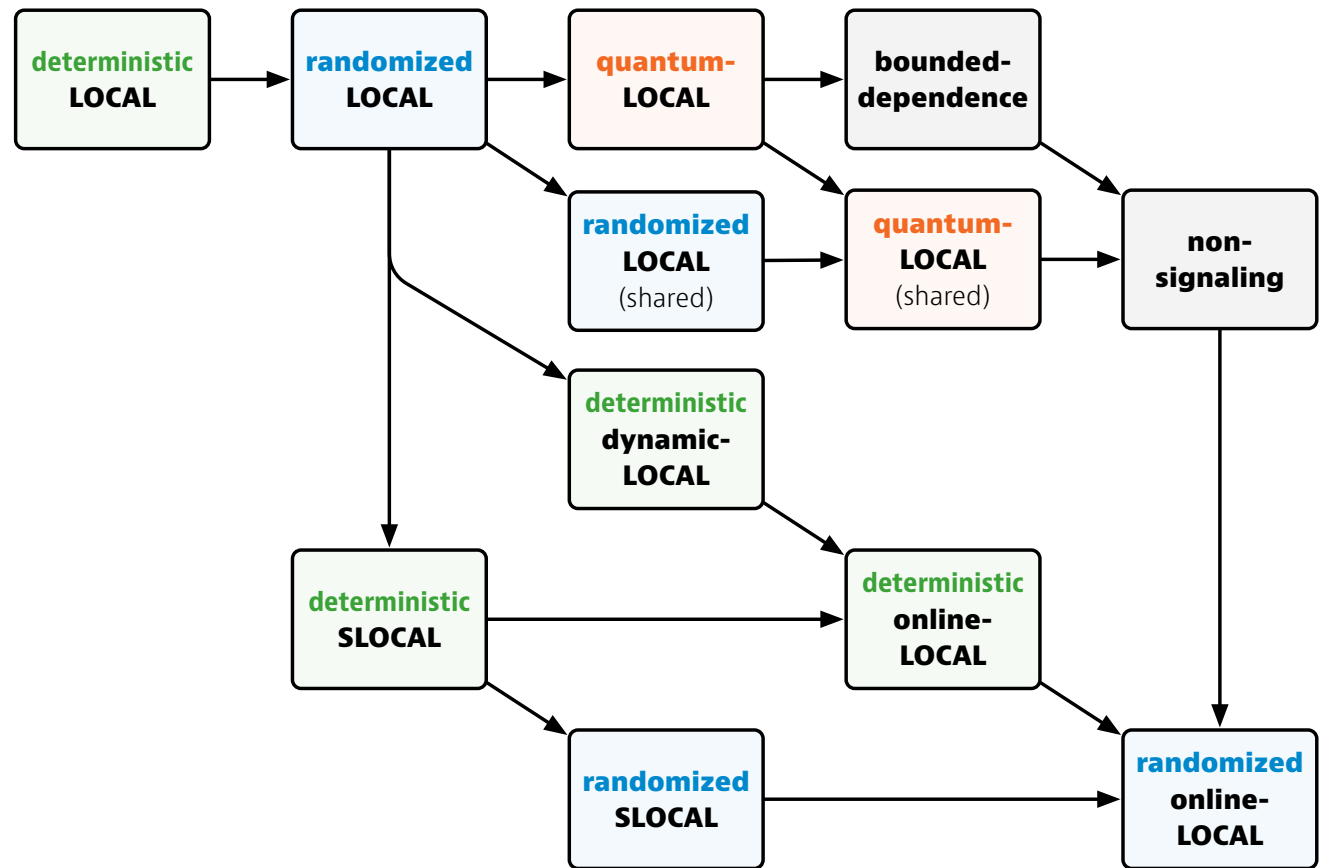


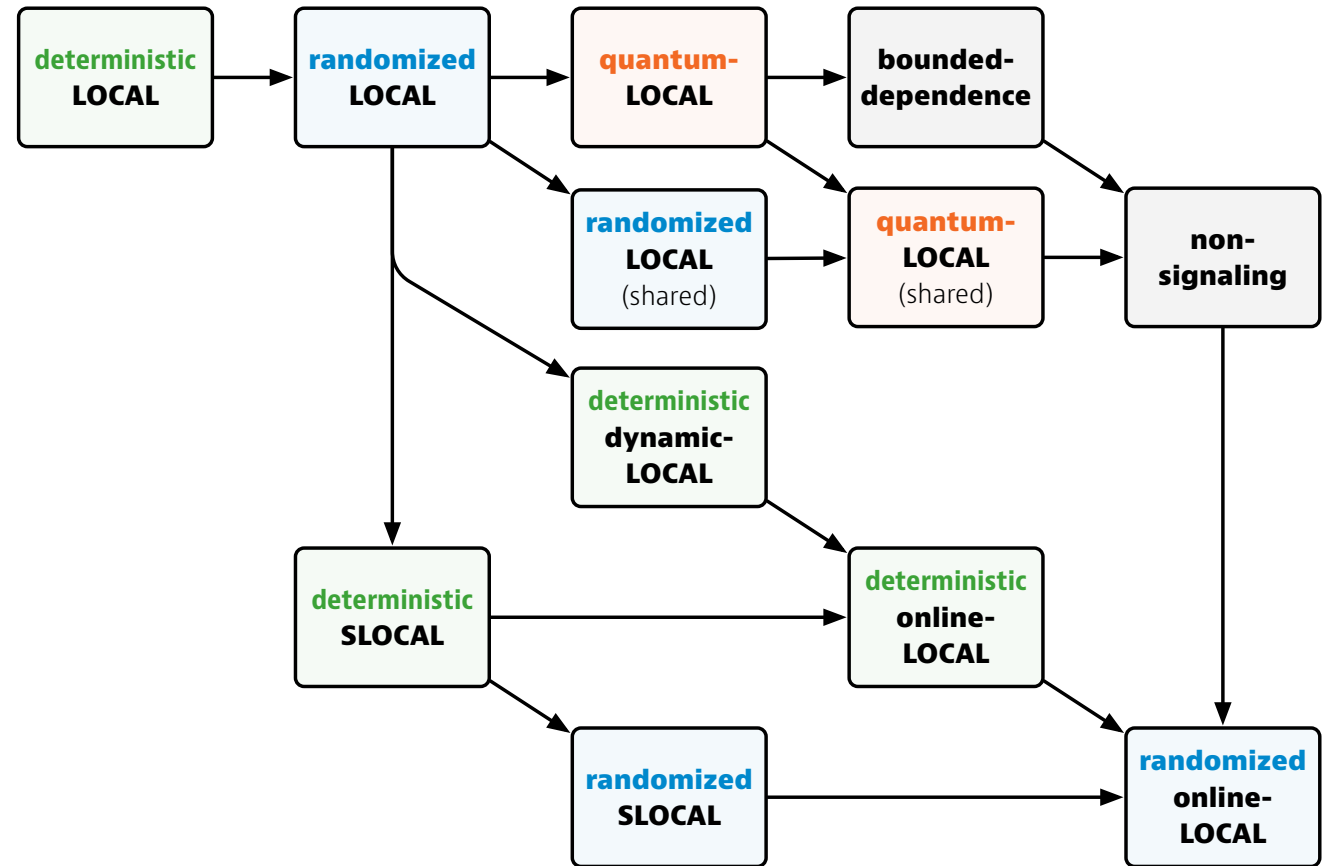
*Fun stuff &
latest news:
Locality
in all kinds
of weird settings*



Jukka Suomela
Aalto University

Joint work with e.g.

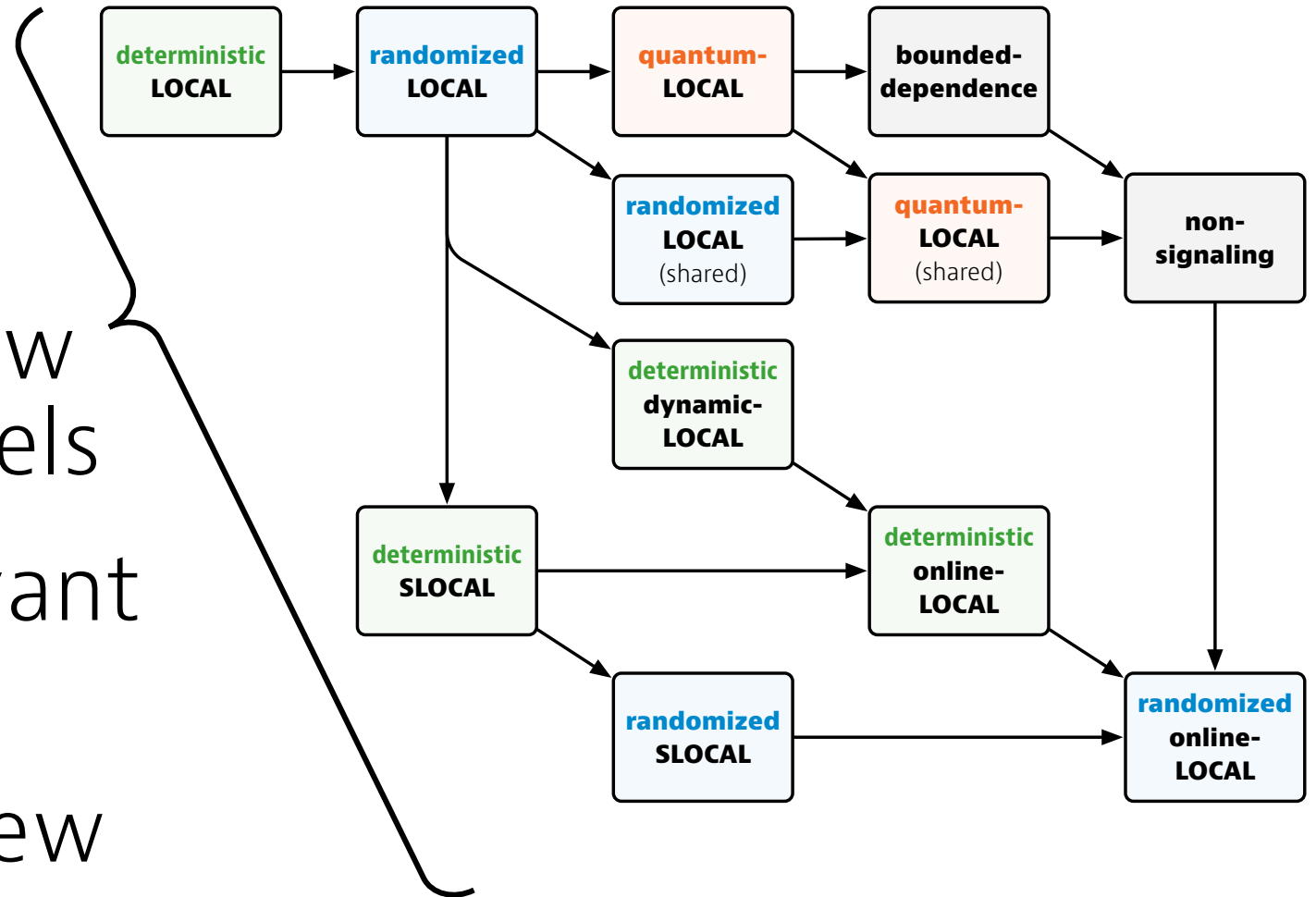
- Amirreza Akbari
- Francesco d'Amore
- Xavier Coiteux-Roy
- Navid Eslami
- Rishikesh Gajjala
- Fabian Kuhn
- François Le Gall
- Henrik Lievonen
- Darya Melnyk
- Augusto Modanese
- Shreyas Pai
- Marc-Olivier Renou
- Václav Rozhoň
- Gustav Schmid
- Joonas Särkijärvi

