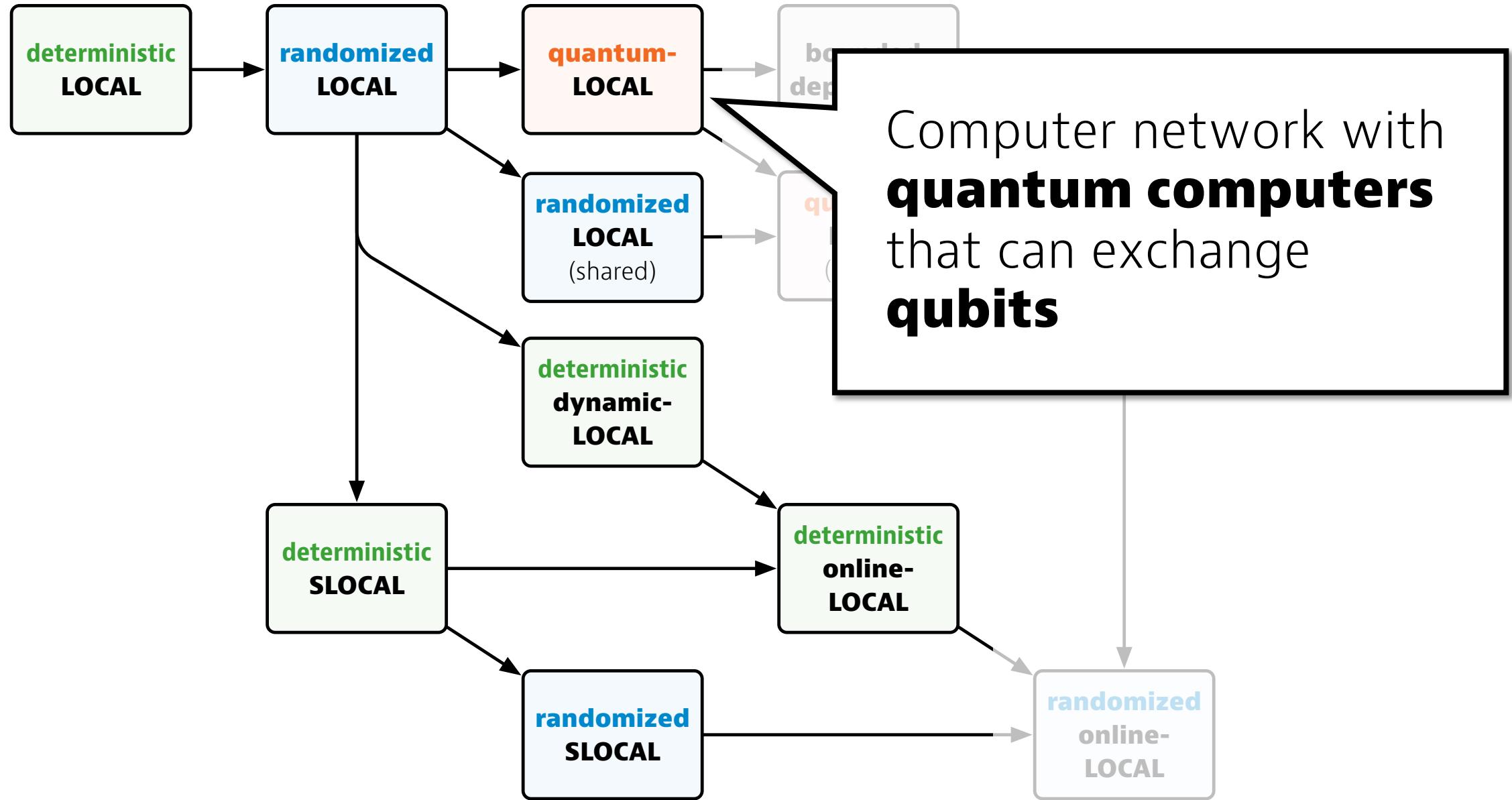


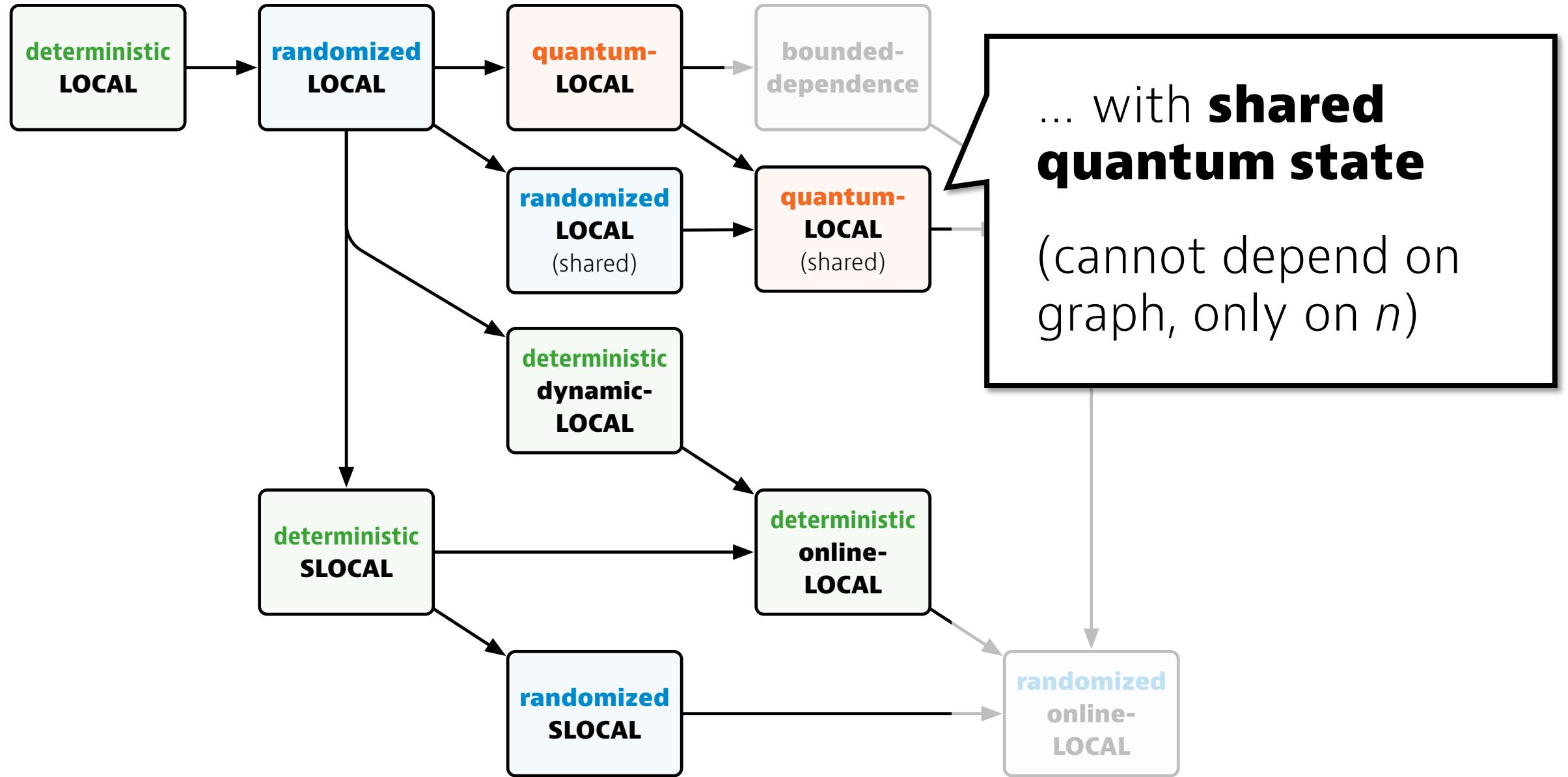
Online graph algorithms

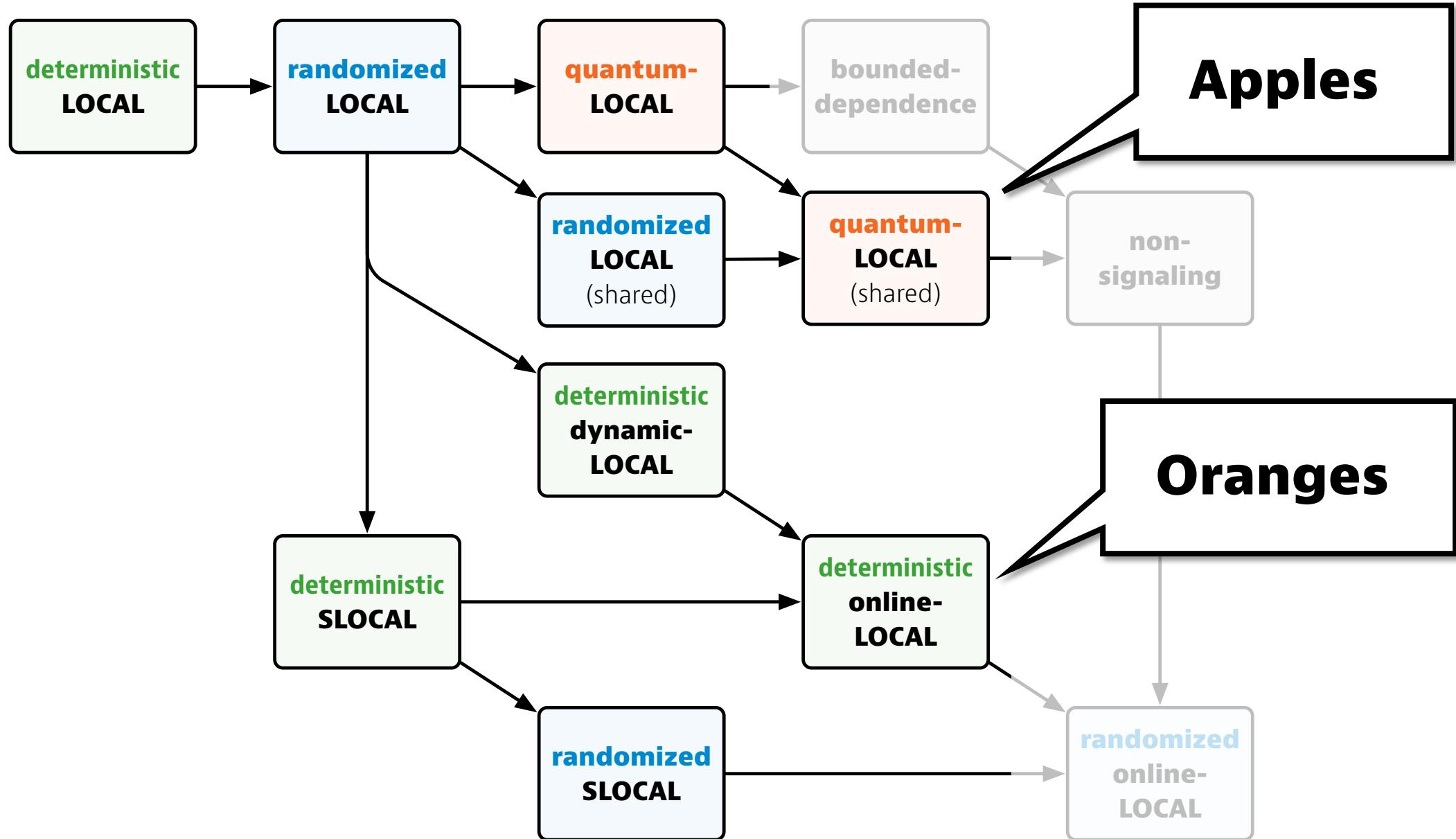
Graph revealed by adversary

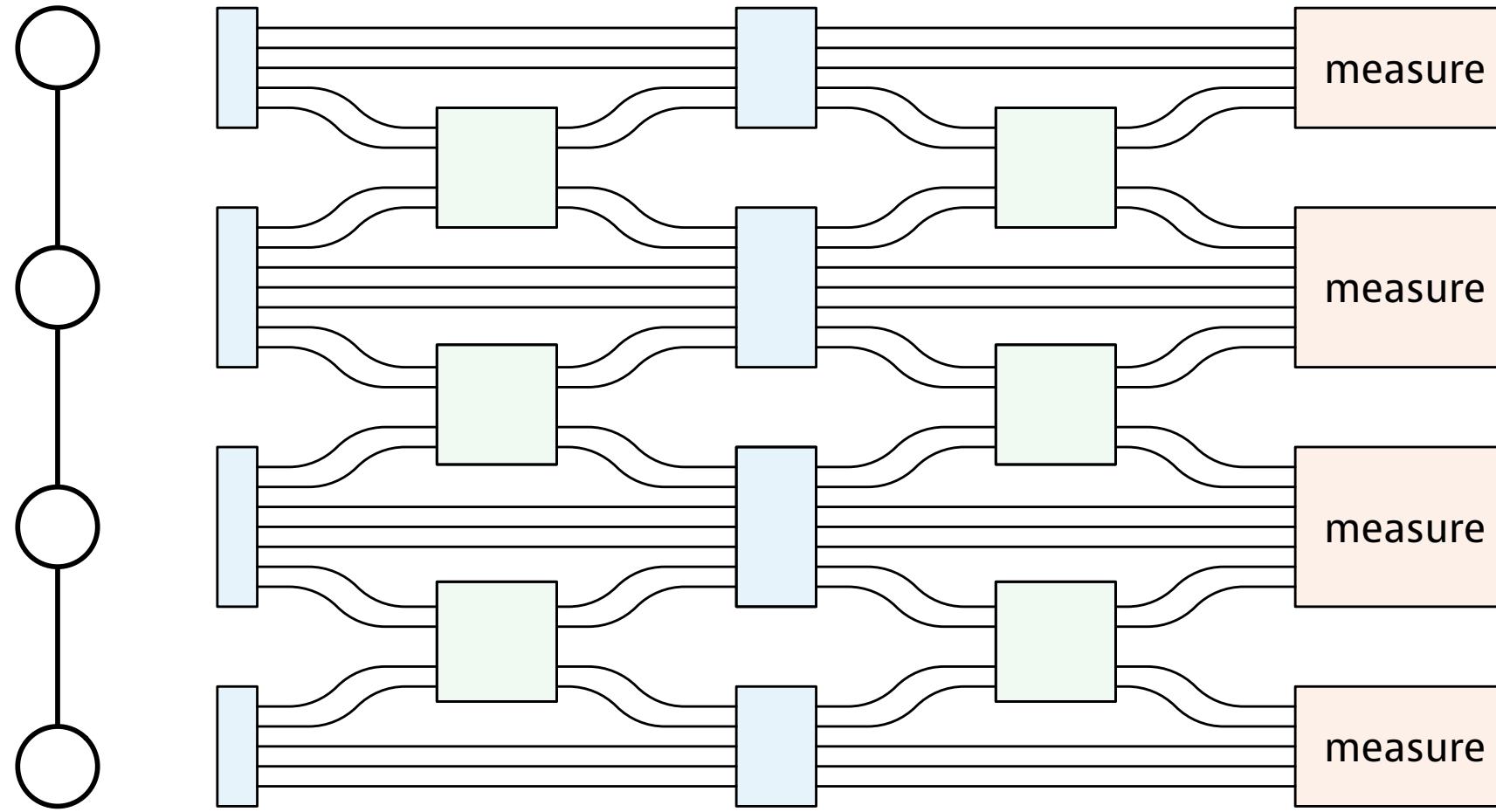
For each node we can **delay decisions** until we see its radius- $T(n)$ neighborhood









