Jukka Suomela Aalto University

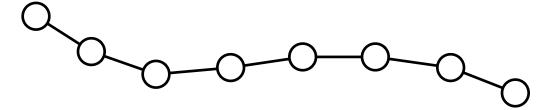
Locality in online, dynamic, sequential, and distributed graph algorithms

Joint work with:

- Amirreza Akbari
- Navid Eslami
- Henrik Lievonen
- Darya Melnyk
- Joona Särkijärvi

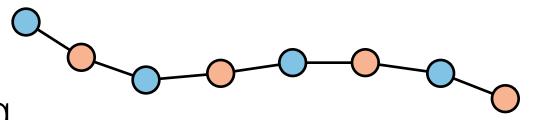
arxiv.org/abs/2109.06593

Informal introduction to locality

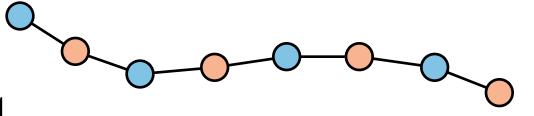


• Given: a path graph

• Find: a proper 2-coloring



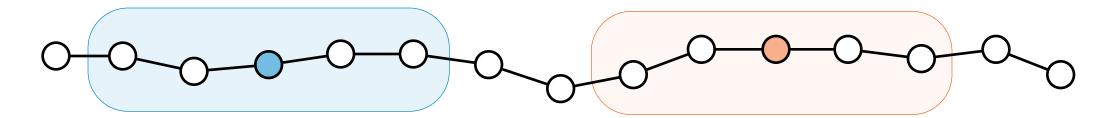
• Given: a path graph



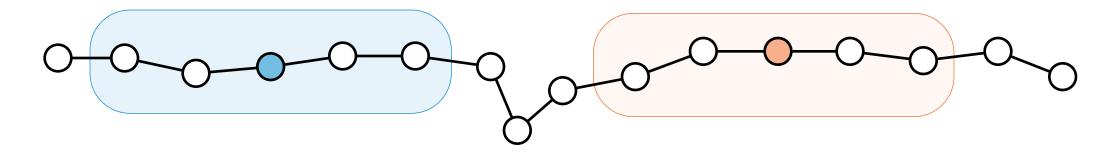
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 - say, radius–o(n) neighborhood

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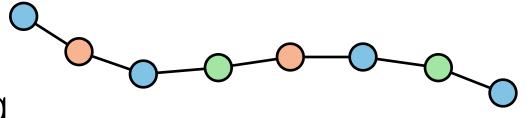
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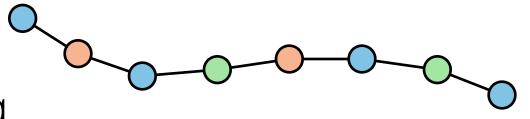
Conclusion:

Paths can't be 2-colored with any local strategy

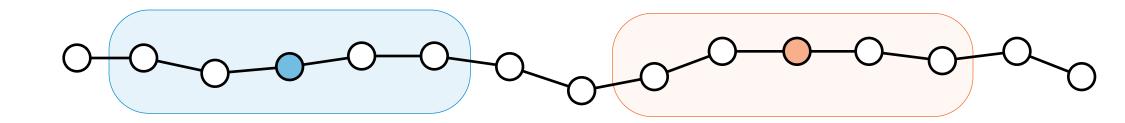
... and it doesn't really depend on exactly how we define "local"

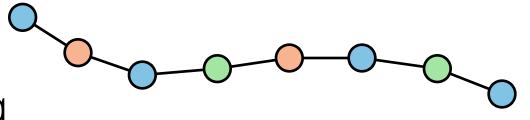


- Find: a proper 3-coloring
- Restriction: the color of each node only depends on its local neighborhood

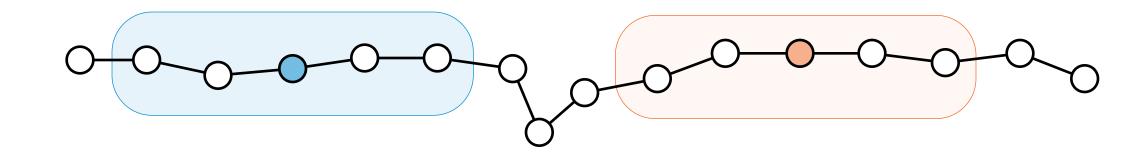


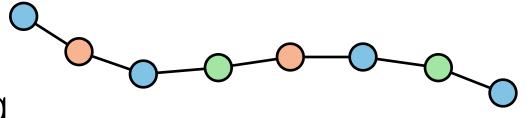
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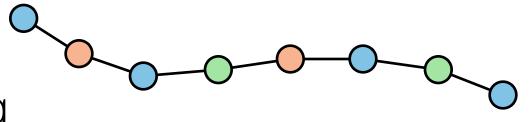


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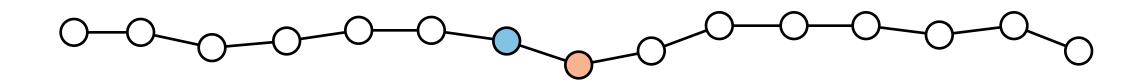




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- randomness
- unique node identifiers
- sequential ordering ...