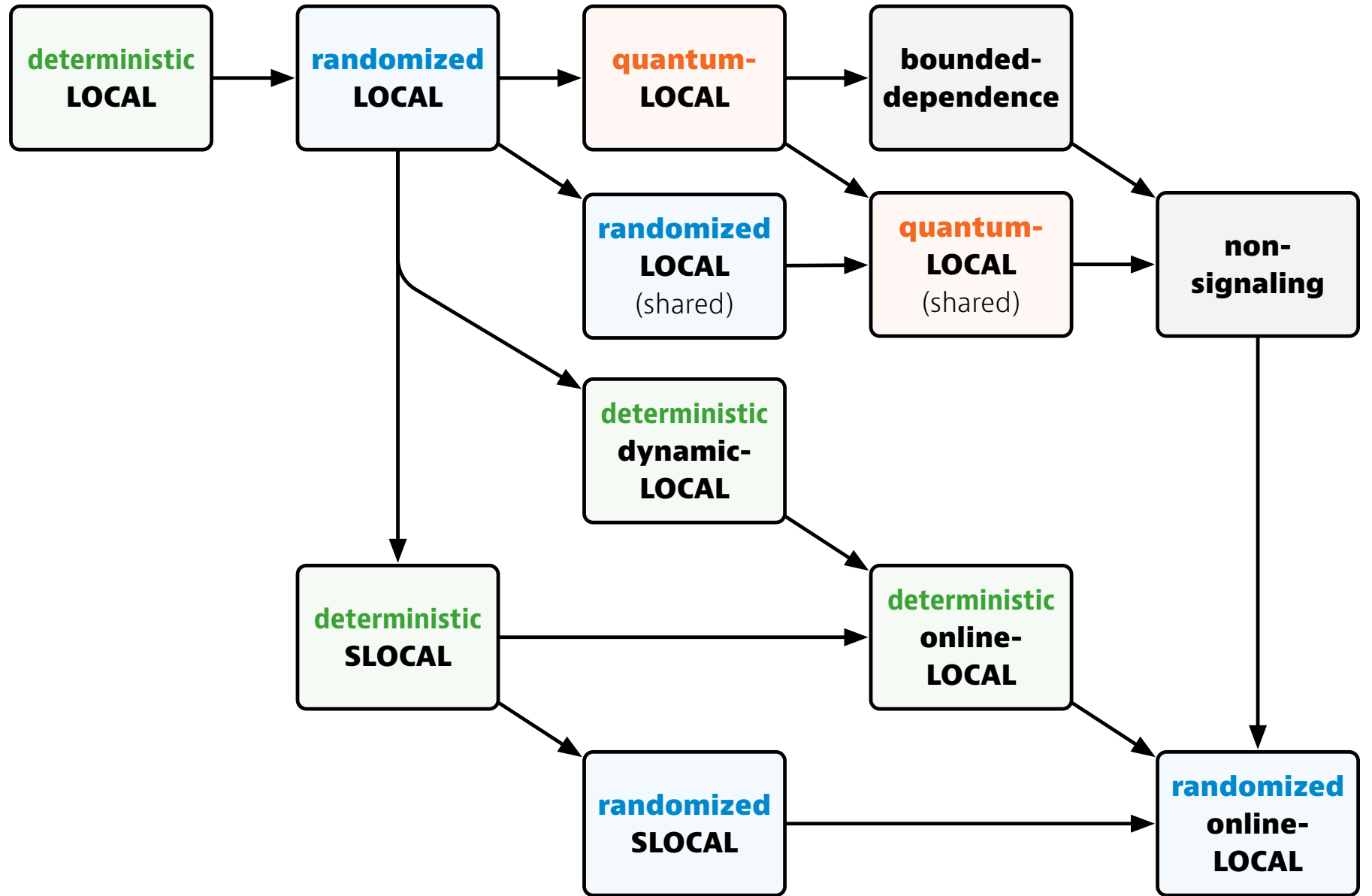


+ many results & concepts by others

Goals

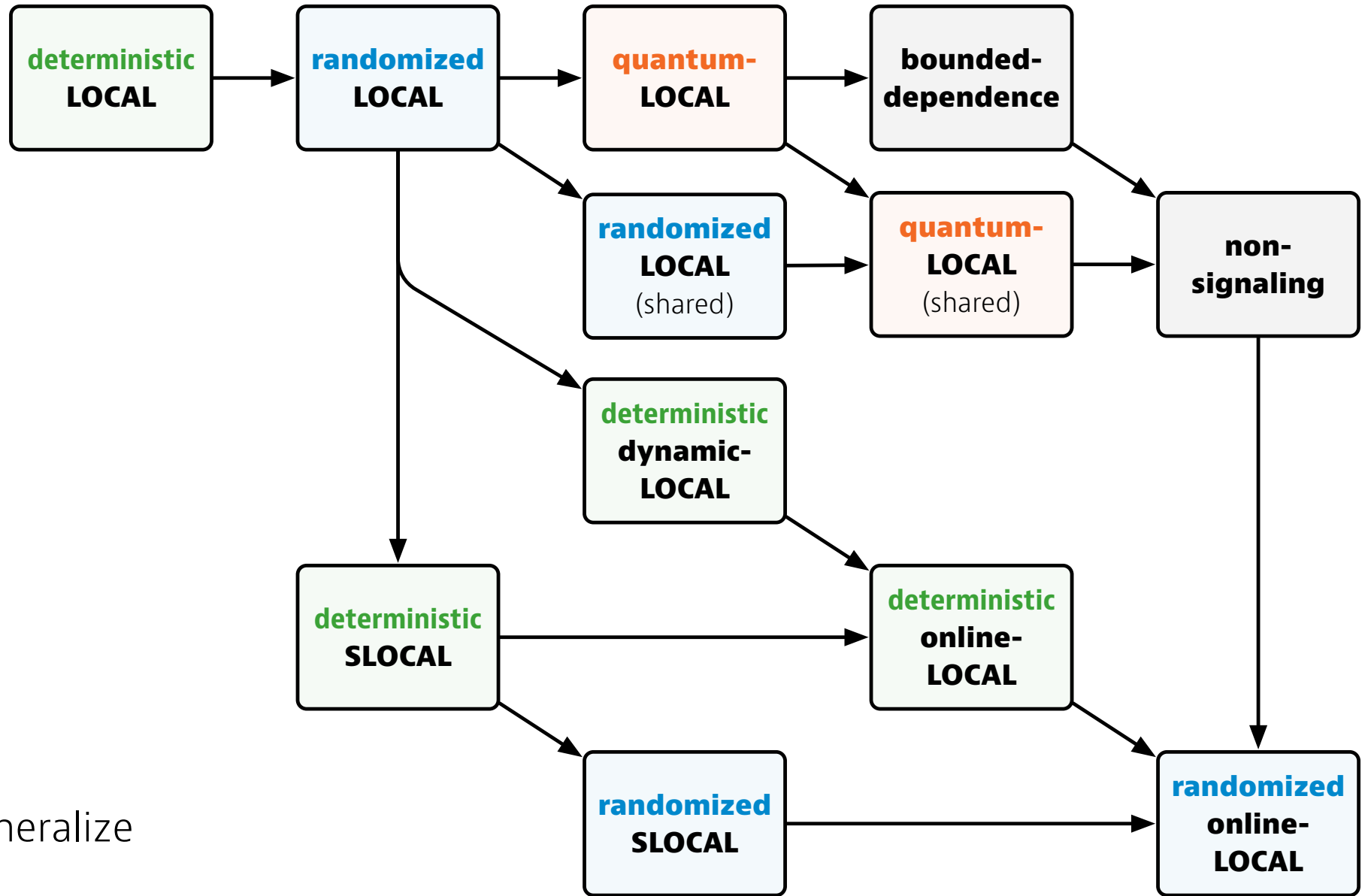
1. Informal overview of all these models
2. Why is this relevant and interesting?
3. Inspiration for new connections with e.g. measurable combinatorics?

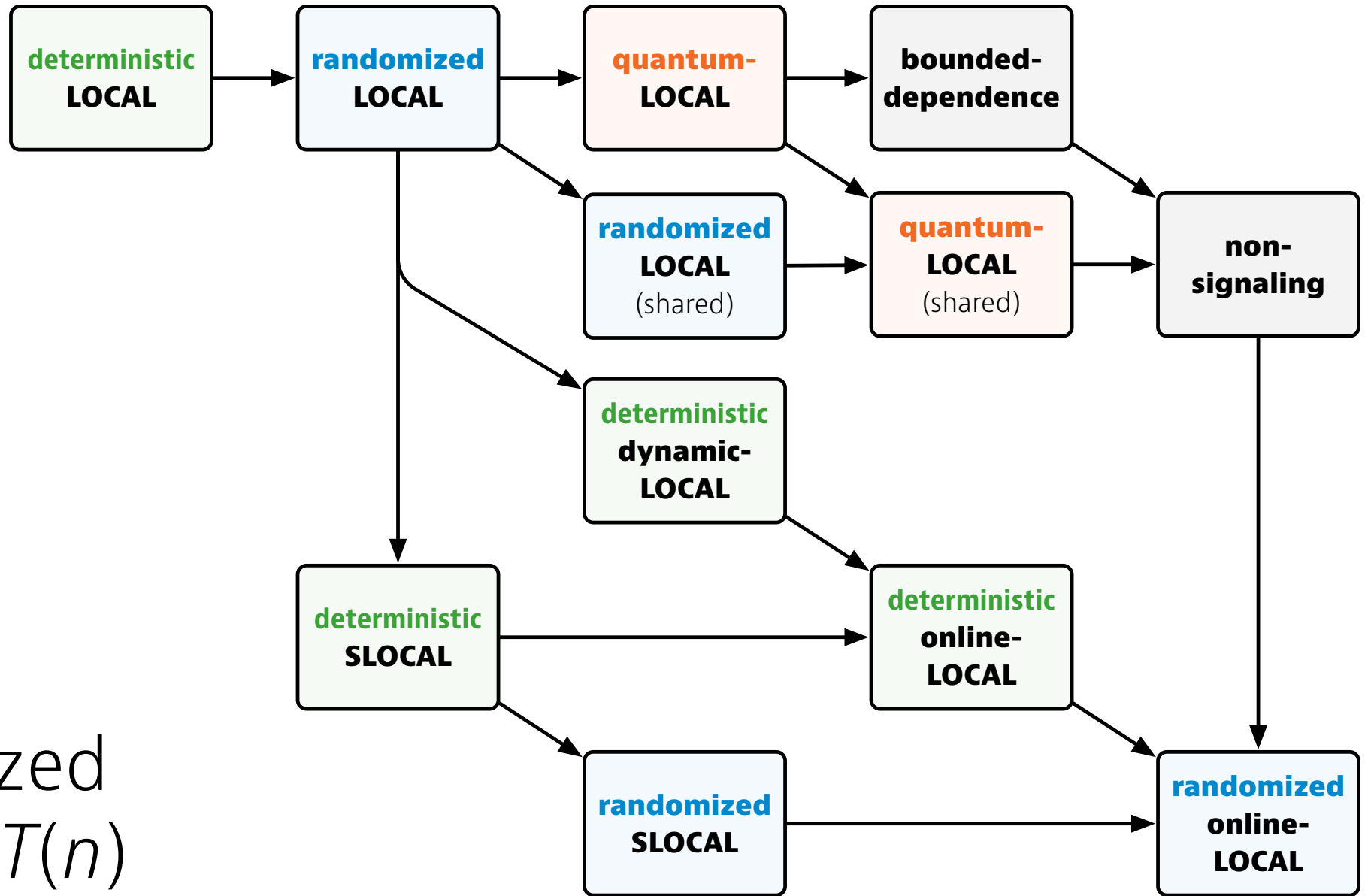




Context:
locally
checkable
labelings
(LCL)