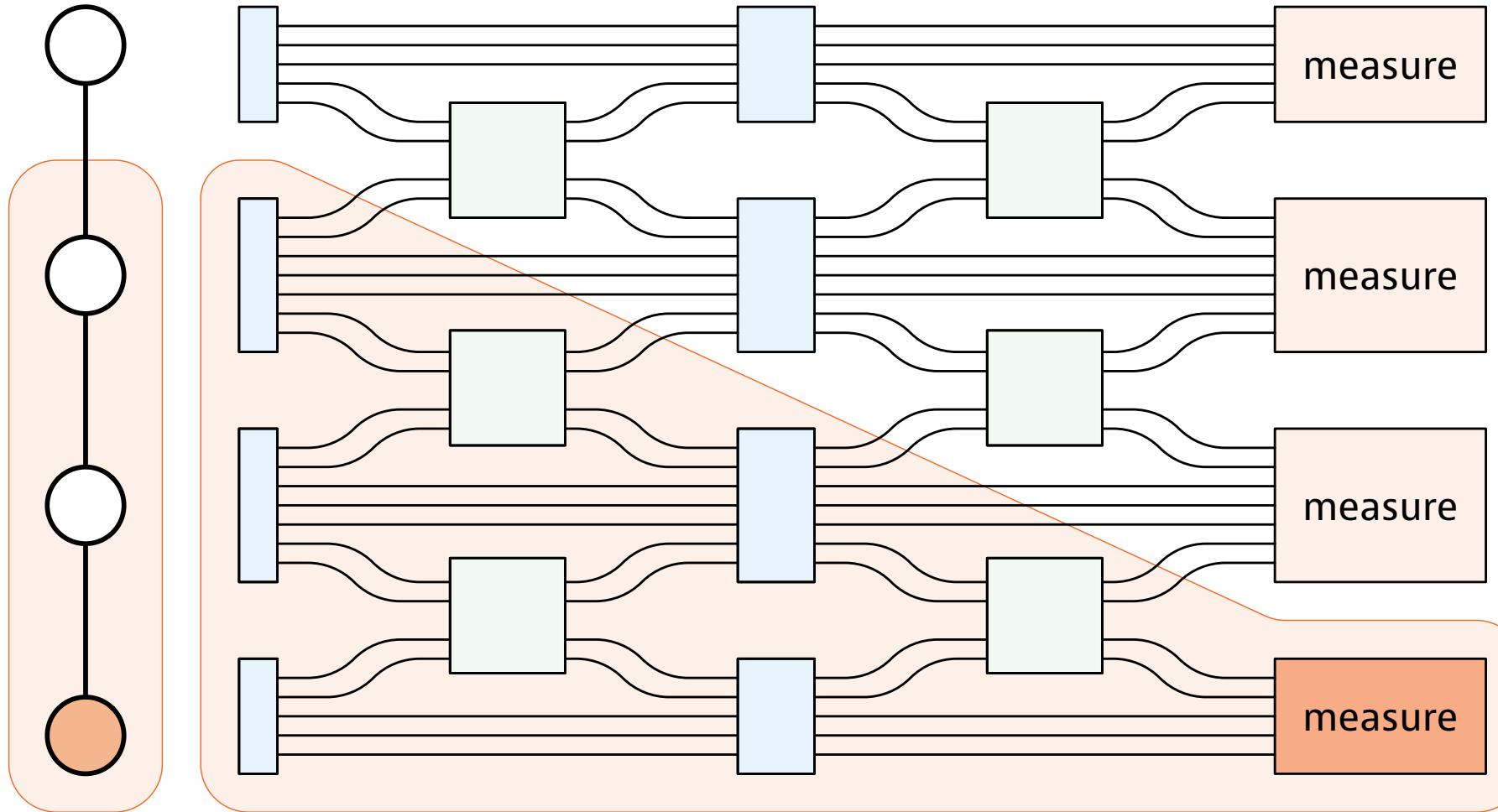


2 rounds

communication

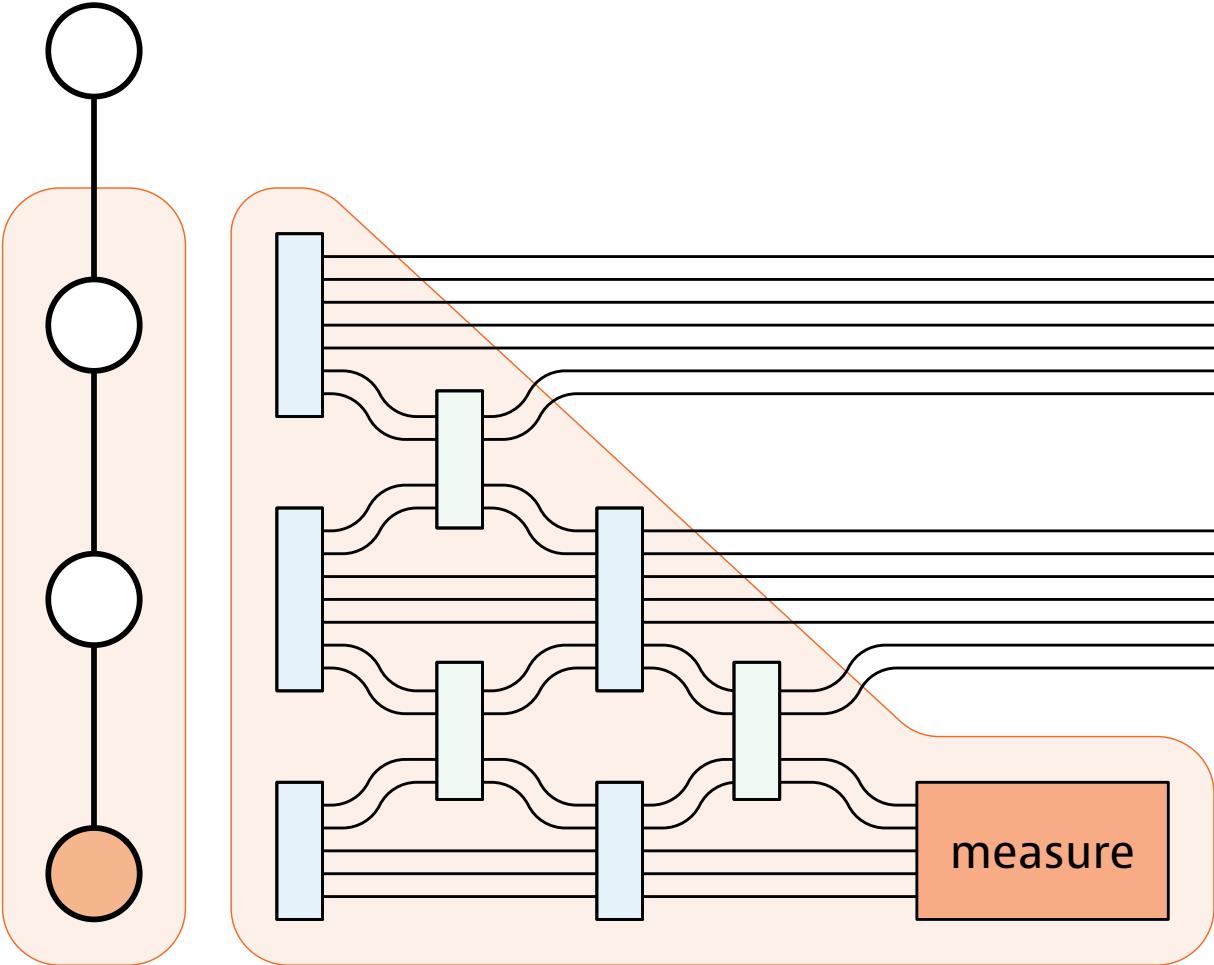
local
computation

communication



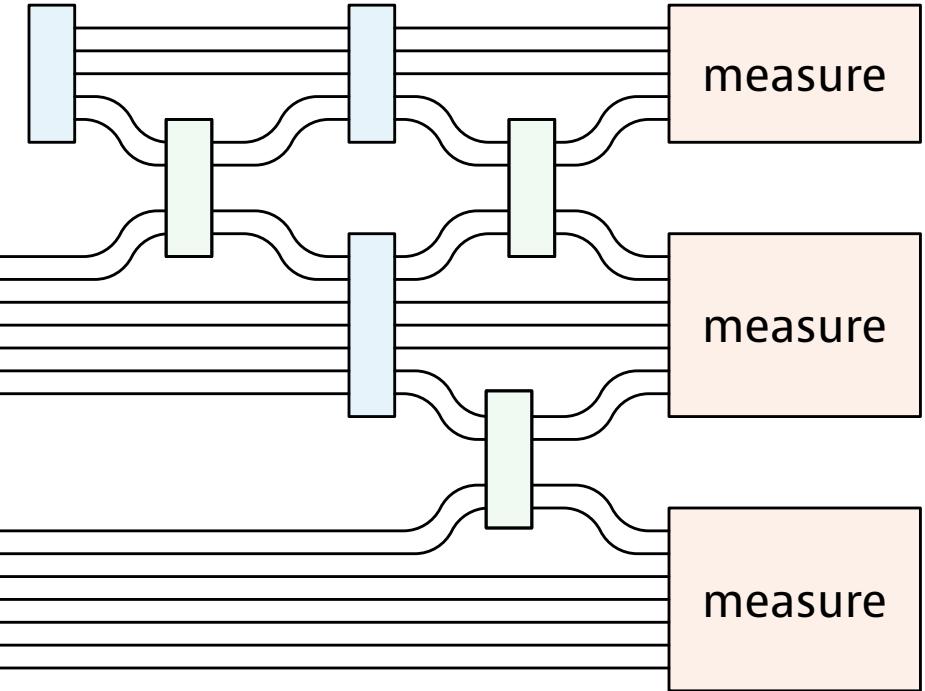
2 rounds

light cone



2 rounds

light cone



Non-signaling model

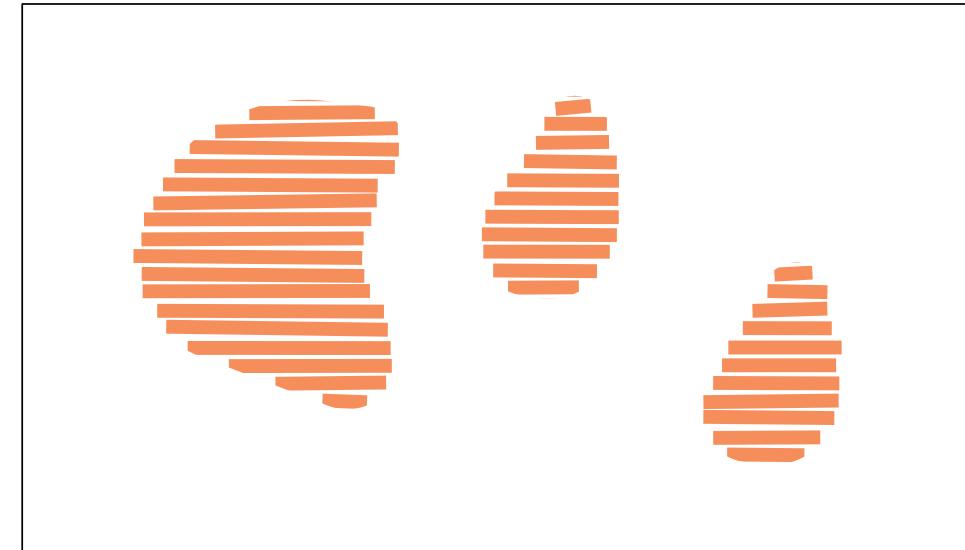
- Quantum LOCAL can't violate causality
- Key idea: **define** a model so that it can do **anything** except violating causality

Non-signaling model

Definition (*non-signaling distribution*):

- fix any **set of nodes X** ...

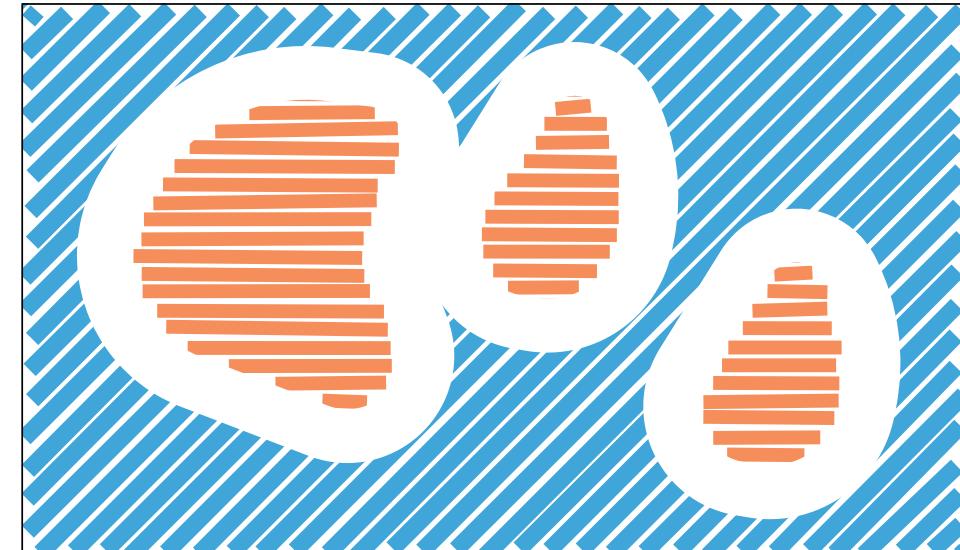