Nodes labeled with (small) unique identifiers:

locality $\approx \frac{1}{2} \log^* n$

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[Cole & Vishkin 1986, Linial 1992, Naor 1991]

Four models of computing

LOCAL distributed, parallel



SLOCAL

distributed, sequential

LOCAL distributed, parallel

online LOCAL centralized

LOCAL distributed, parallel

Each node in parallel:

- looks at its radius-T neighborhood
- picks its output based on this information

(nodes have unique identifiers)

SLOCAL

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Each node in a sequential, adversarial order:

- looks at its radius-T neighborhood
- picks its output & state based on this information

LOCAL distributed, parallel

online LOCAL centralized

LOCAL distributed, parallel



Graph **constructed** by an adversary that adds nodes and edges one by one

We can see everything

We can **change** our output only within distance *T* from a point of change

LOCAL distributed, parallel



LOCAL distributed, parallel



Some unknown input graph is revealed piece by piece:

- adversary points at a node v
- we can see the radius-T neighborhood of v
- we have to choose the label for v

We can **remember** everything



LOCAL distributed, parallel



Genuinely different models

LOCAL distributed, parallel



coloring

SLOCAL

distributed, sequential

LOCAL distributed, parallel

online LOCAL centralized

dynamic LOCAL

centralized

coloring

SLOCAL

distributed, sequential

LOCAL distributed, parallel

online LOCAL centralized

cycle detection dynamic LOCAL centralized

coloring

SLOCAL

distributed, sequential

LOCAL distributed, parallel

dynamic LOCAL

centralized

online LOCAL

centralized

leader election

cycle detection

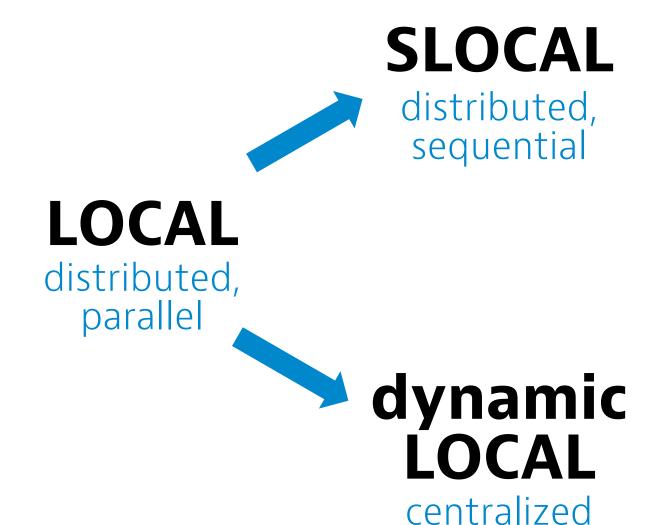
Closely related models



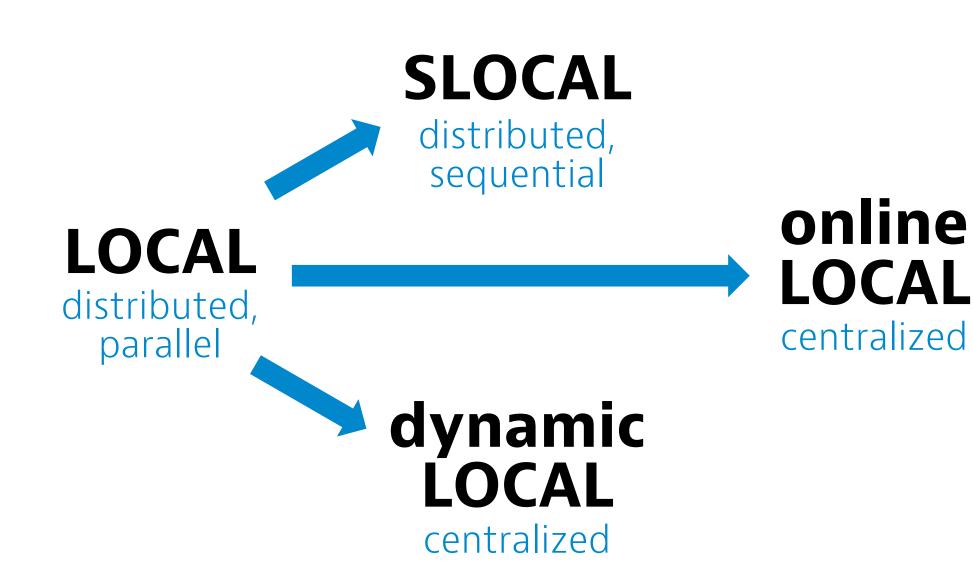




dynamic LOCAL centralized



online LOCAL centralized





distributed, sequential

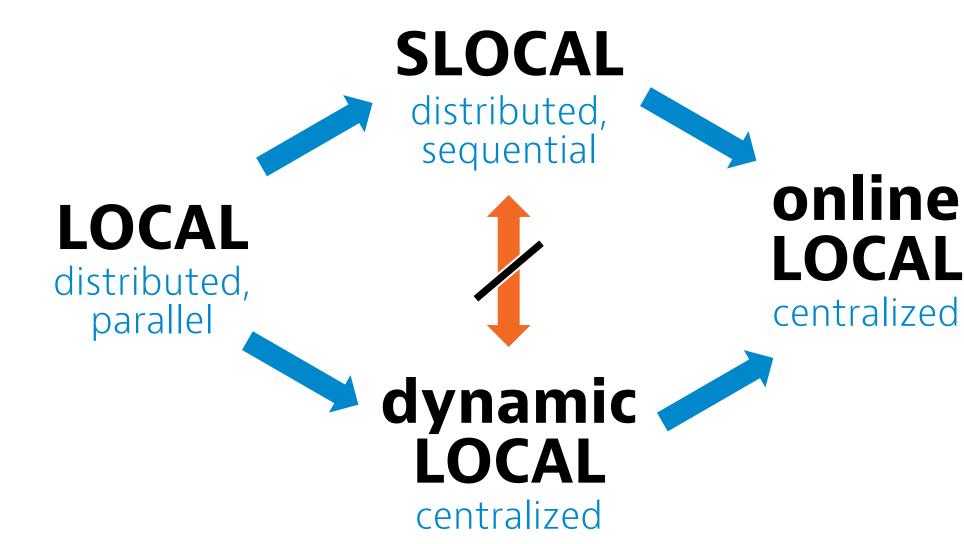
LOCAL distributed, parallel



centralized

dynamic LOCAL

centralized

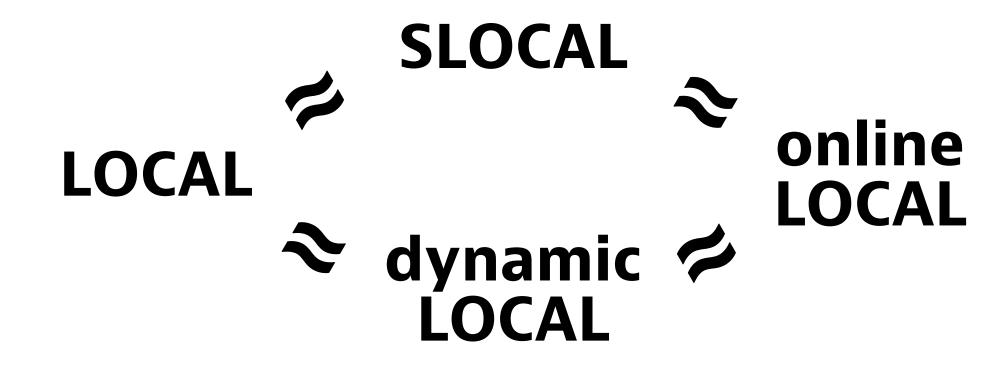


Collapse in rooted trees

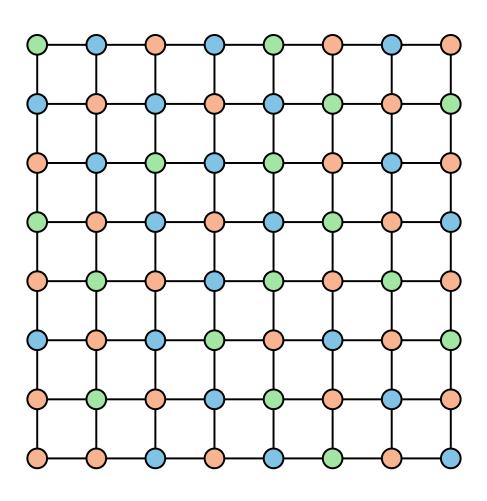
LCLs in rooted trees

- Rooted regular trees
- Locally checkable labelings (LCLs)
 - solution valid if it "looks good everywhere"
 - example: 3-coloring
- In this setting all models equally strong!

LCLs in rooted trees



Case study: grids



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 - online-LOCAL: O(log n)

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- 4-coloring: local in all models (hard to see)
- 3-coloring:
 - LOCAL, SLOCAL: global
 - online-LOCAL: $O(\log n)$ is this tight?
 - dynamic-LOCAL: open