but many things generalize
far beyond LCLs

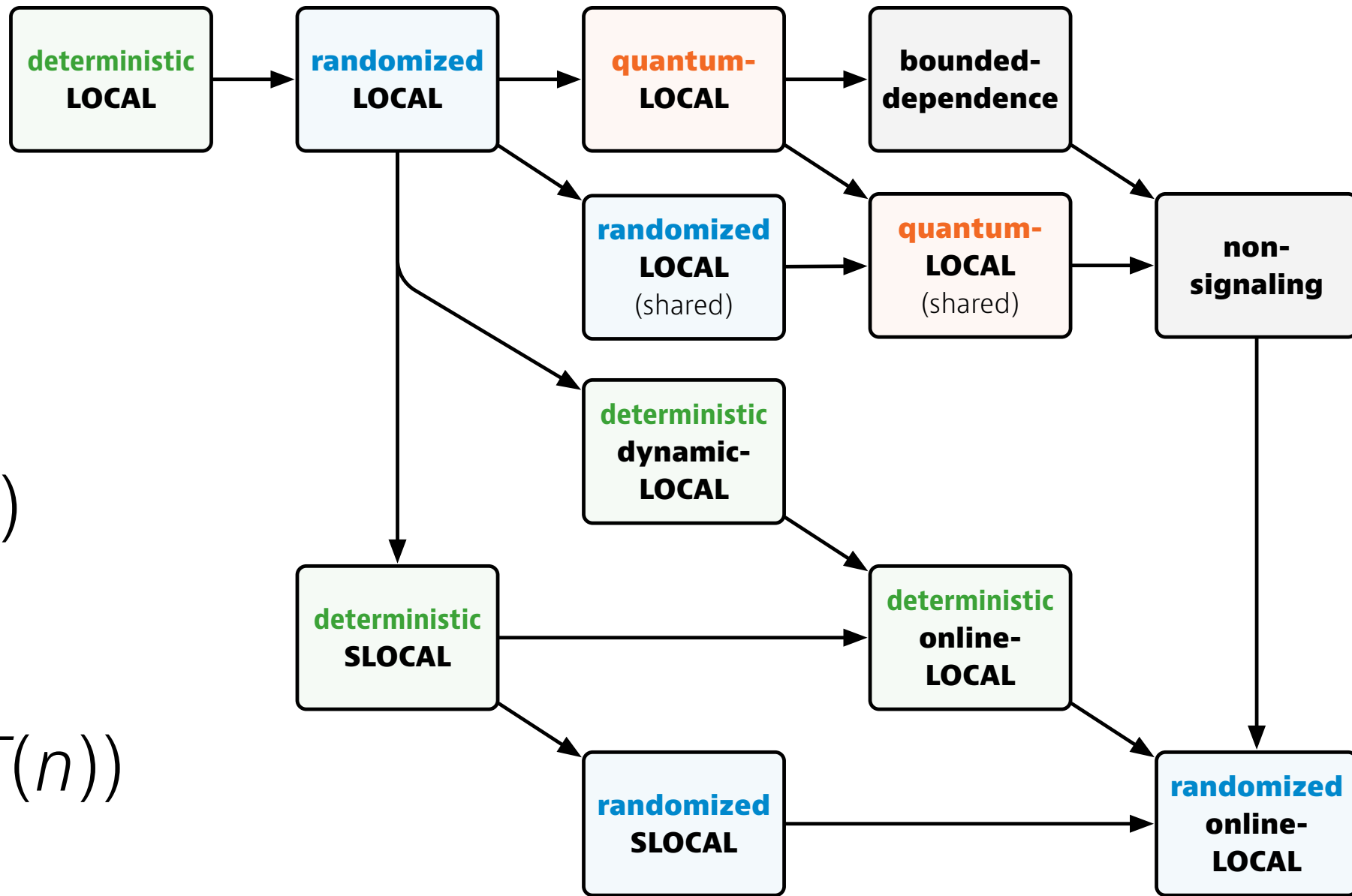




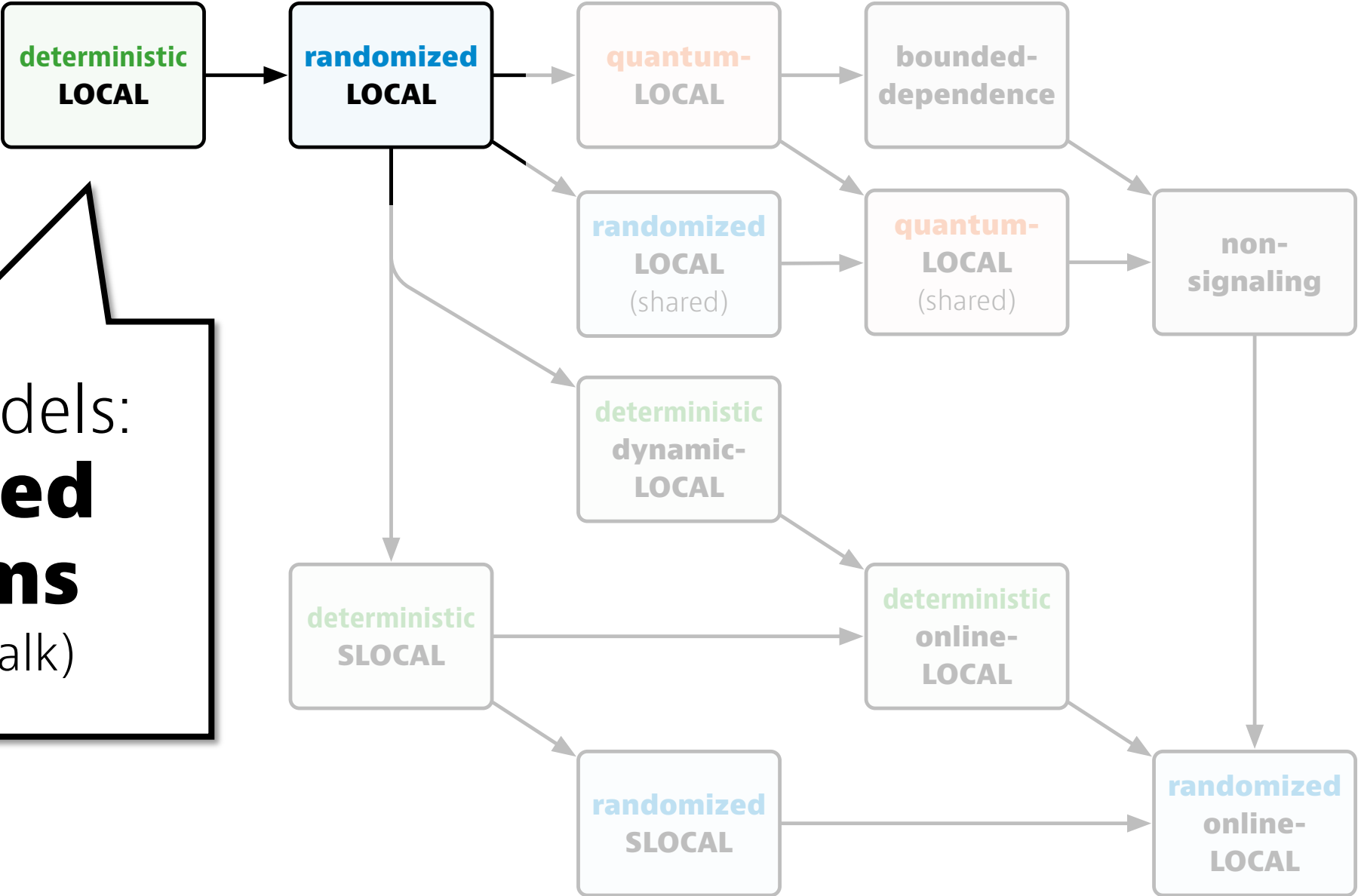
Each box:
 model of
 computing
 parameterized
 by locality $T(n)$

A → B:

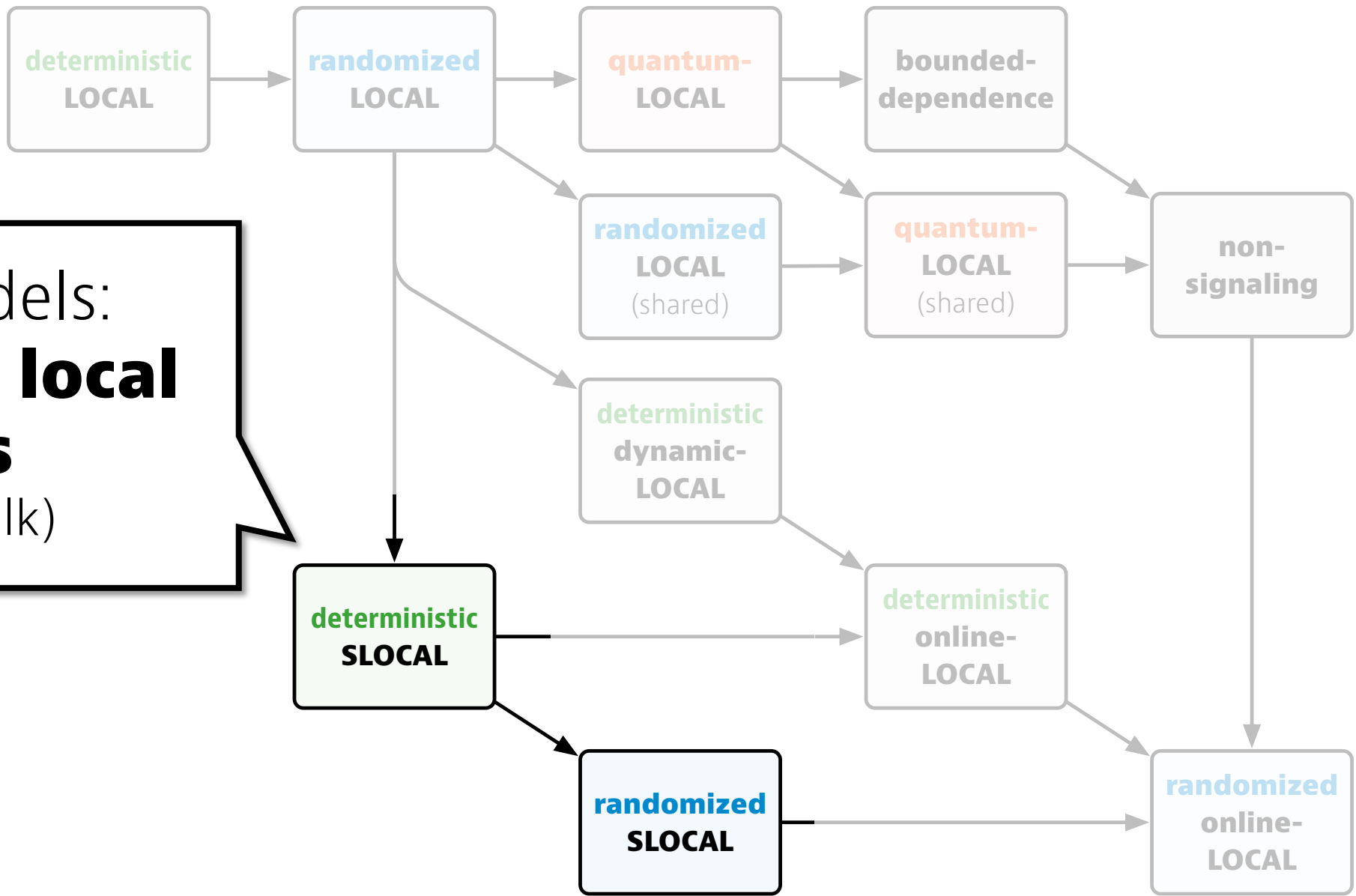
locality $T(n)$
in model A
implies
locality $O(T(n))$
in model B



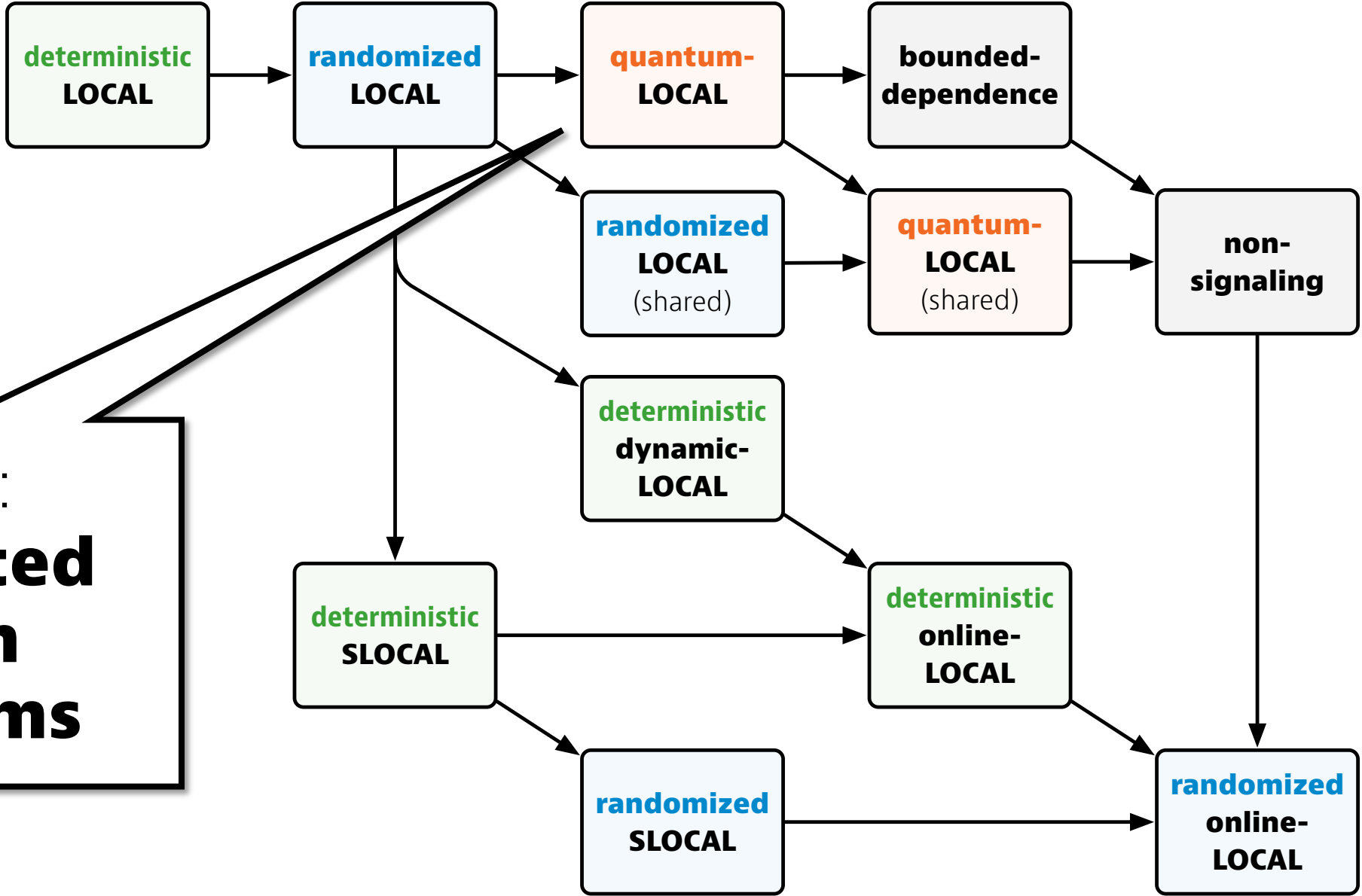
Familiar models:
distributed algorithms
(recall Vašek's talk)



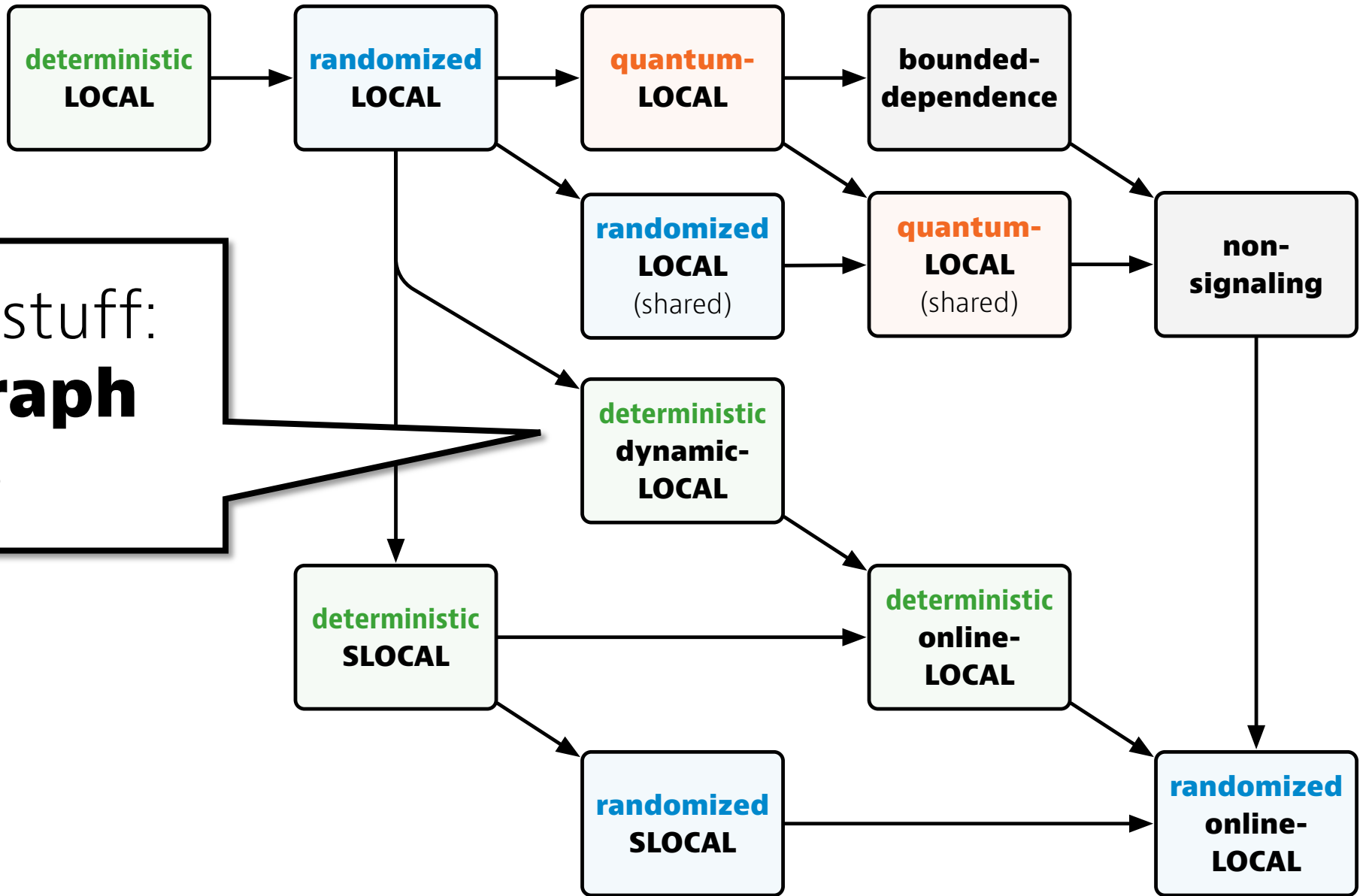
Familiar models:
sequential local algorithms
(recall Vašek's talk)