Gavoille, Kosowski, Markiewicz 2009
Arfaoui, Fraigniaud 2014



Non-signaling model

Definition (*non-signaling distribution*):

- fix any **set of nodes X**
- changes in the input **more than T hops away** from X do not influence the output distribution of X



Gavoille, Kosowski, Markiewicz 2009
Arfaoui, Fraigniaud 2014

Classical
probability
theory

Classical (randomized) distributed algorithms



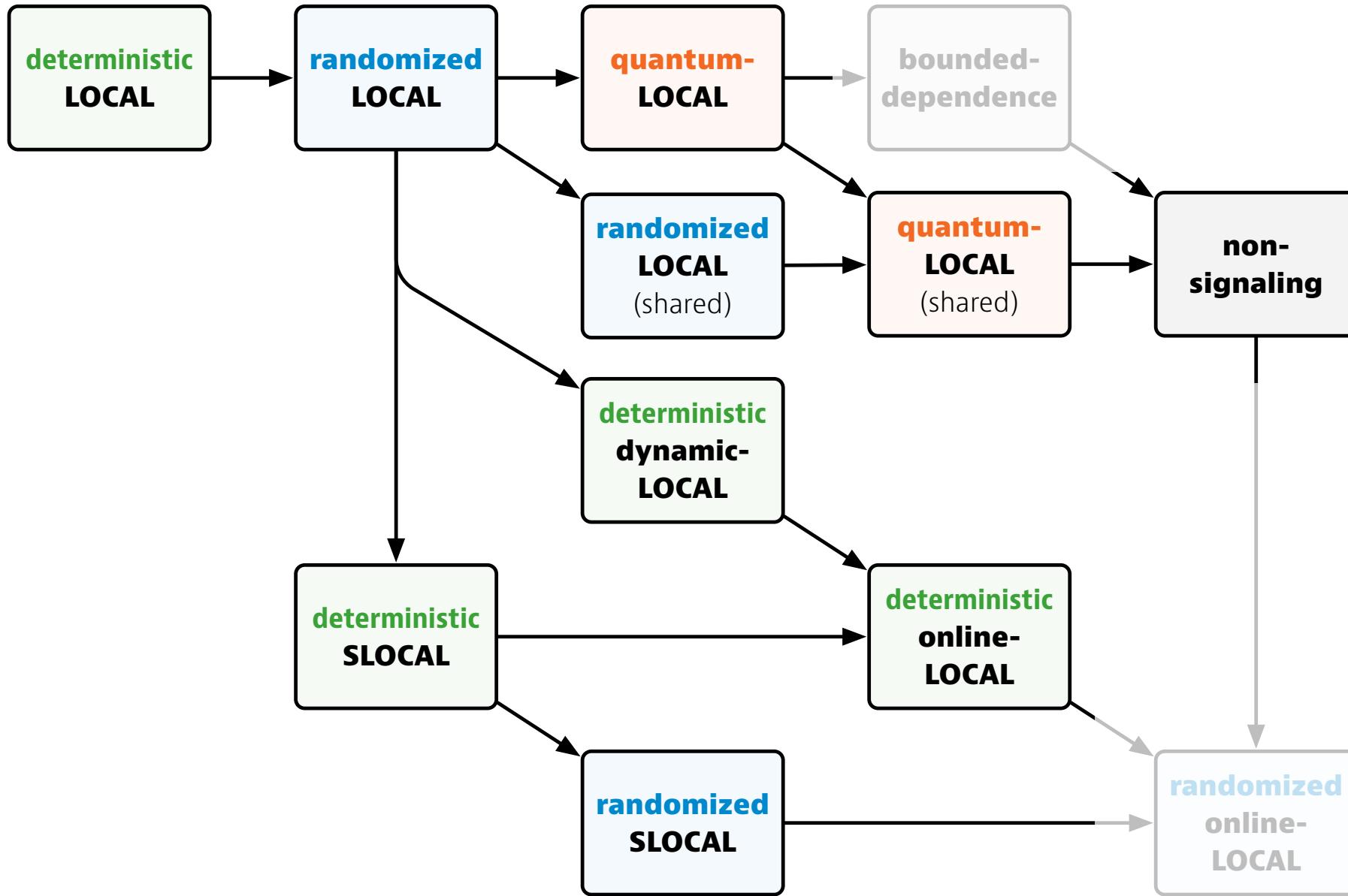
Quantum distributed algorithms

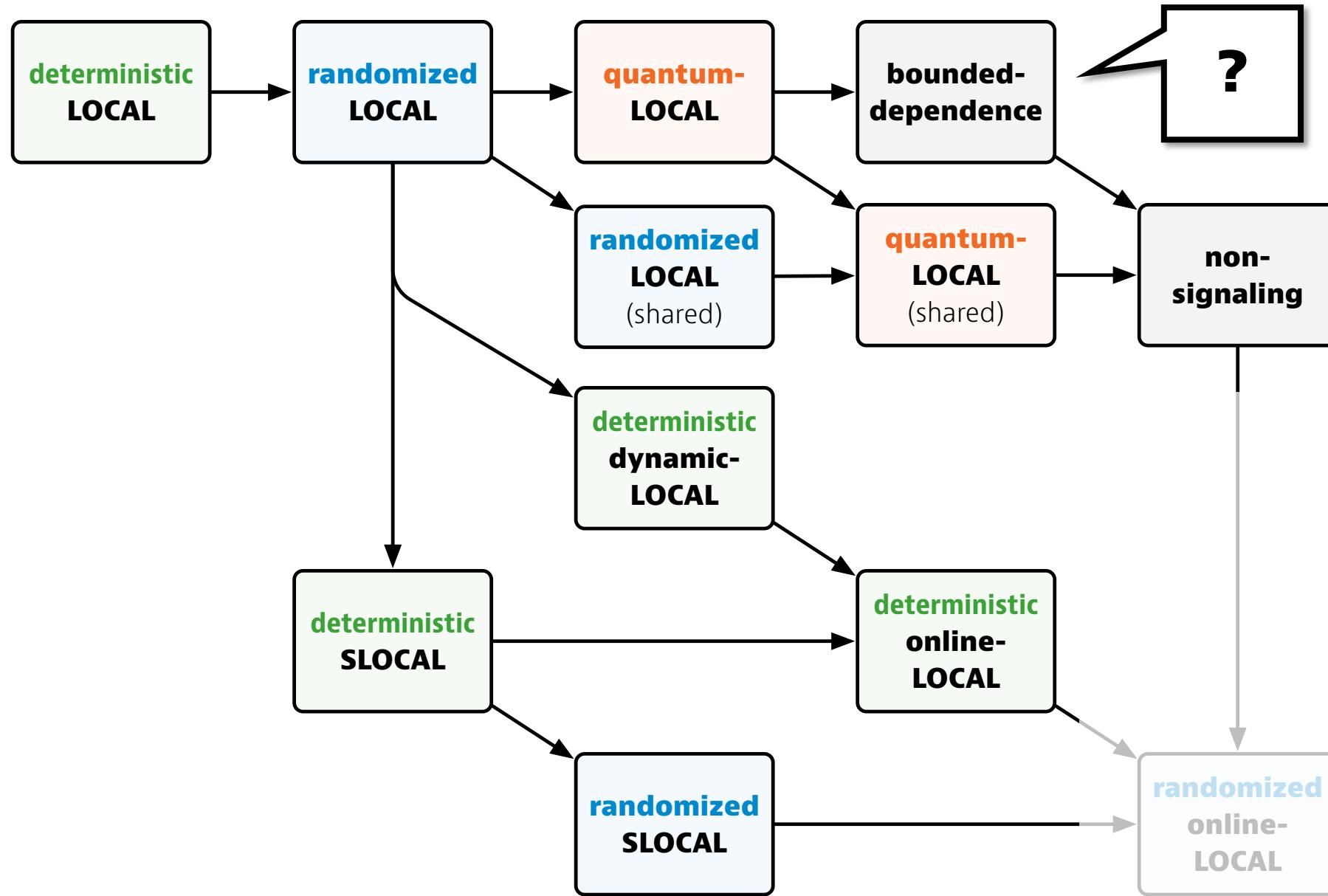


Non-signaling “algorithms”