Cool stuff:
**distributed
quantum
algorithms**

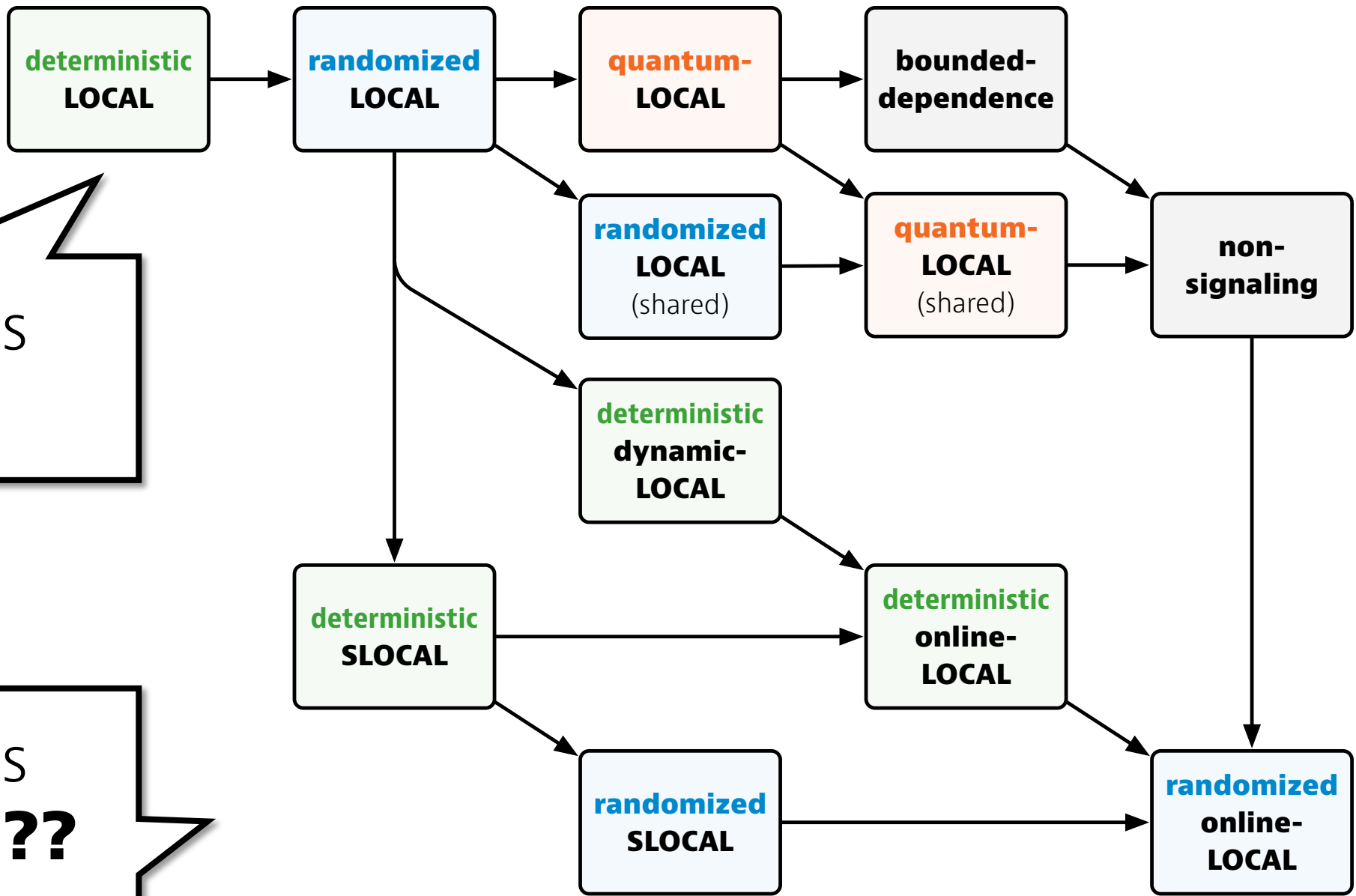


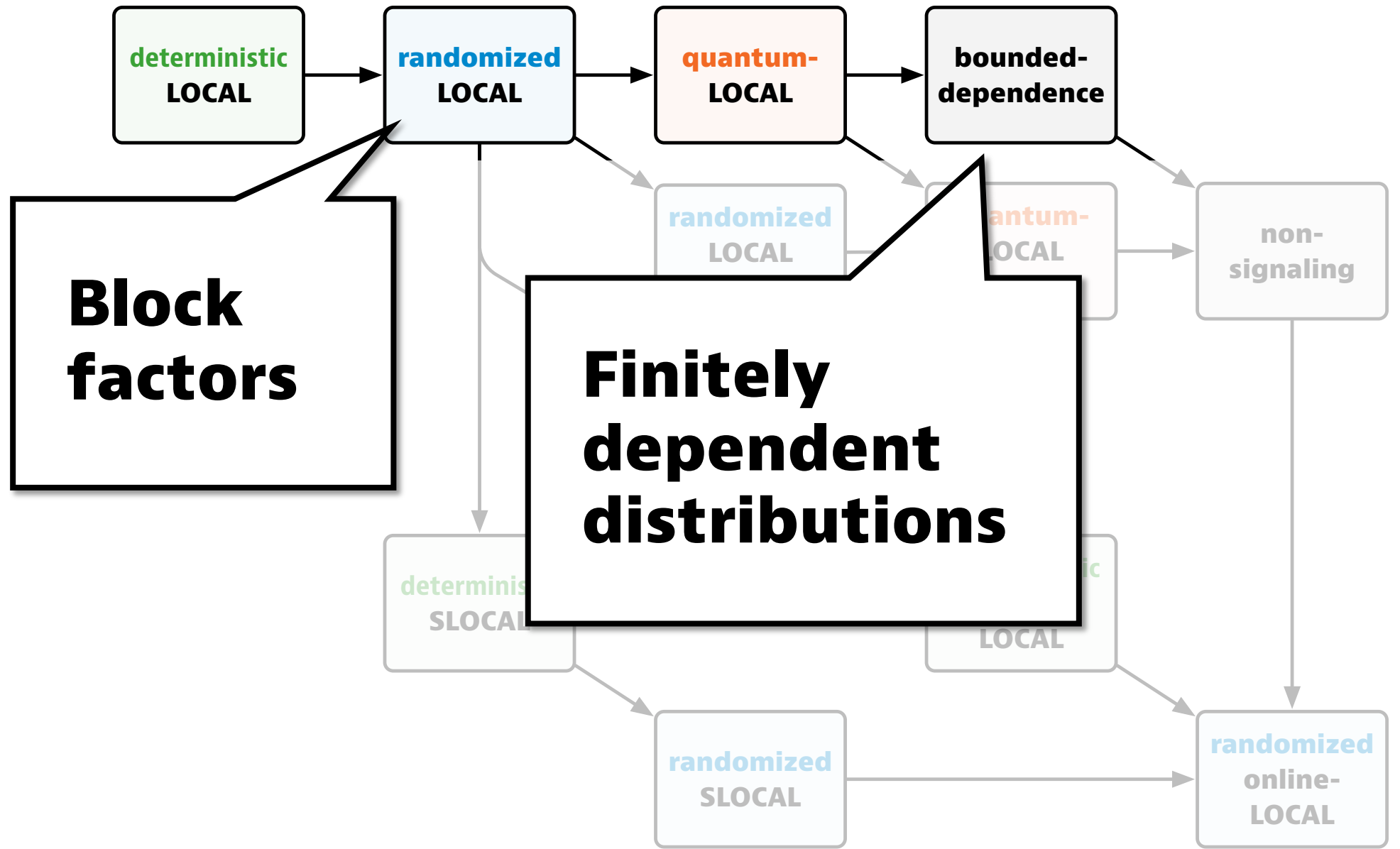
Unexpected stuff:
dynamic graph algorithms



Connections with **Borel**

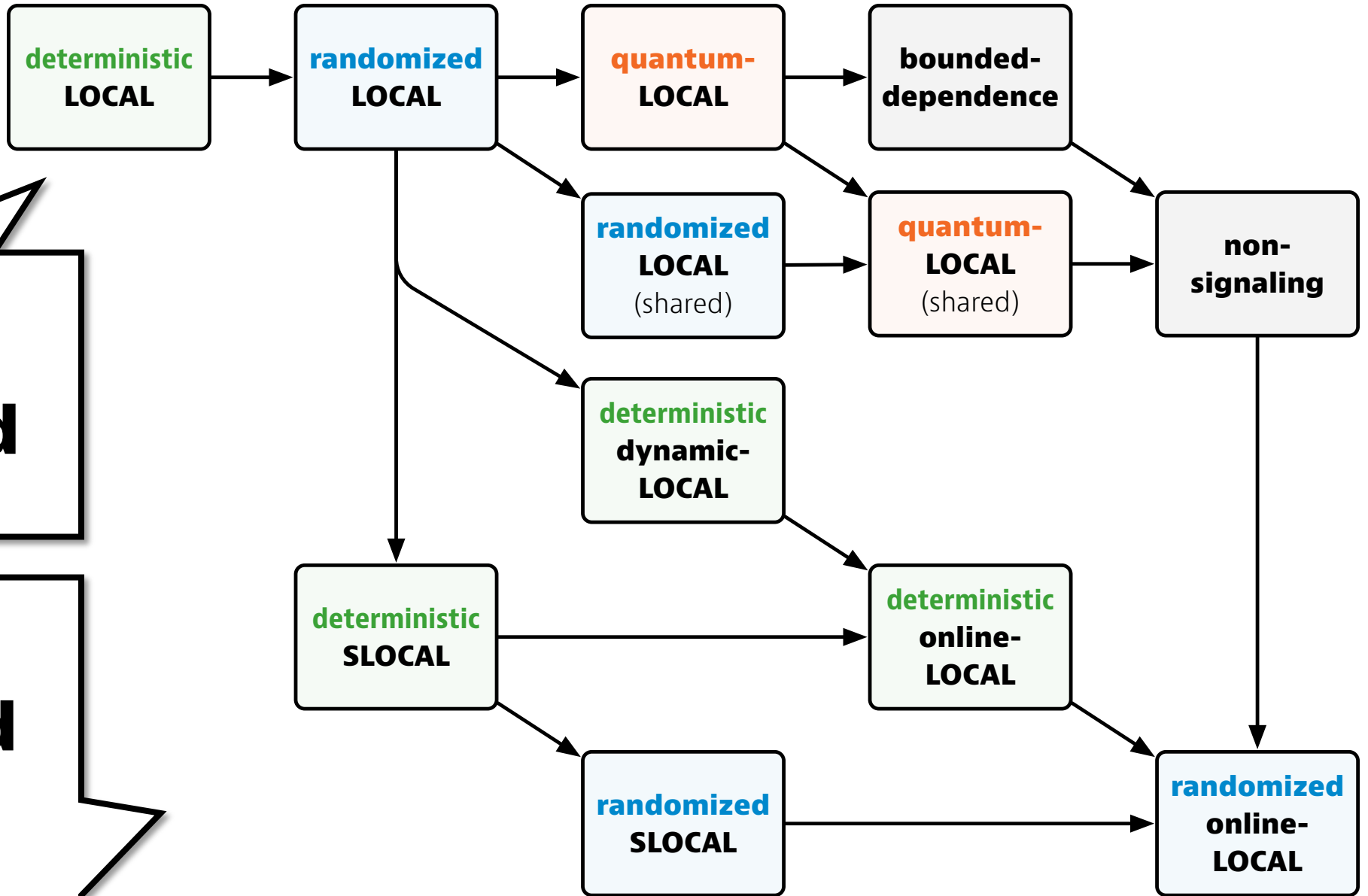
Connections with **Borel??**



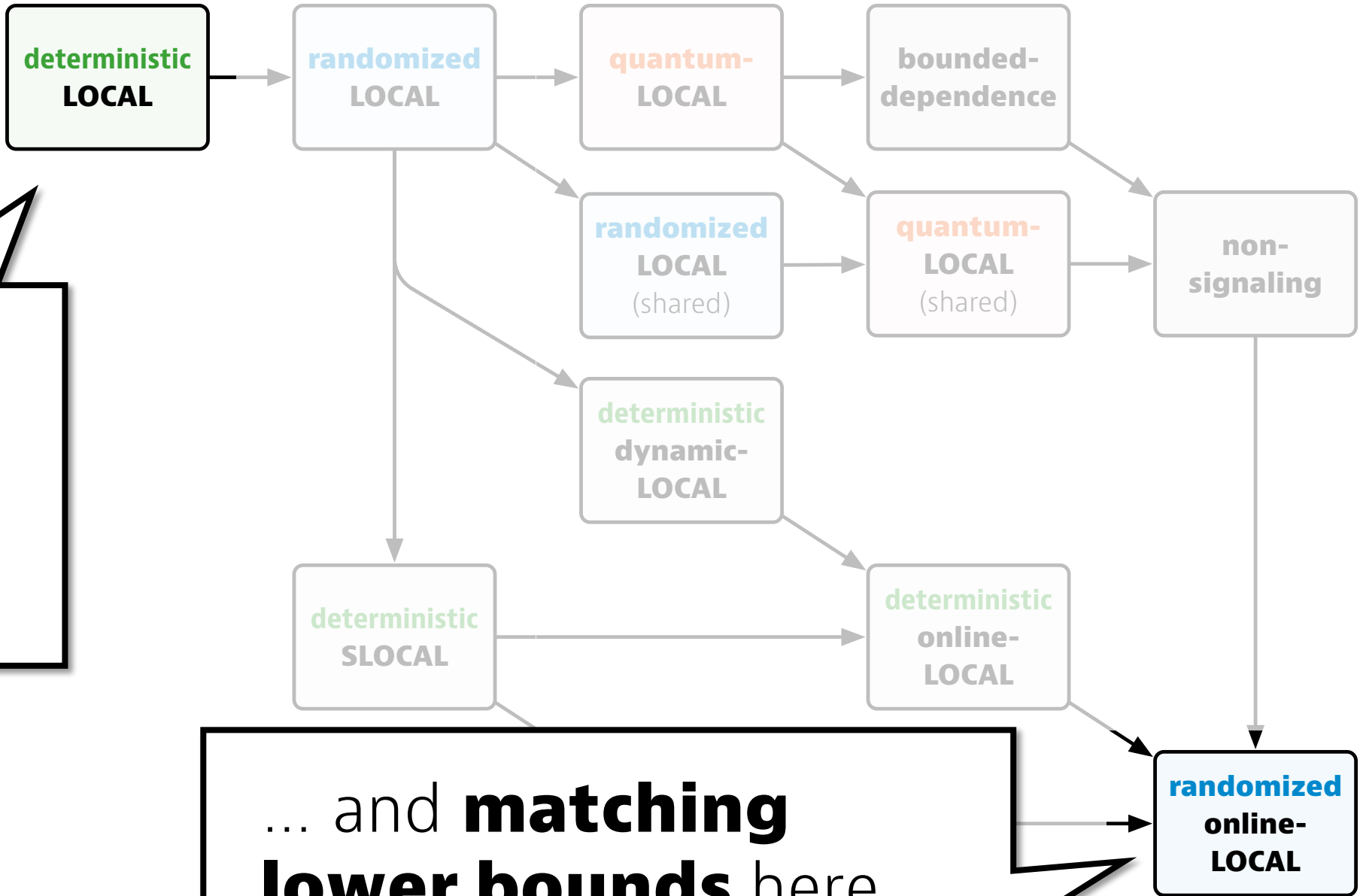


This is **well understood**

Could we **understand** also all the rest??

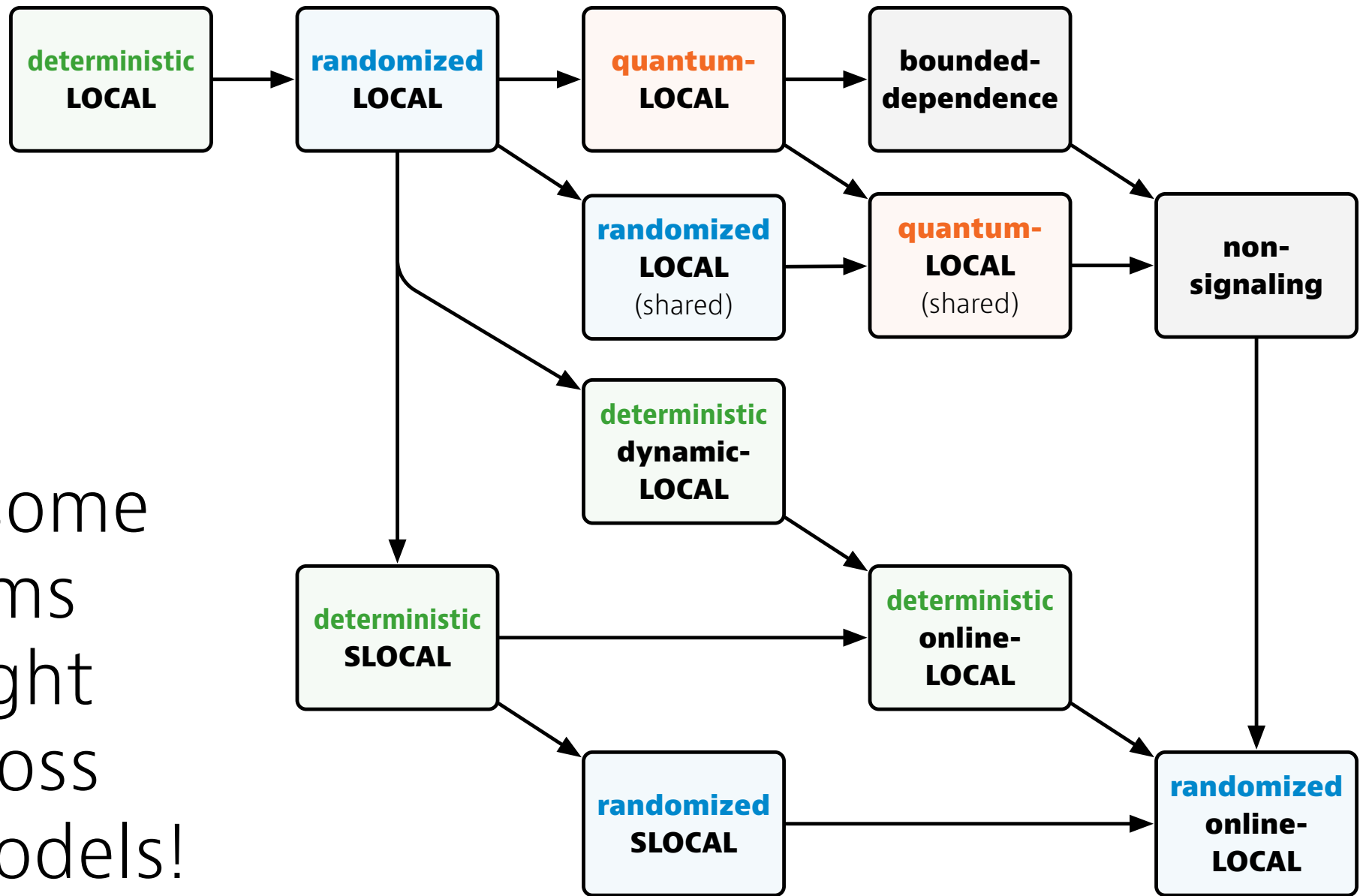


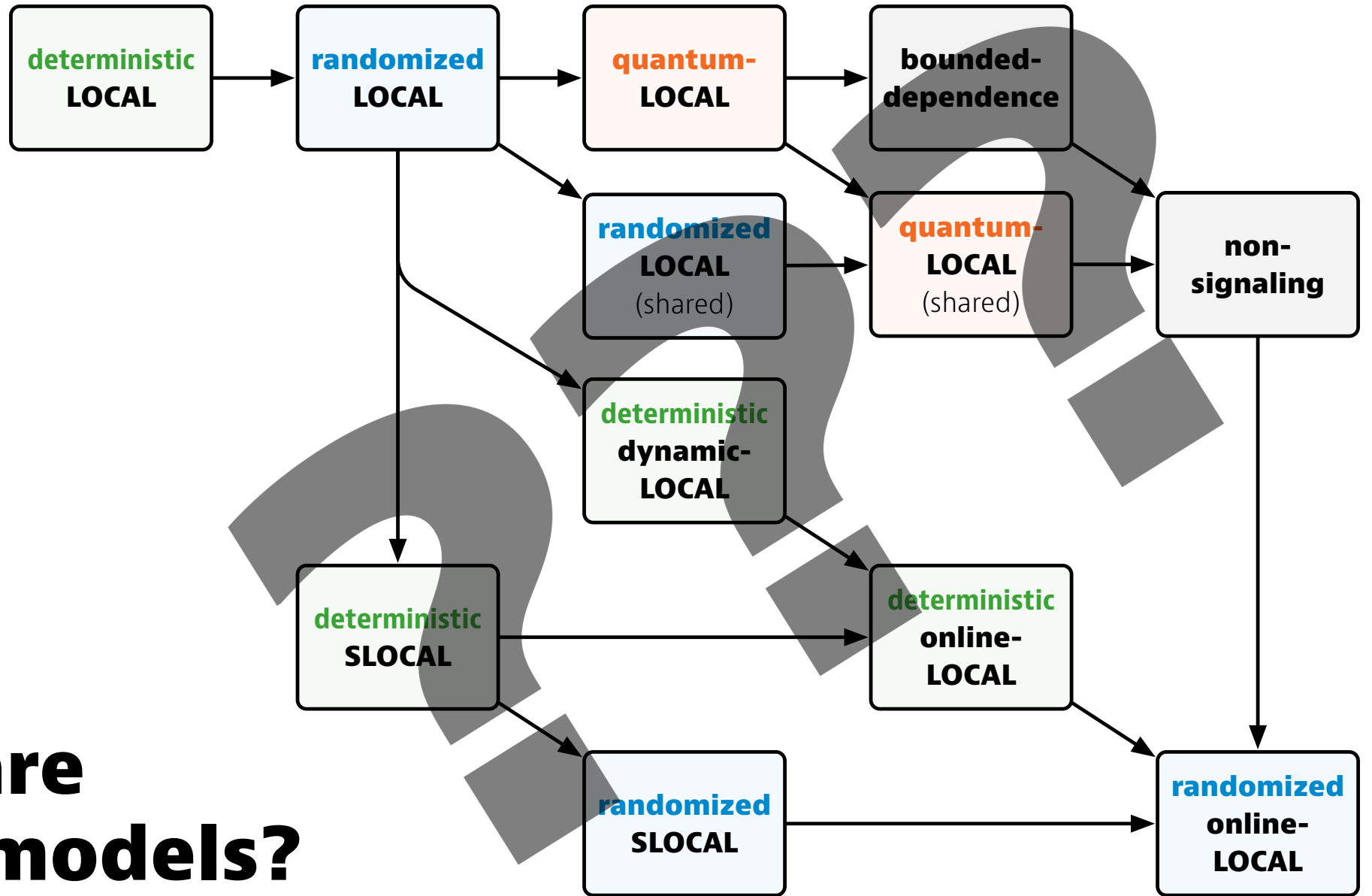
Prove **upper bounds** here...



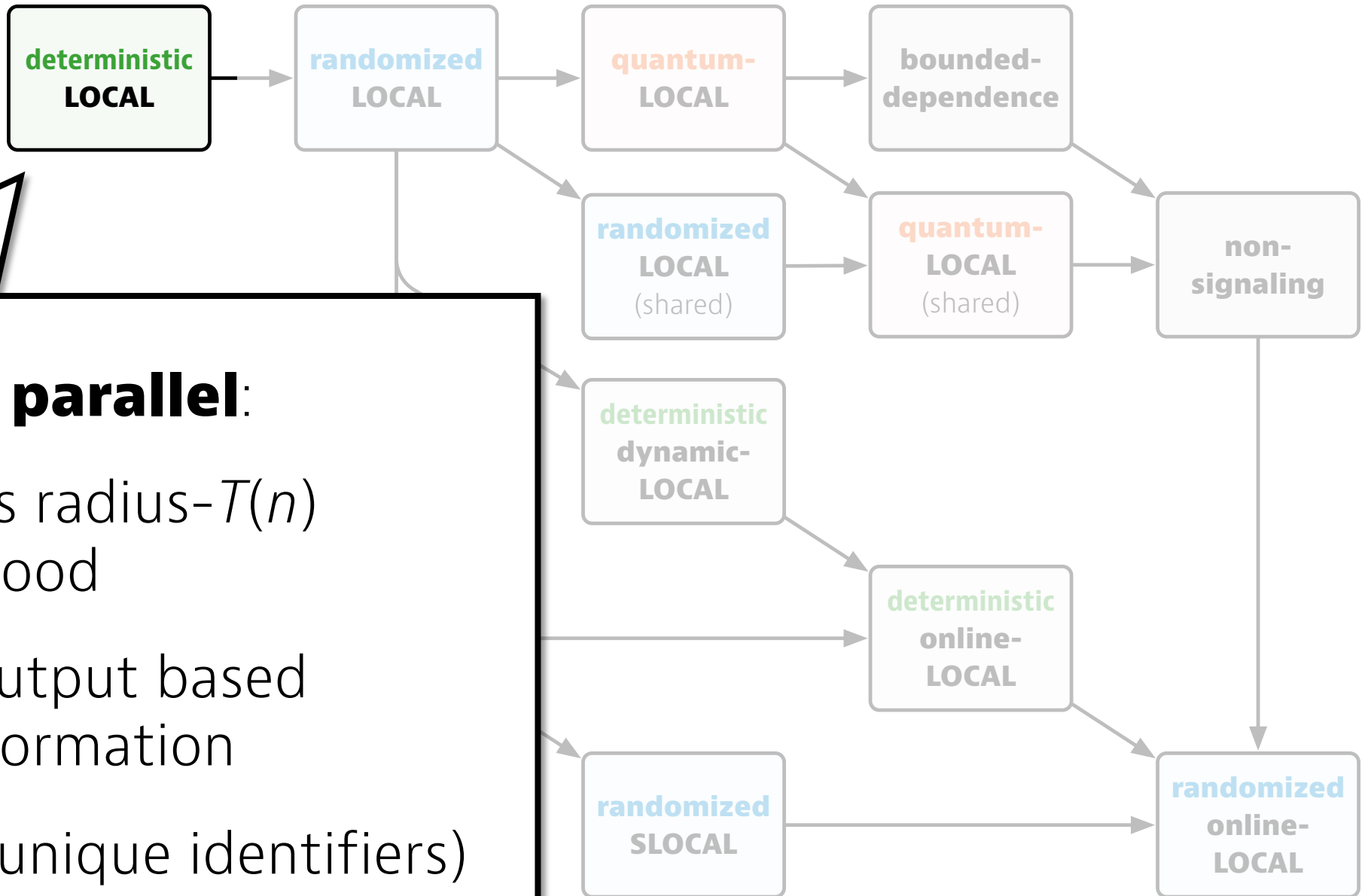
... and **matching lower bounds** here ...

... and for some LCL problems we have tight bounds across all these models!