Weird
quantum
things

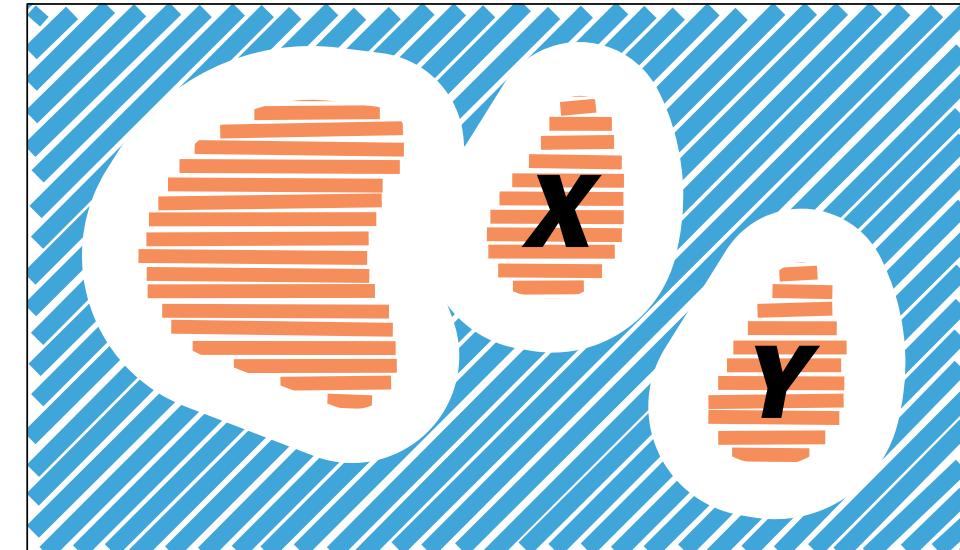
Classical
probability
theory





Bounded dependence

- **Finitely dependent distributions**
 - X and Y far from each others \rightarrow independent
 - usually “far” = some constant
- For clarity, we call it here **bounded dependence model** when “far” = $T(n)$



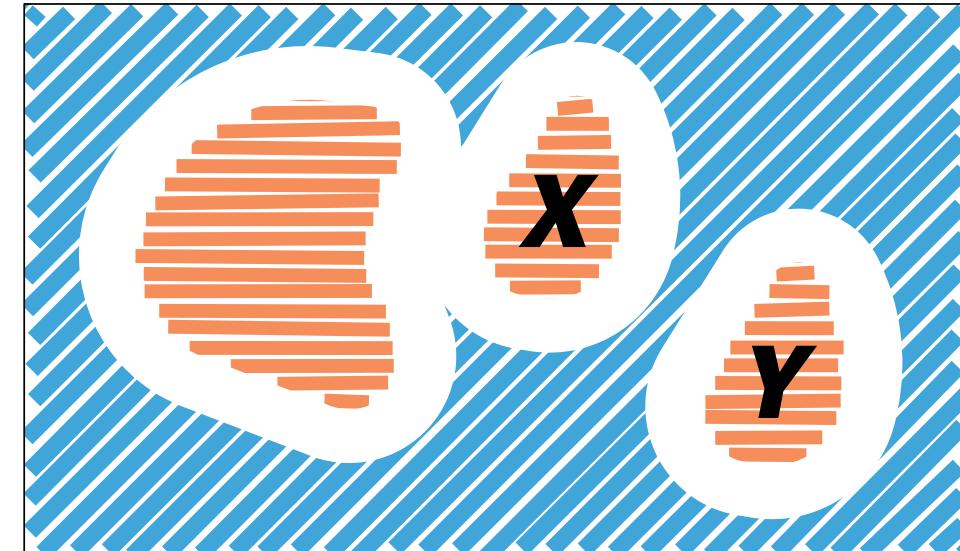
Bounded dependence

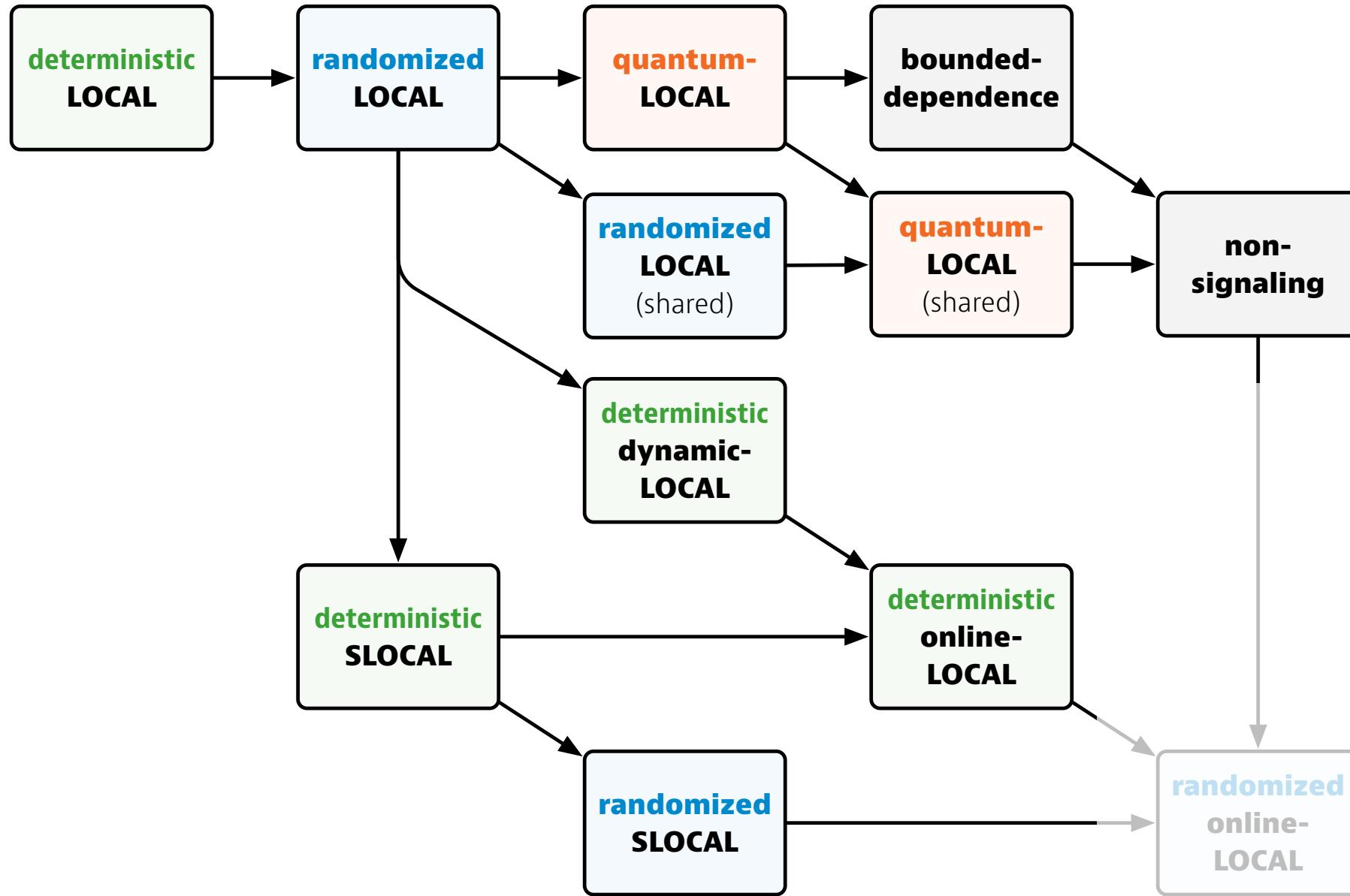
**Quantum-LOCAL without
shared quantum state:**

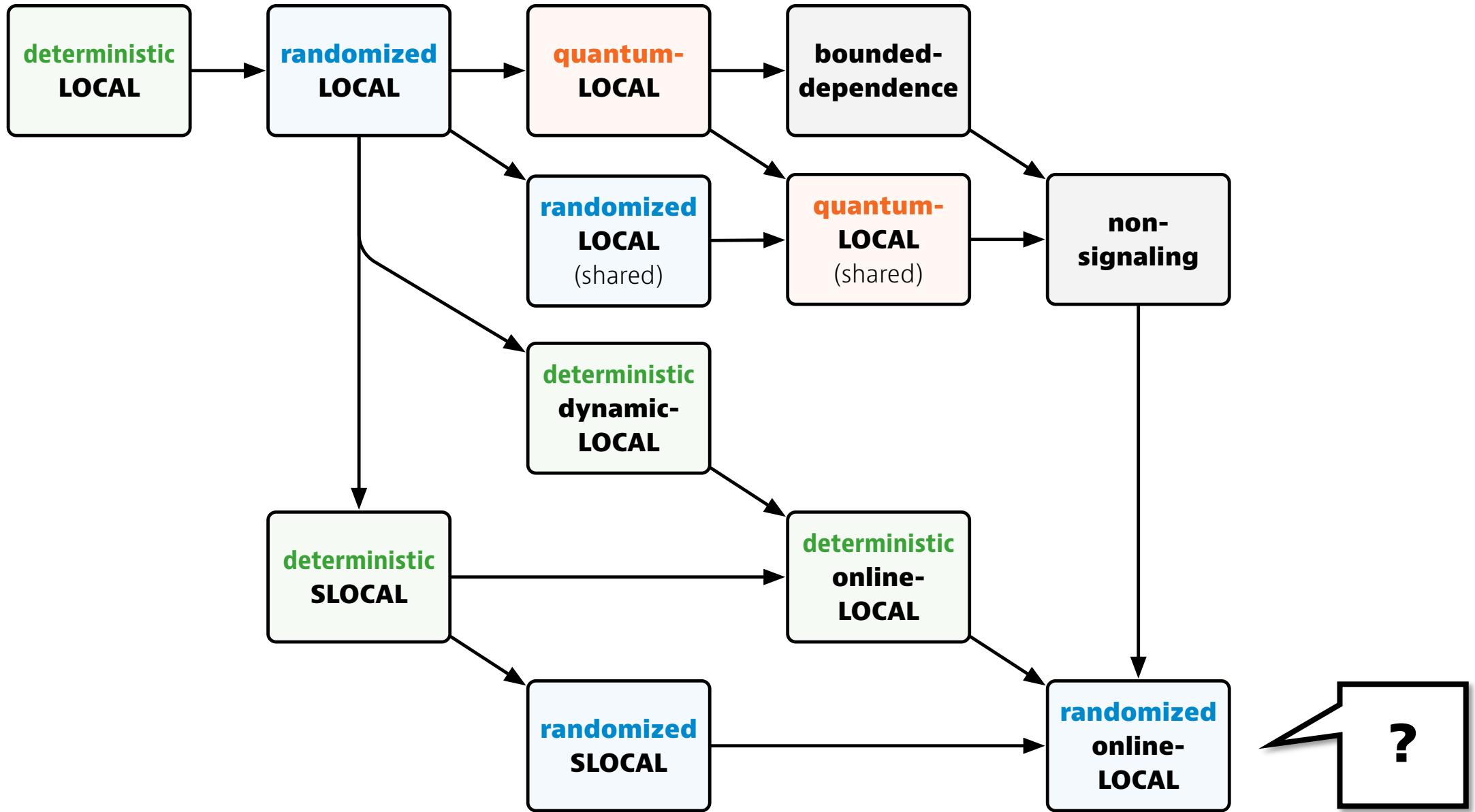
Output of $T(n)$ -round algorithm



Not just non-signaling but
also bounded dependence
for distance $\approx T(n)$







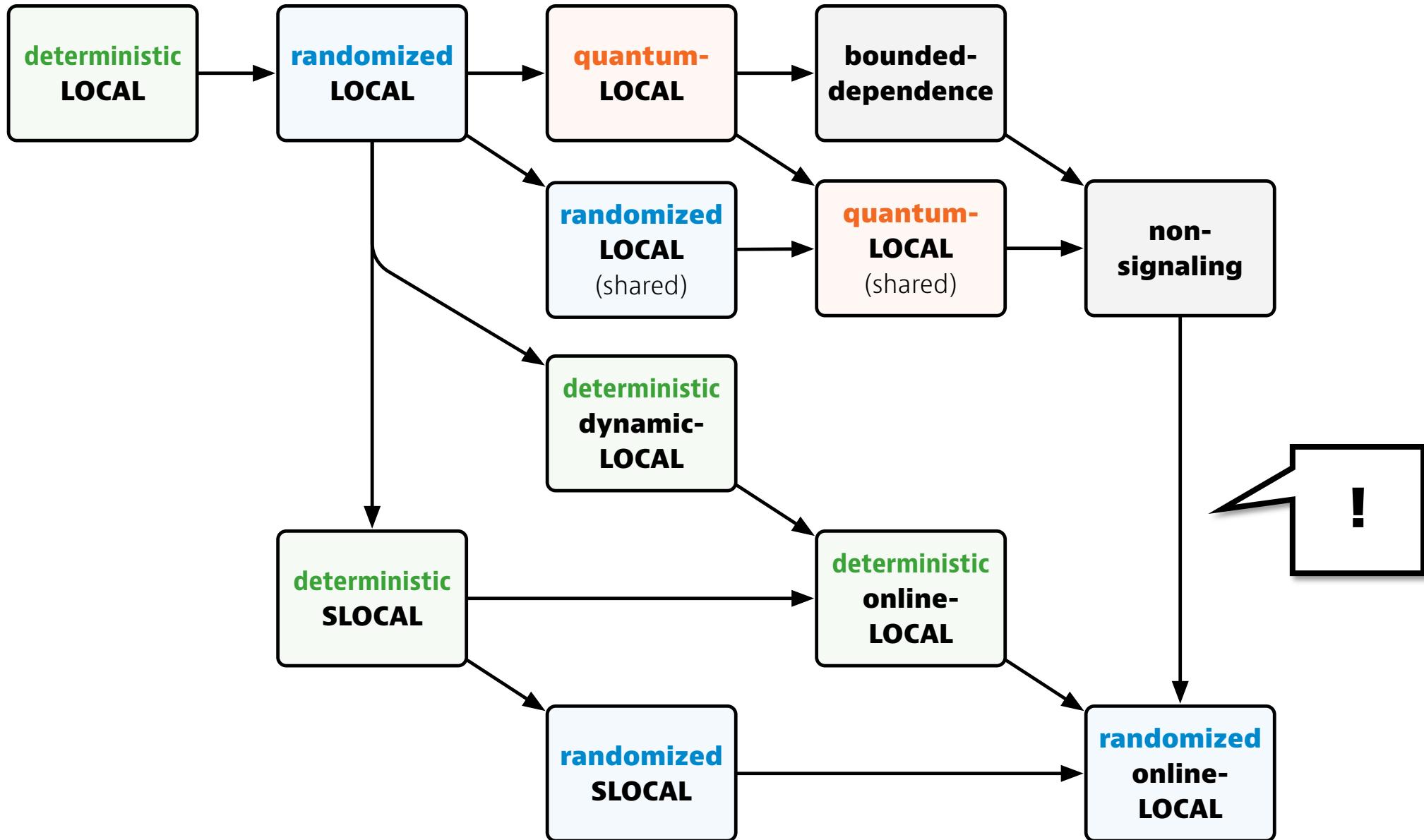
Rand. online-LOCAL

- Adversary fixes a **graph** + **order** in which nodes are revealed
- For each node v
 - algorithm sees radius- T neighborhood of v
 - algorithm must choose the label of v
- Algorithm can **remember** everything,
algorithm can use **randomness**

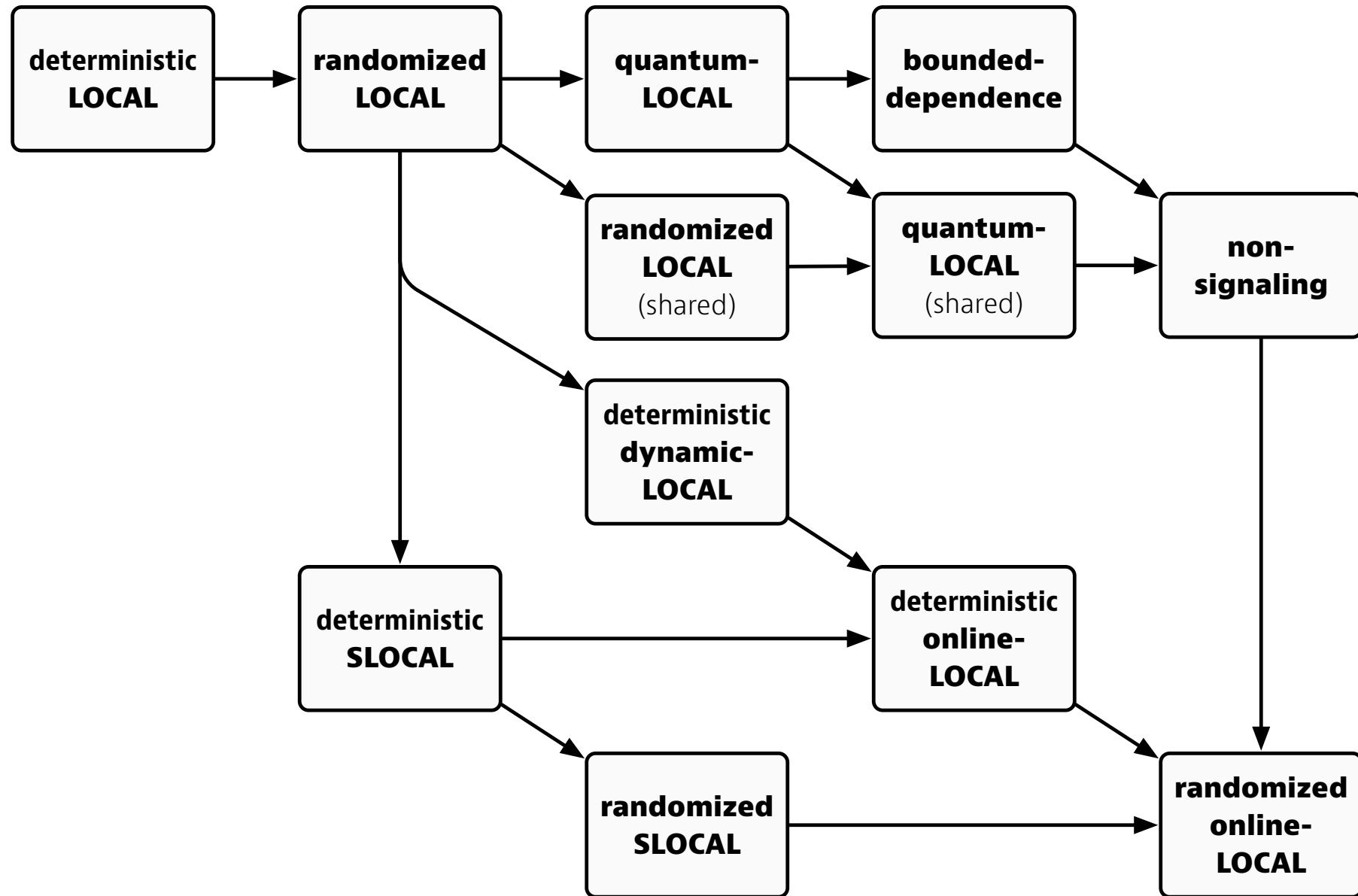
Oblivious adversary

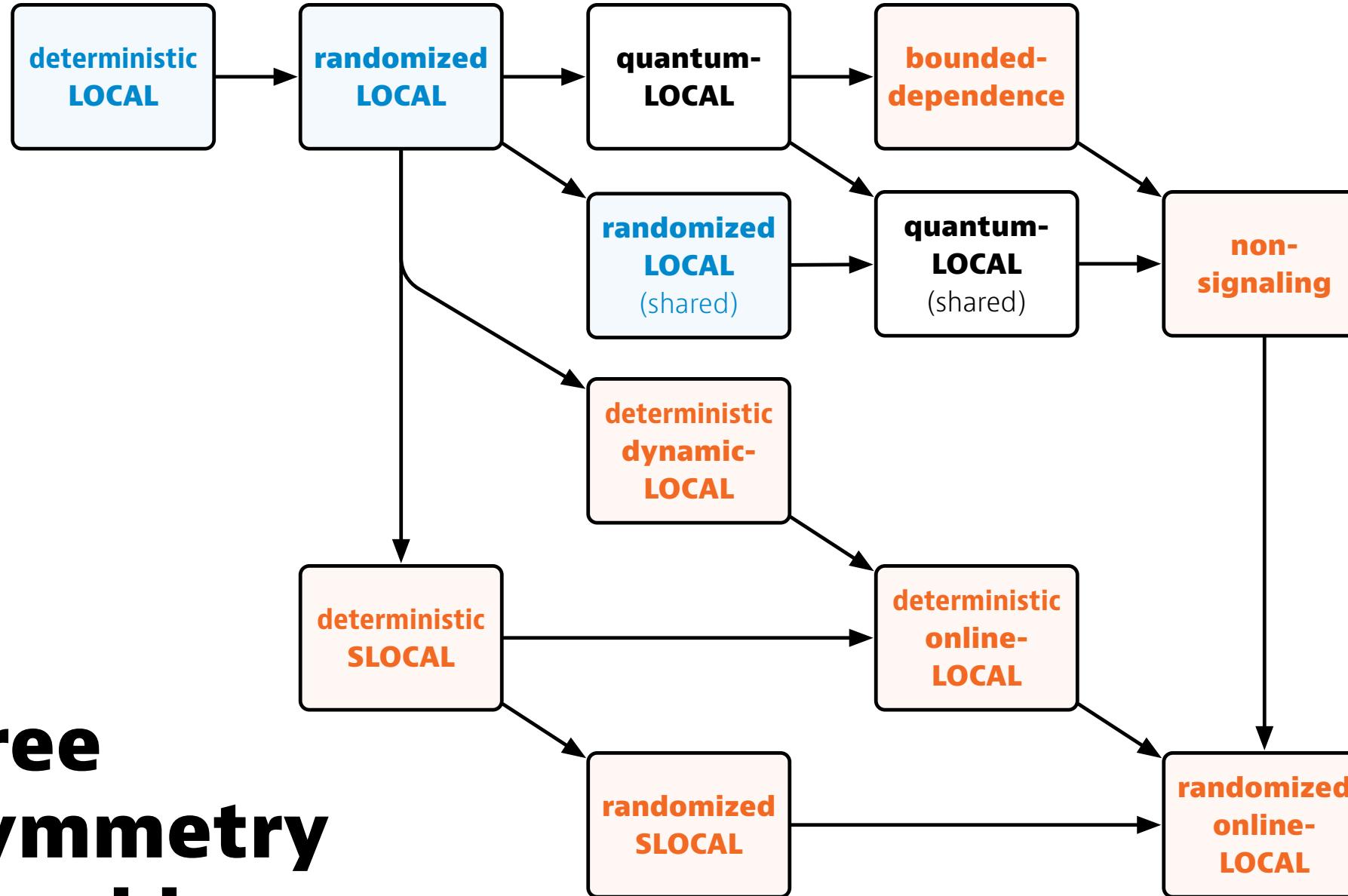
Rand. online-LOCAL

- **Trivial:**
 - randomize online-LOCAL can simulate deterministic online-LOCAL
- **Surprise:**
 - randomized online-LOCAL *can simulate any non-signaling distribution*
(with the same asymptotic locality)