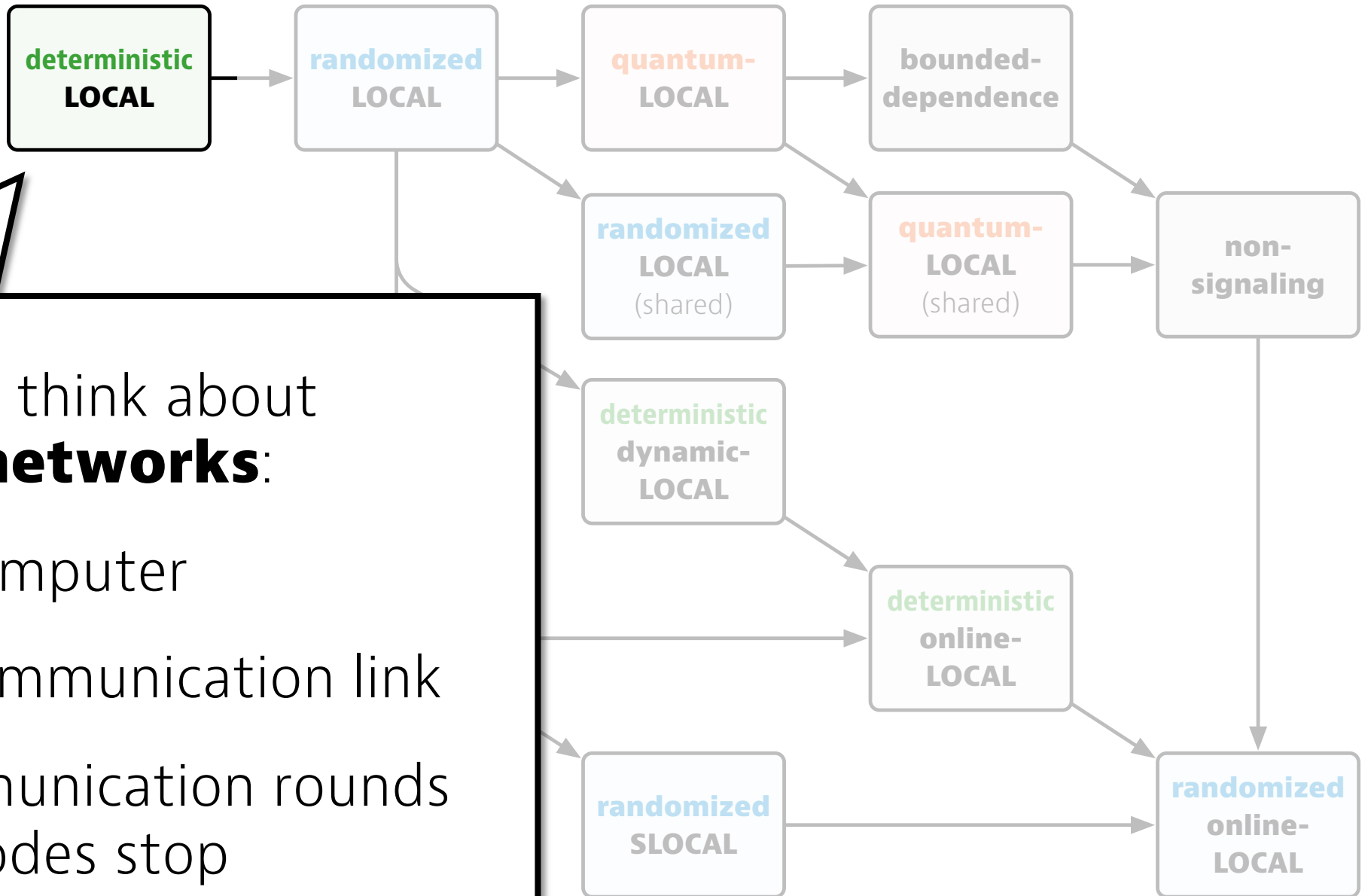


**So what are
all these models?**



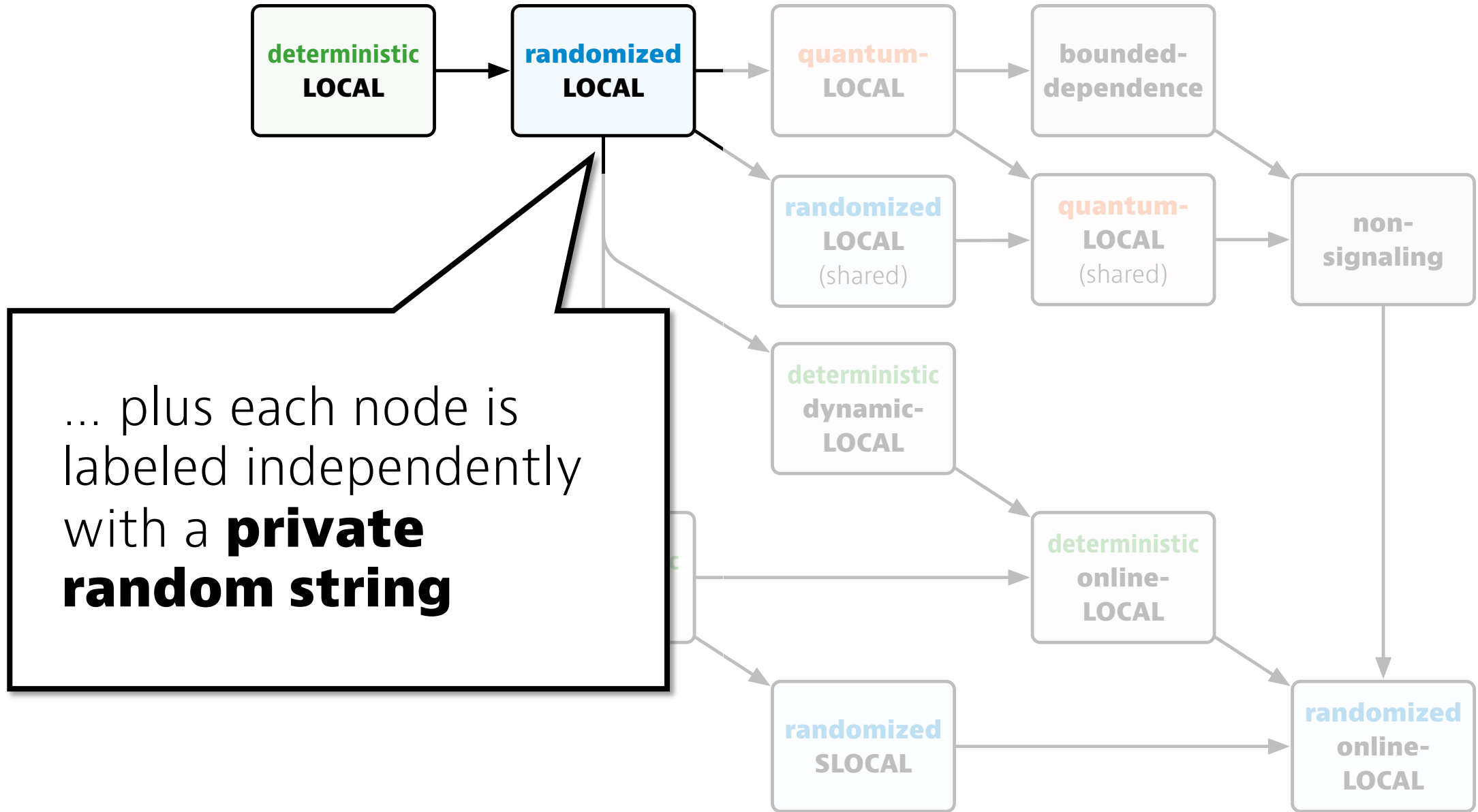
Each node in **parallel**:

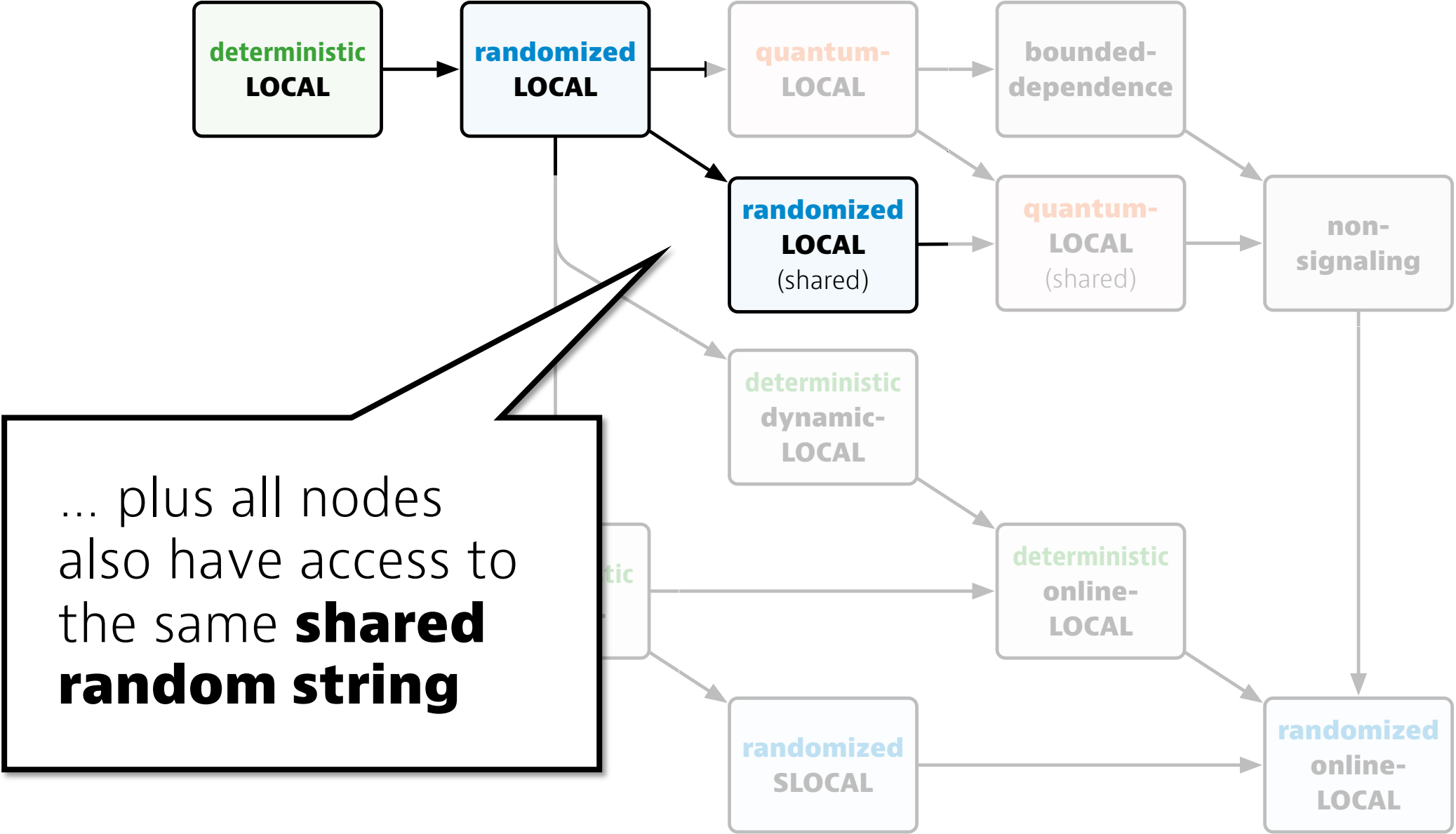
- looks at its radius- $T(n)$ neighborhood
 - picks its output based on this information
- (nodes have unique identifiers)



... or you can think about **computer networks:**

- node = computer
- edge = communication link
- $T(n)$ communication rounds until all nodes stop