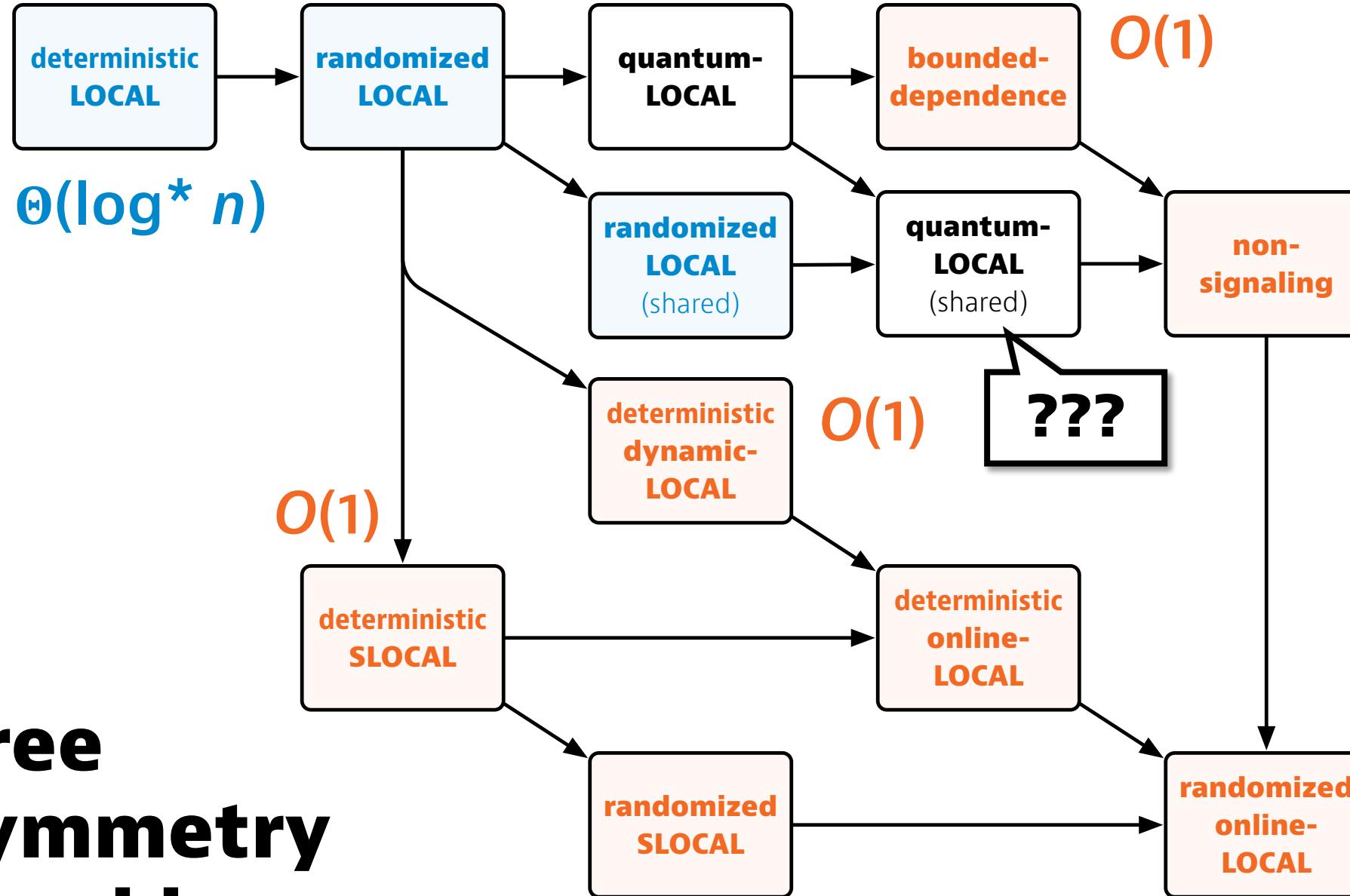


**Some
separations**

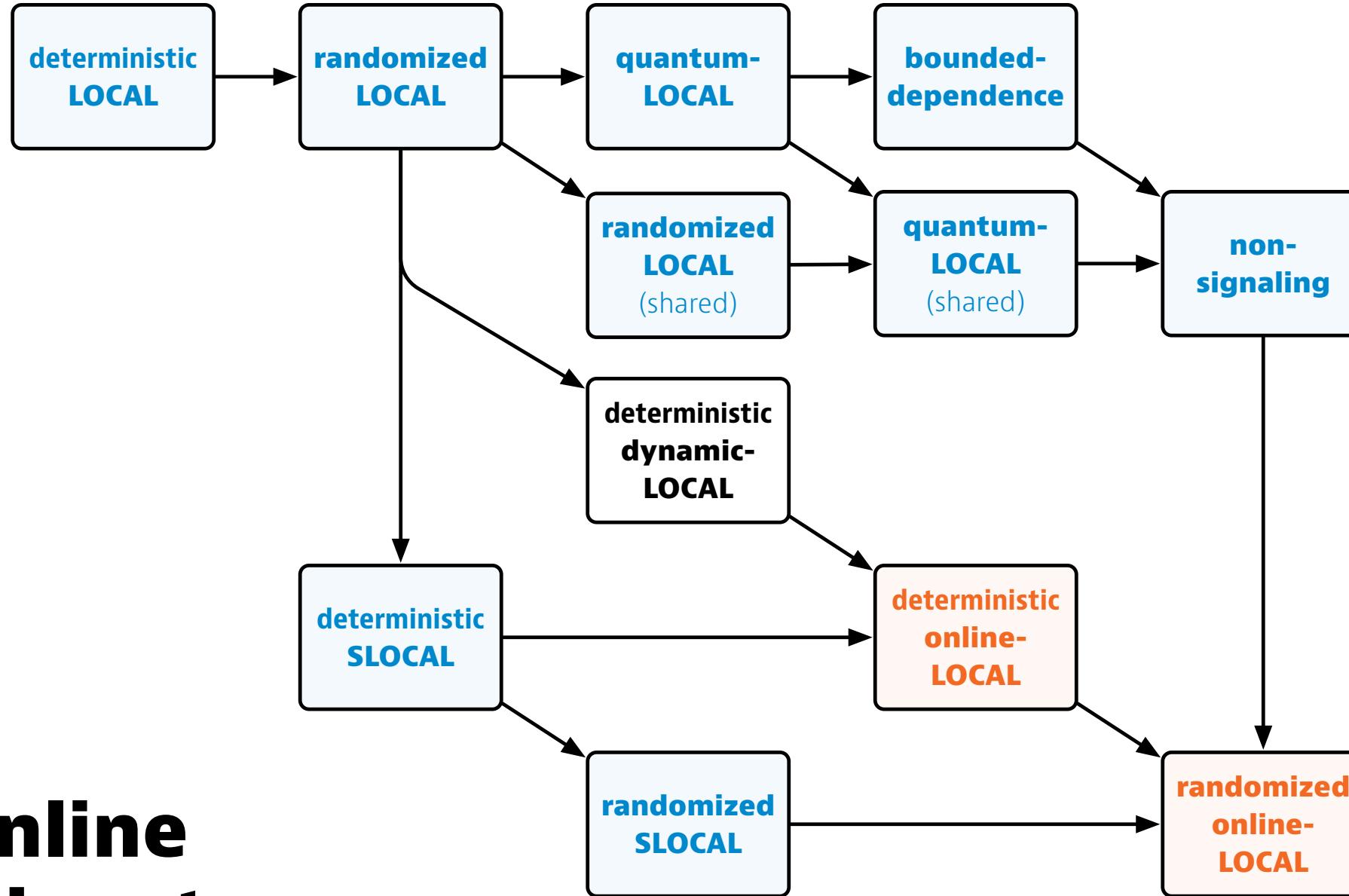




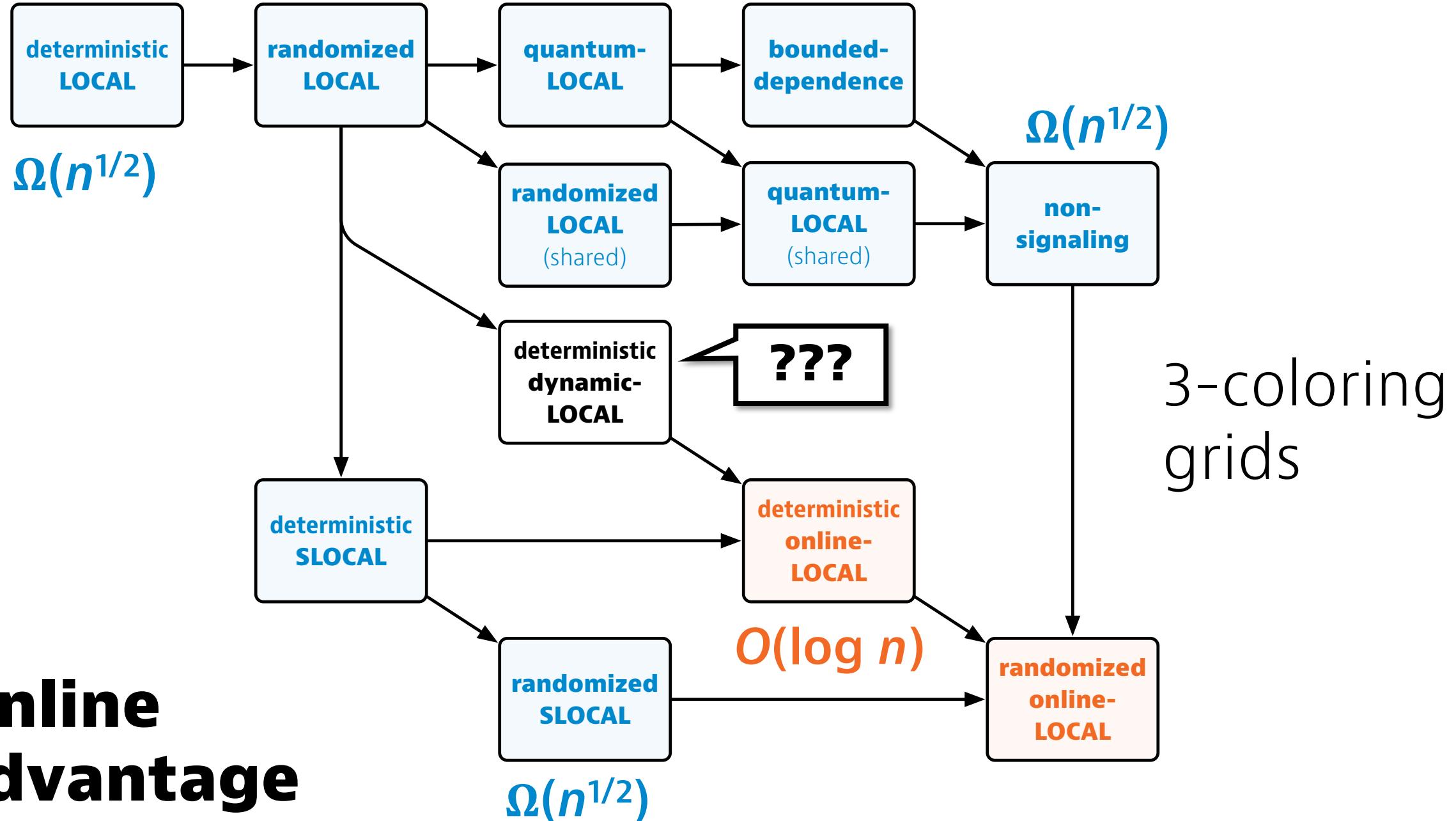
**Free
symmetry
breaking**

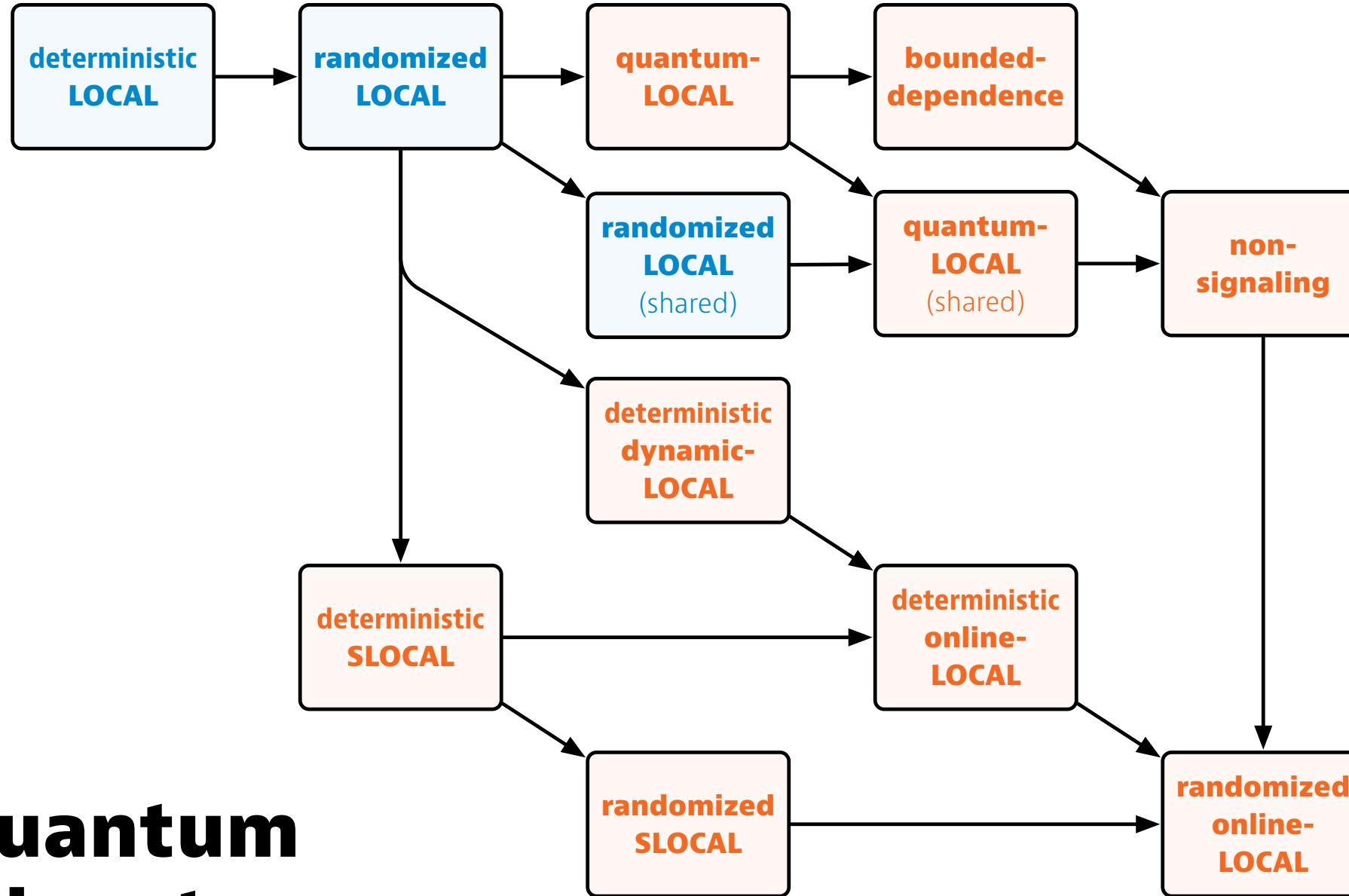


Free
symmetry
breaking

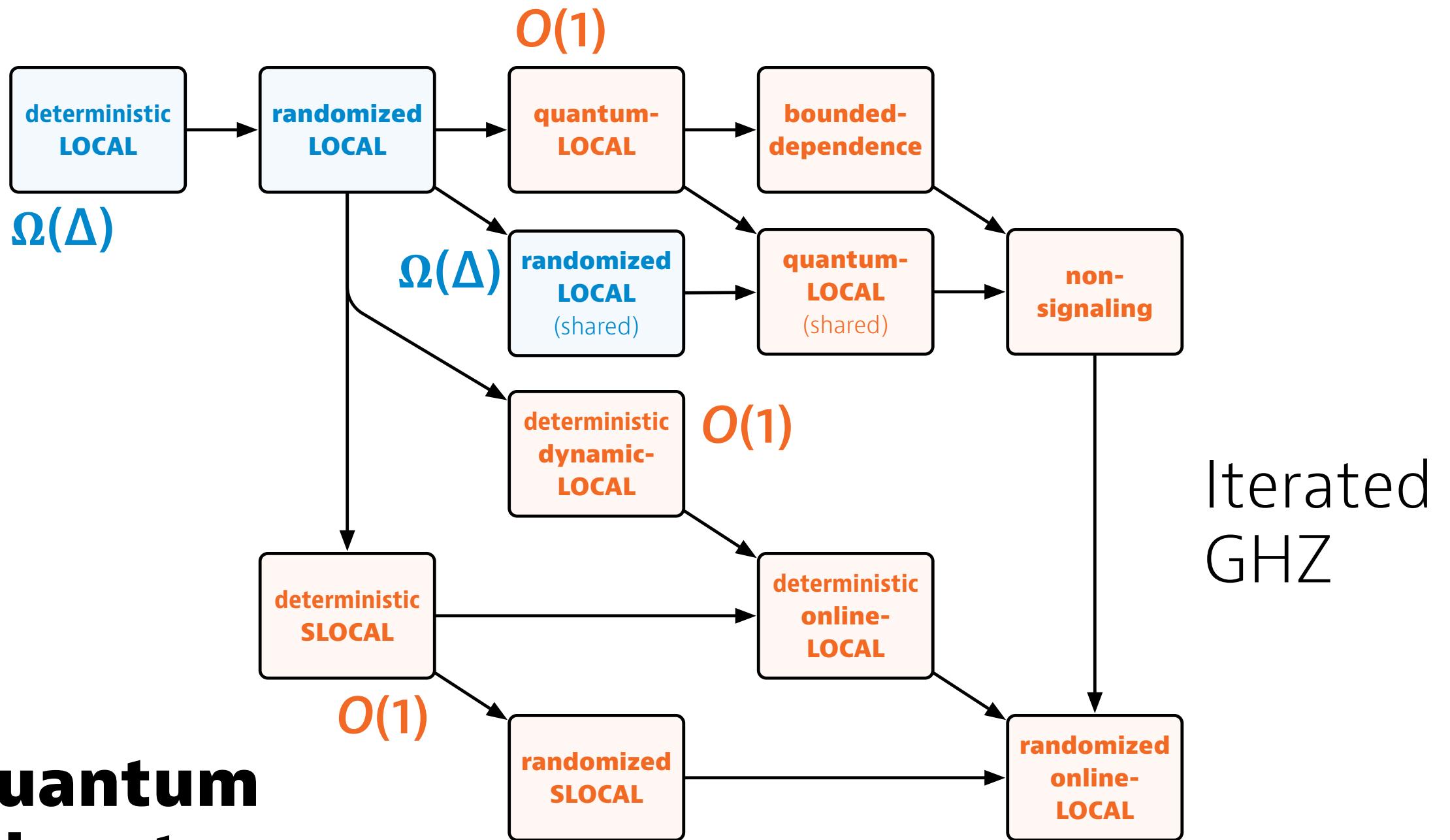


Online advantage

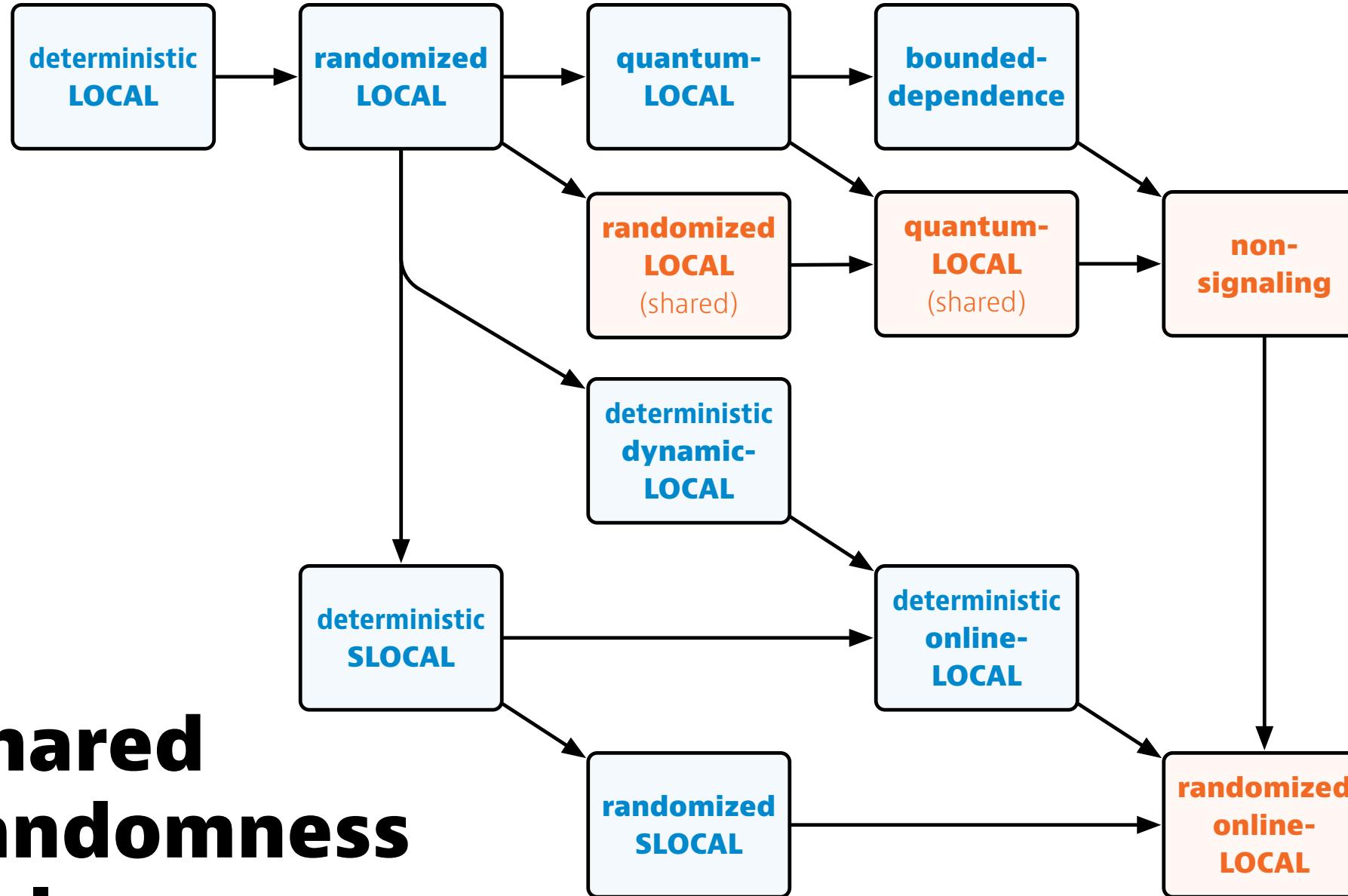




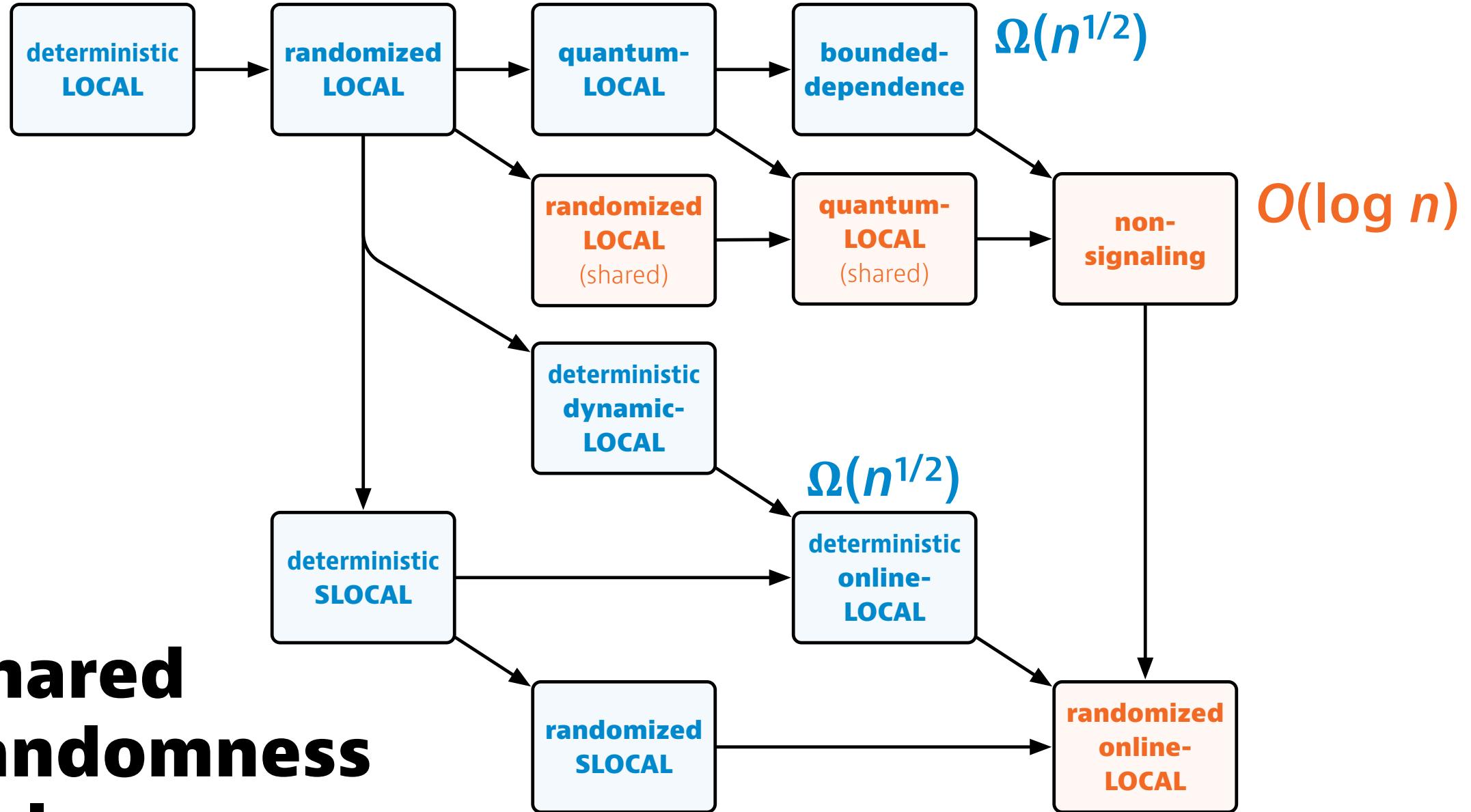
Quantum advantage



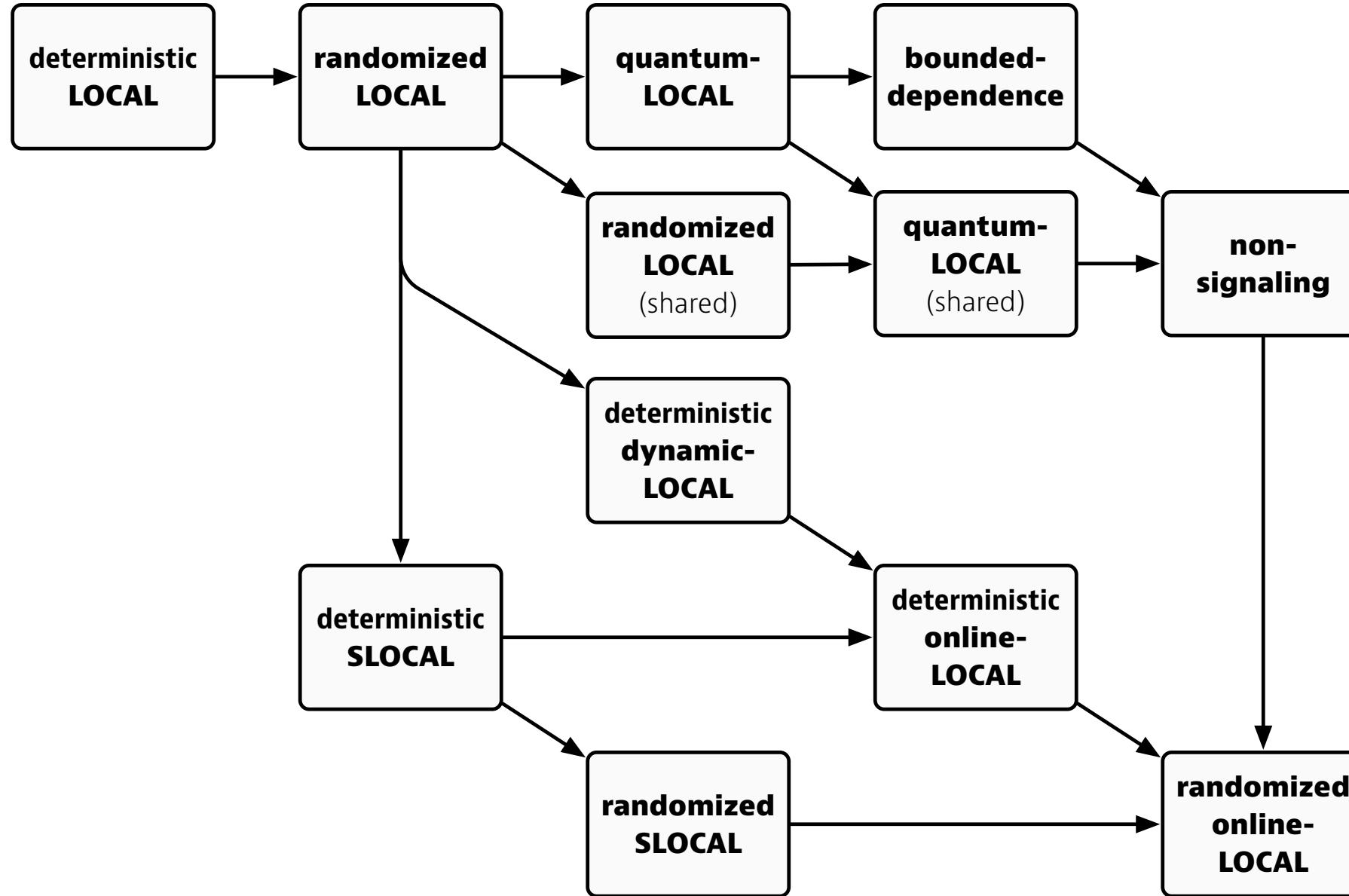
Quantum advantage



**Shared
randomness
helps**



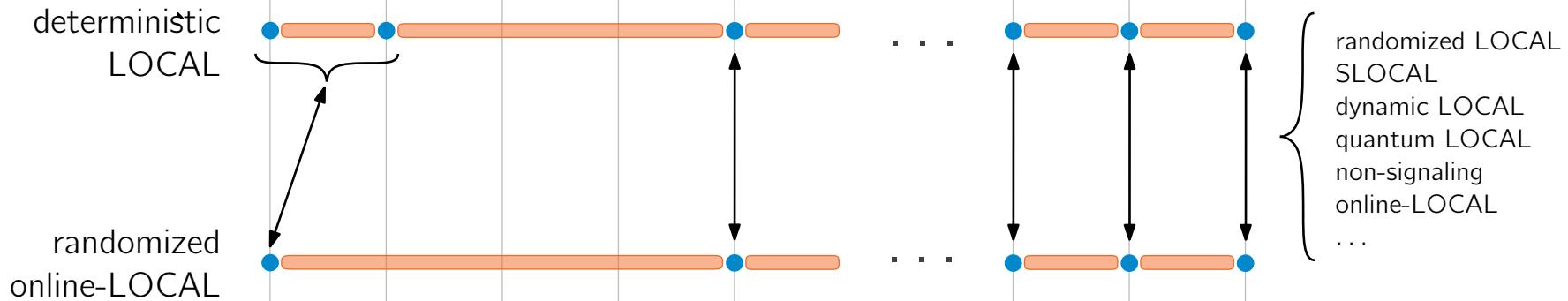
Shared randomness helps



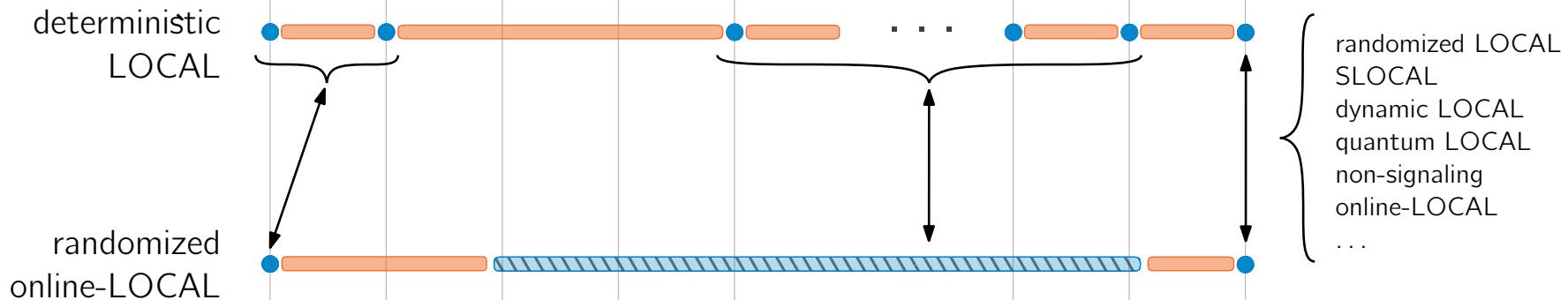
Some equivalences

$\Theta(1)$ $\Theta(\log^* n)$ $\Theta(\log \log \log n)$
 $\Theta(\log \log n)$ $\Theta(\log n)$... $\Theta(n^{1/3})$ $\Theta(n^{1/2})$ $\Theta(n)$

Rooted Regular Trees

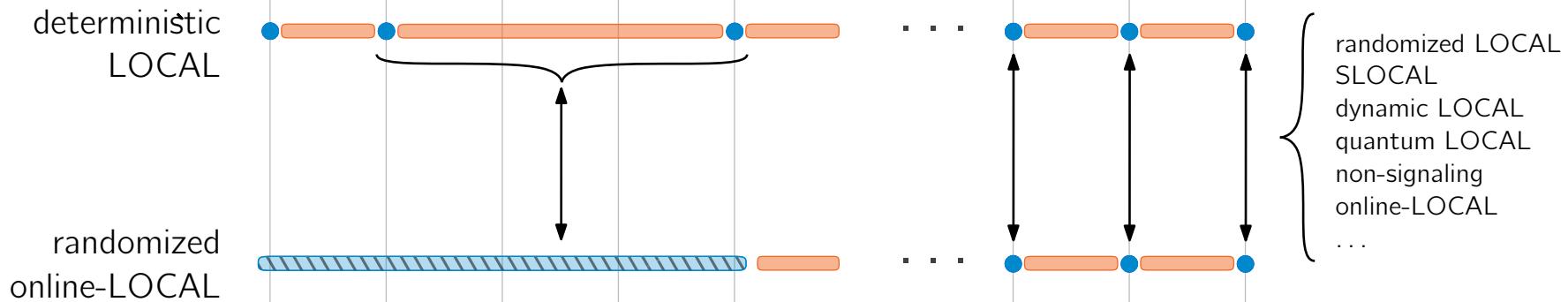


Rooted Trees

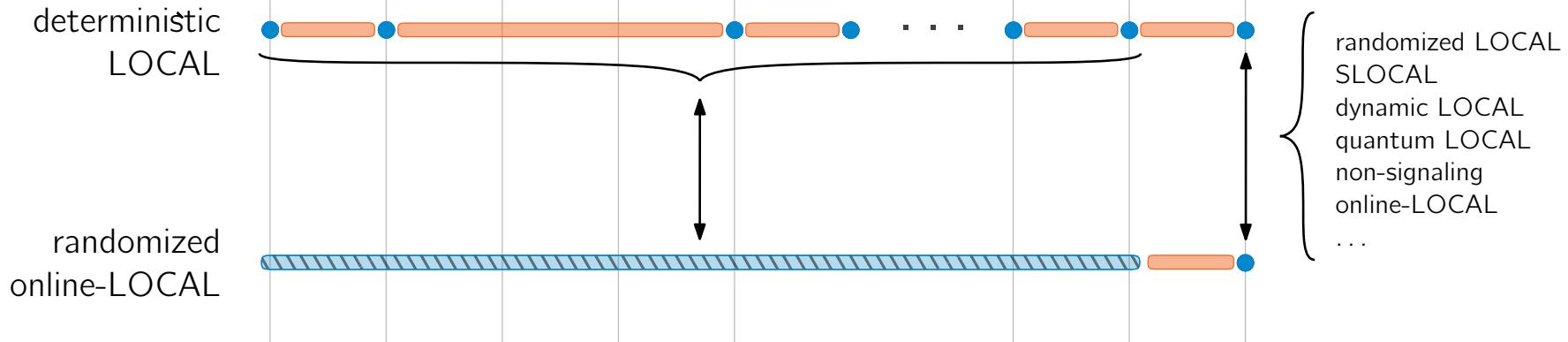


$\Theta(1)$ $\Theta(\log^* n)$ $\Theta(\log \log \log n)$
 $\Theta(\log \log n)$ $\Theta(\log n)$... $\Theta(n^{1/3})$ $\Theta(n^{1/2})$ $\Theta(n)$

Unrooted Regular Trees



Unrooted Trees



**Lots of open
questions**

Open questions