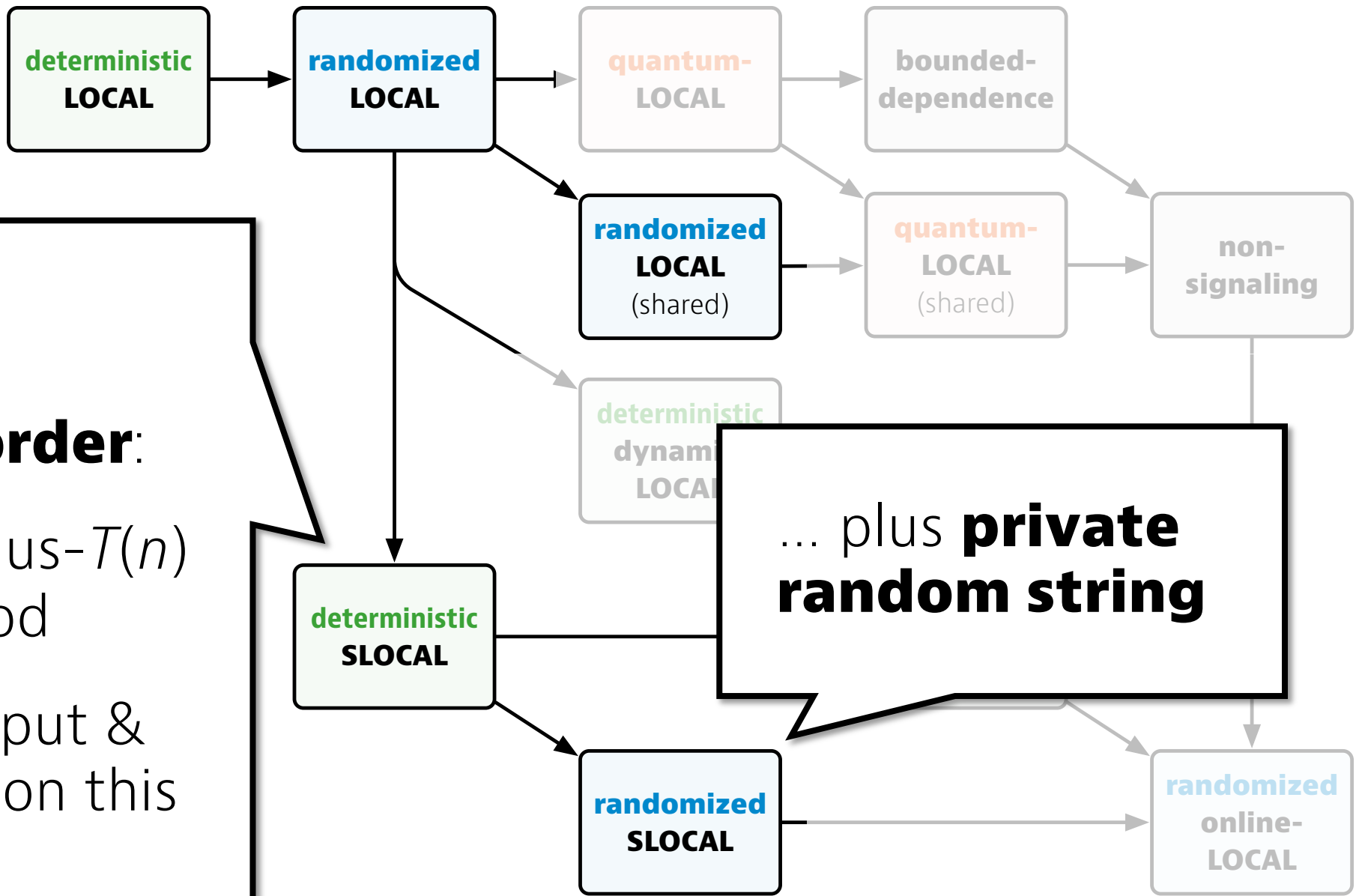




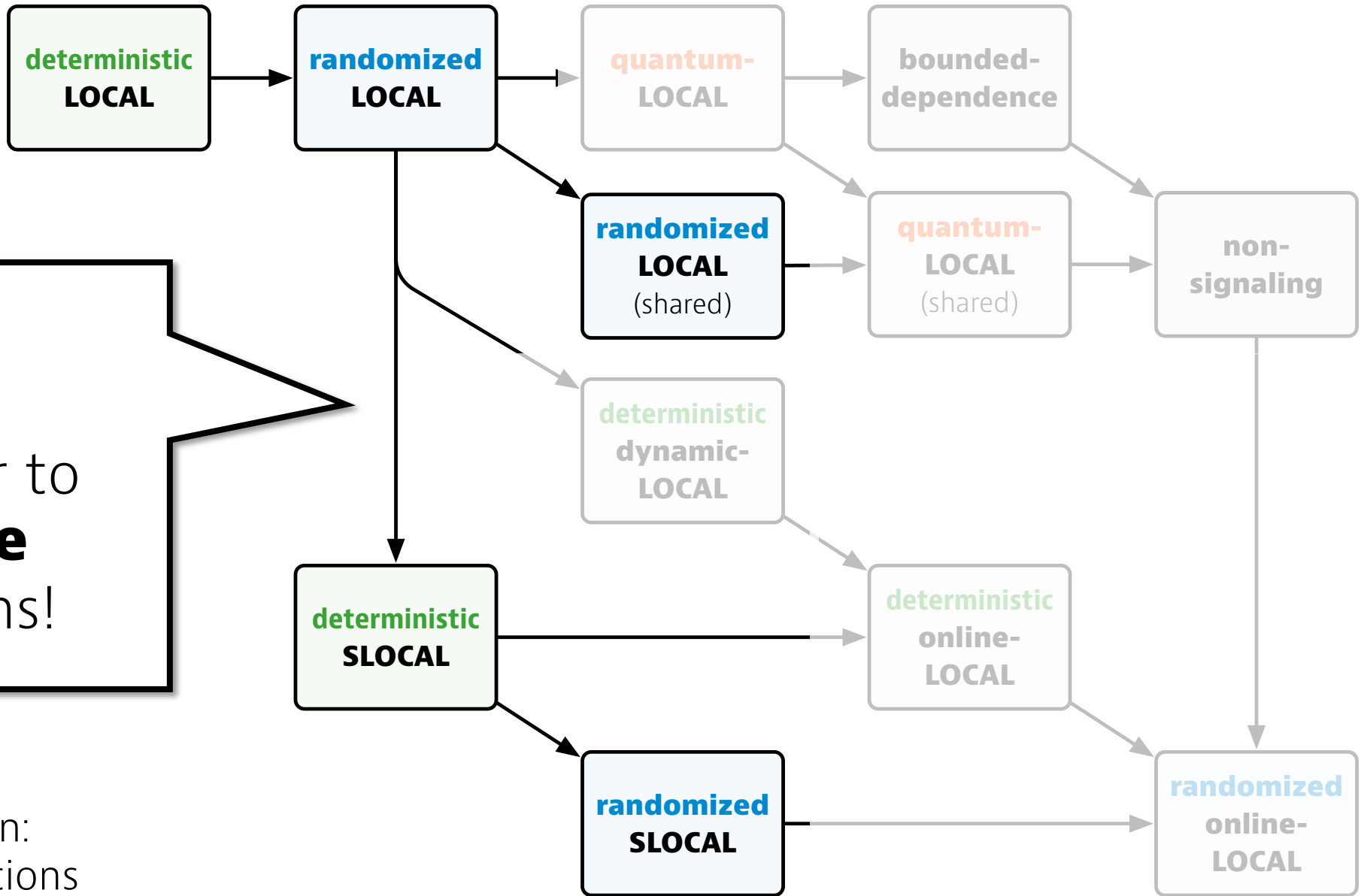
... plus all nodes also have access to the same **shared random string**

Each node in a **sequential, adversarial order**:

- looks at radius- $T(n)$ neighborhood
- picks its output & state based on this information



Sequential ordering gives enough power to **derandomize** local algorithms!



Ghaffari, Harris, Kuhn:
conditional expectations

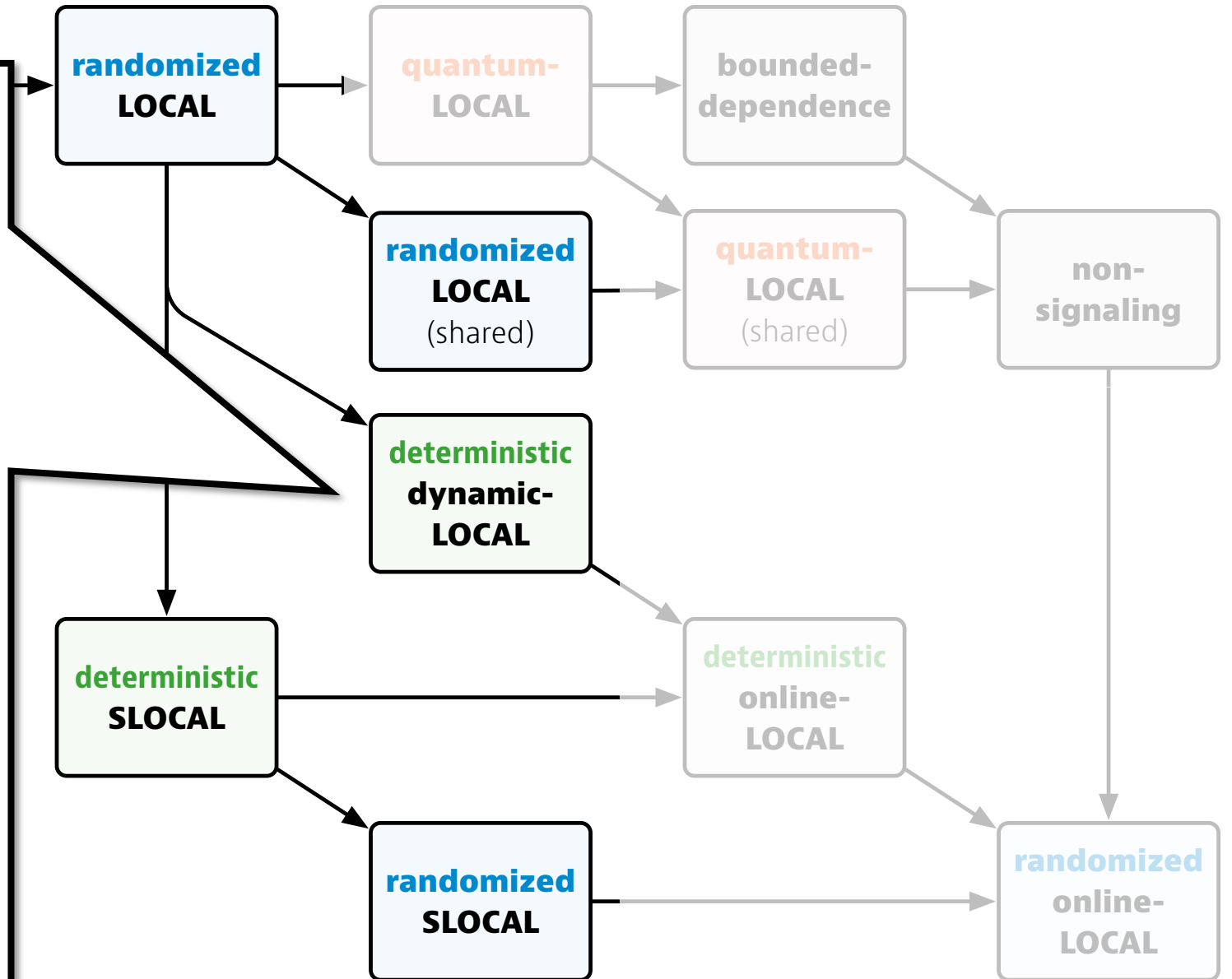
Adversary adds nodes and edges one by one