- **Does dynamic-LOCAL every help with LCLs?**
 - what is the complexity of *3-coloring grids* in dynamic-LOCAL?
 - what is the complexity of *3-coloring bipartite graphs* in dynamic-LOCAL?

Open questions

- **Can you make online-LOCAL algorithms promise-free?**
 - currently we need some kind of a *promise*:
“3-coloring grids”, “3-coloring in bipartite graphs”
 - promise-free online-LOCAL algorithms don't seem to be any stronger than SLOCAL algorithms?

Open questions

- **Does global memory ever help in trees?**
 - could we *simulate online-LOCAL algorithms in SLOCAL model* in unrooted trees?
 - could we simulate dynamic-LOCAL algorithms?
 - or can we separate these for some LCL in trees?

Open questions

- **How to prove lower bounds in dynamic-LOCAL or online-LOCAL?**
 - e.g. how to show that you *cannot 3-color trees with locality 10*?

Open questions

- **Is there any function-of- n distributed quantum advantage for LCLs?**
 - recall: we can show a function-of- Δ separation for a family of LCLs
 - can you ever have function-of- n quantum advantage for a single LCL?
 - example: can you *color cycles* in $O(1)$ rounds in *quantum-LOCAL*?

Open questions

- **Lower bounds for sinkless orientation**
 - deterministic LOCAL: $\approx \log n$
 - randomized LOCAL: $\approx \log \log n$
 - deterministic SLOCAL: $\approx \log \log n$
 - randomized SLOCAL: $\approx \log \log \log n$
 - dynamic-LOCAL: ***no lower bounds!***
 - online-LOCAL: ***no lower bounds!***
 - quantum-LOCAL: ***no lower bounds!***

Open questions

- **New lower-bound proof techniques?**
 - *existential graph-theoretic arguments* do not work e.g. in online-LOCAL
 - *round elimination* do not work e.g. in quantum-LOCAL
 - current best hope: new *simulation results*?

- arXiv:2109.06593 **Locality in online, dynamic, sequential, and distributed graph algorithms** (*ICALP* 2023)
- arXiv:2307.09444 **No distributed quantum advantage for approximate graph coloring** (*STOC* 2024)
- arXiv:2403.01903 **Online locality meets distributed quantum computing**
- arXiv:2407.05445 **Shared randomness helps with local distributed problems**
- arXiv:2409.13795 **Local problems in trees across a wide range of distributed models** (*OPODIS* 2024)
- arXiv:2411.03240 **Distributed quantum advantage for local problems**