- We can **see everything**

- Have to maintain valid solution

We can **change** our output only within distance $T(n)$ from a point of change

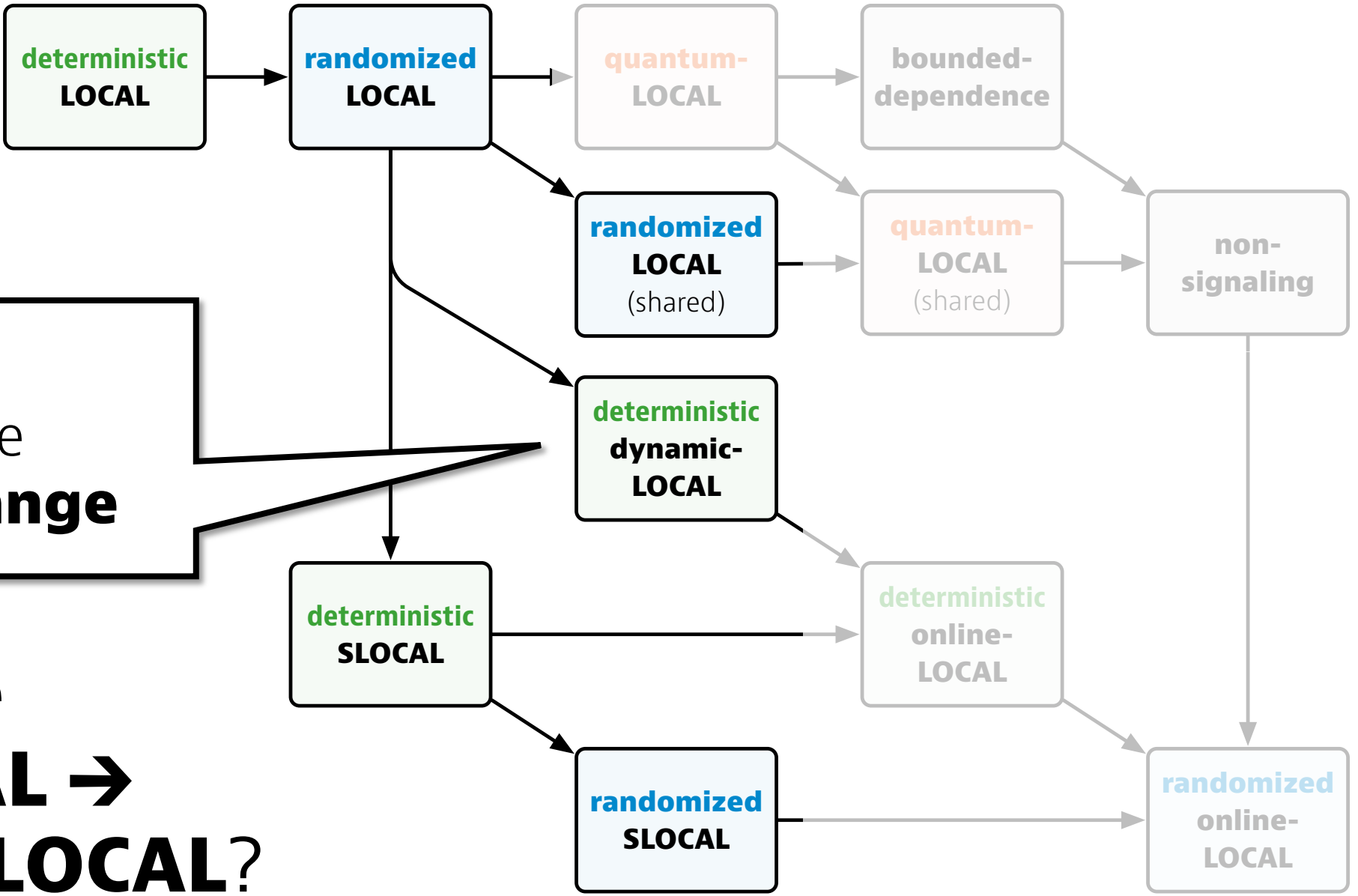
deterministic

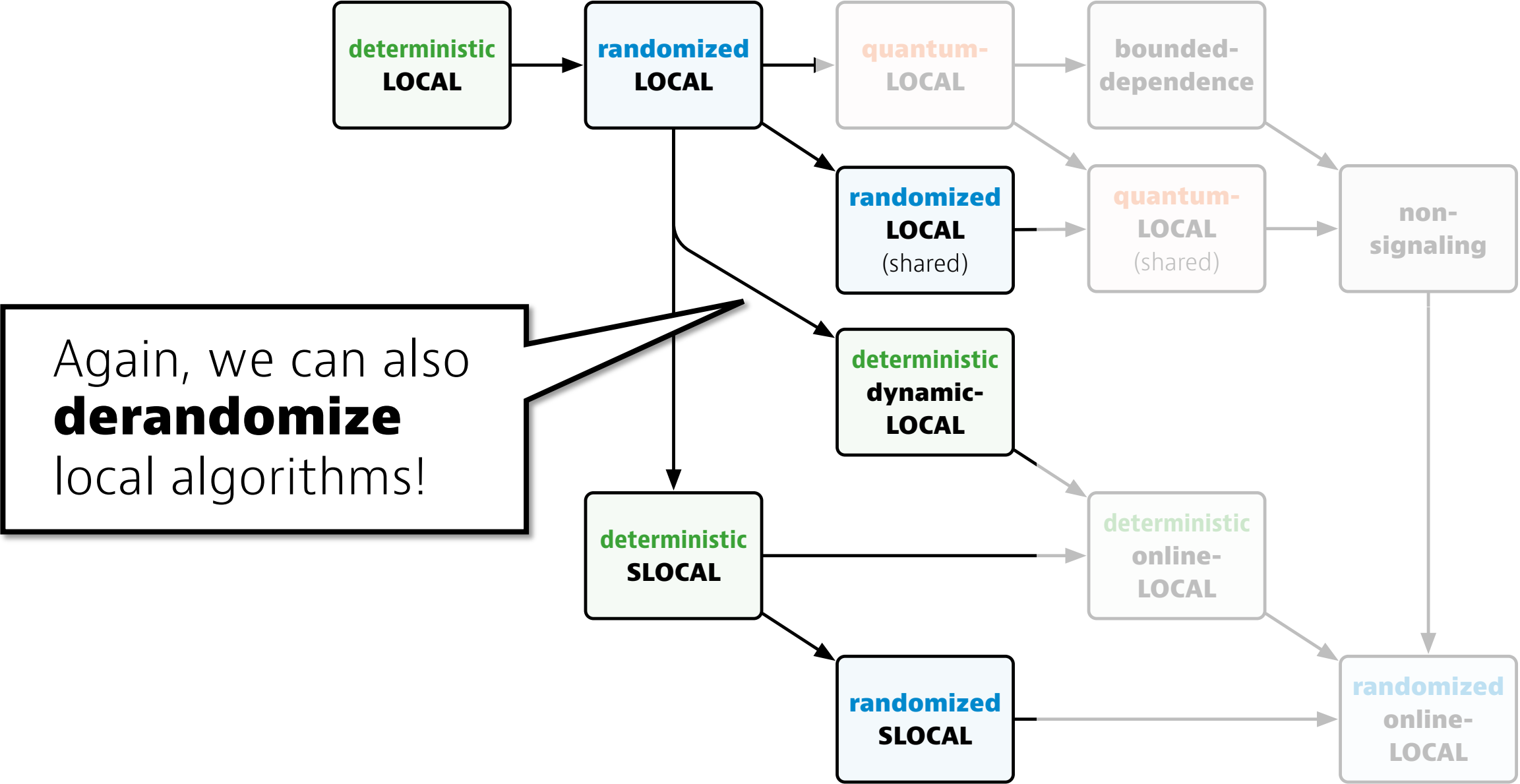


Limited what we can **see**

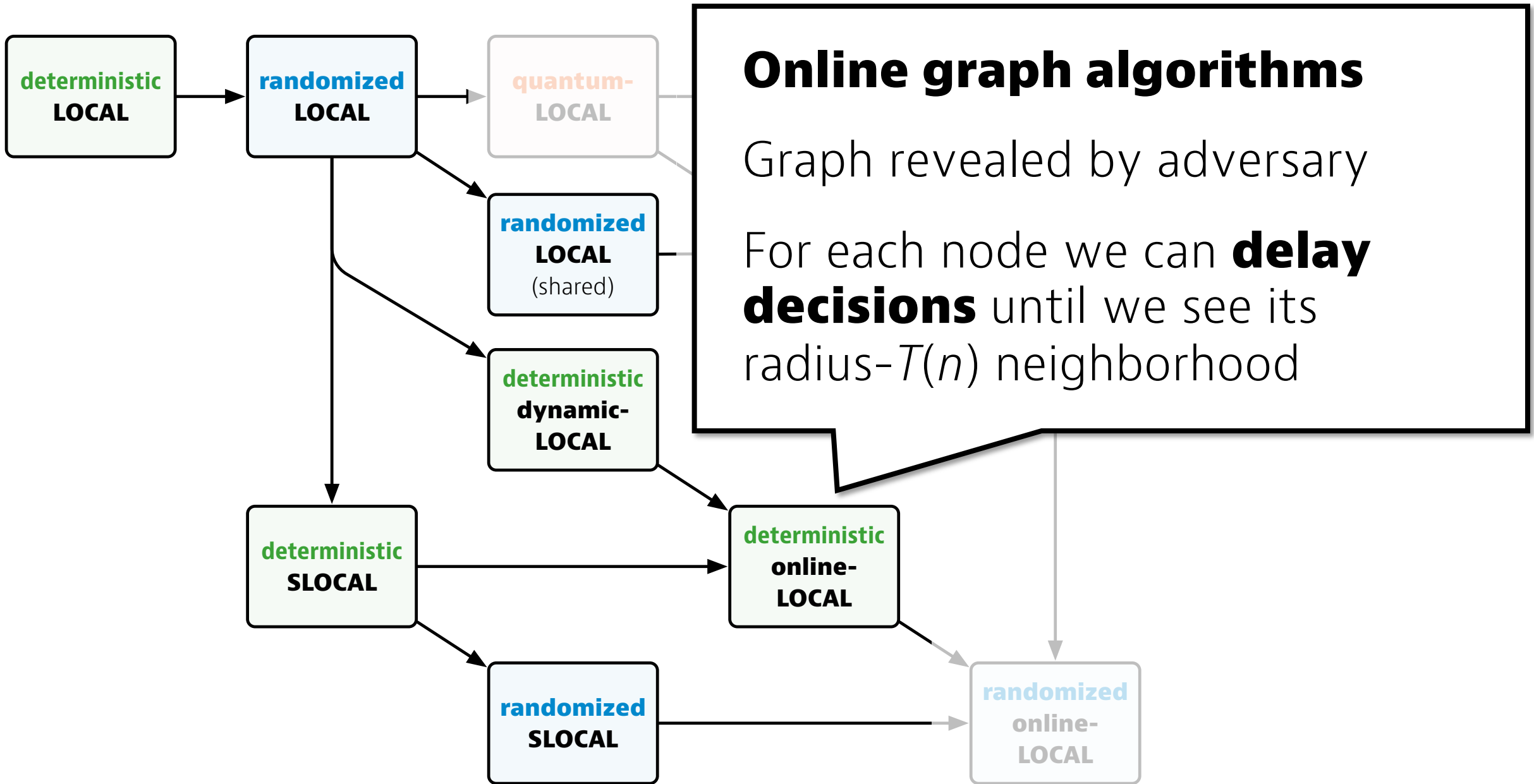
Limited what we can **change**

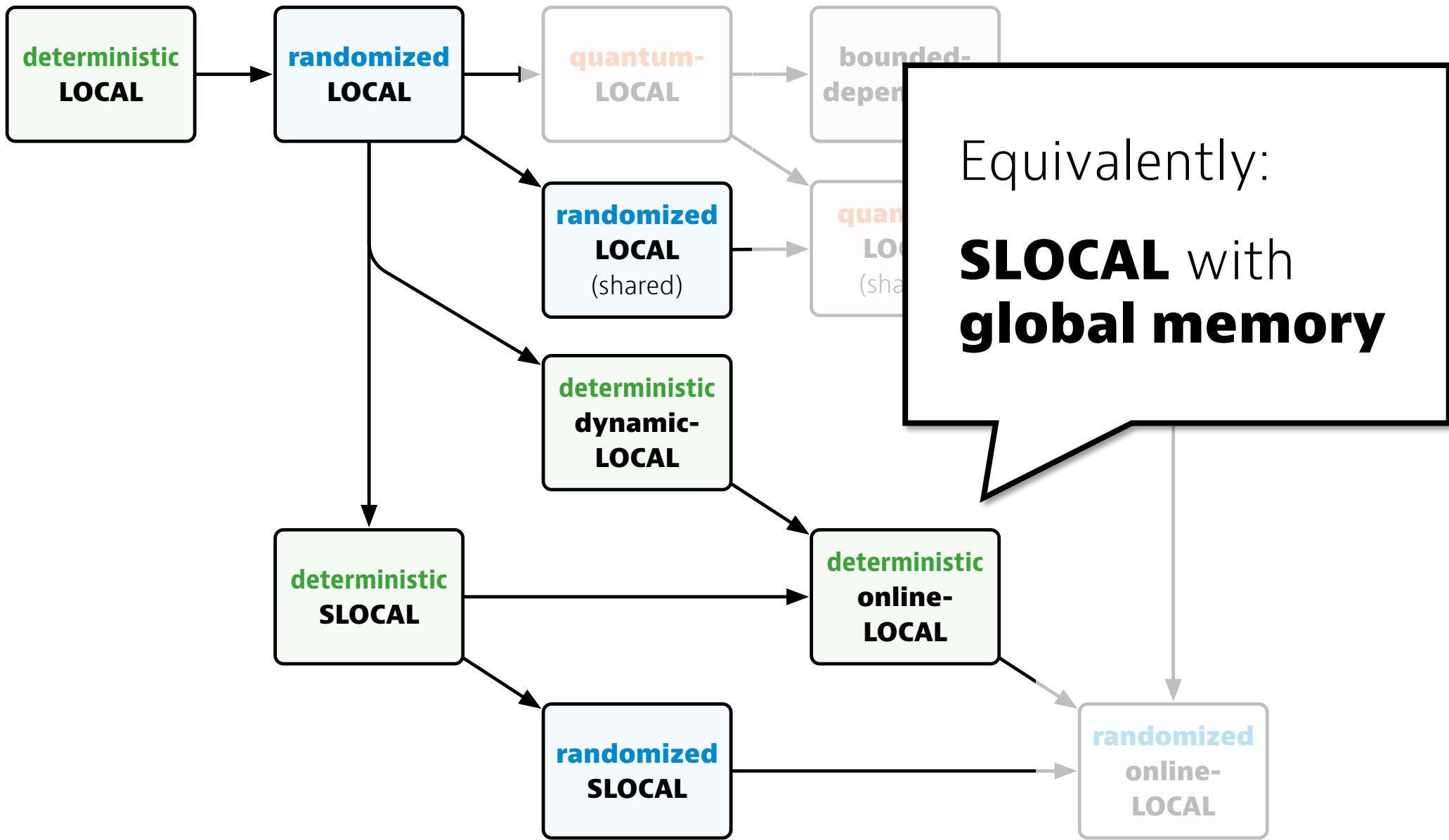
Why do we have **LOCAL** → **dynamic-LOCAL**?



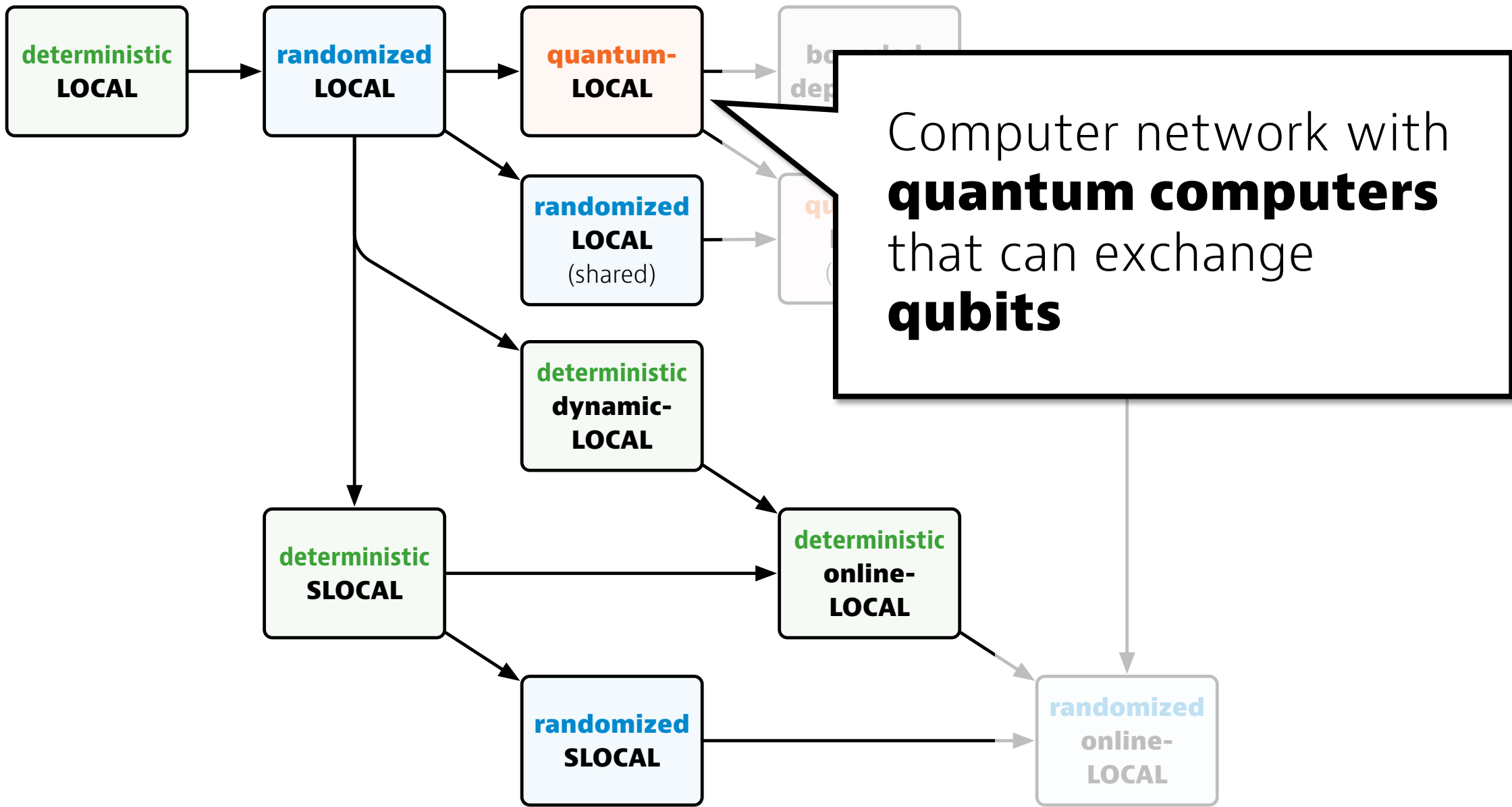


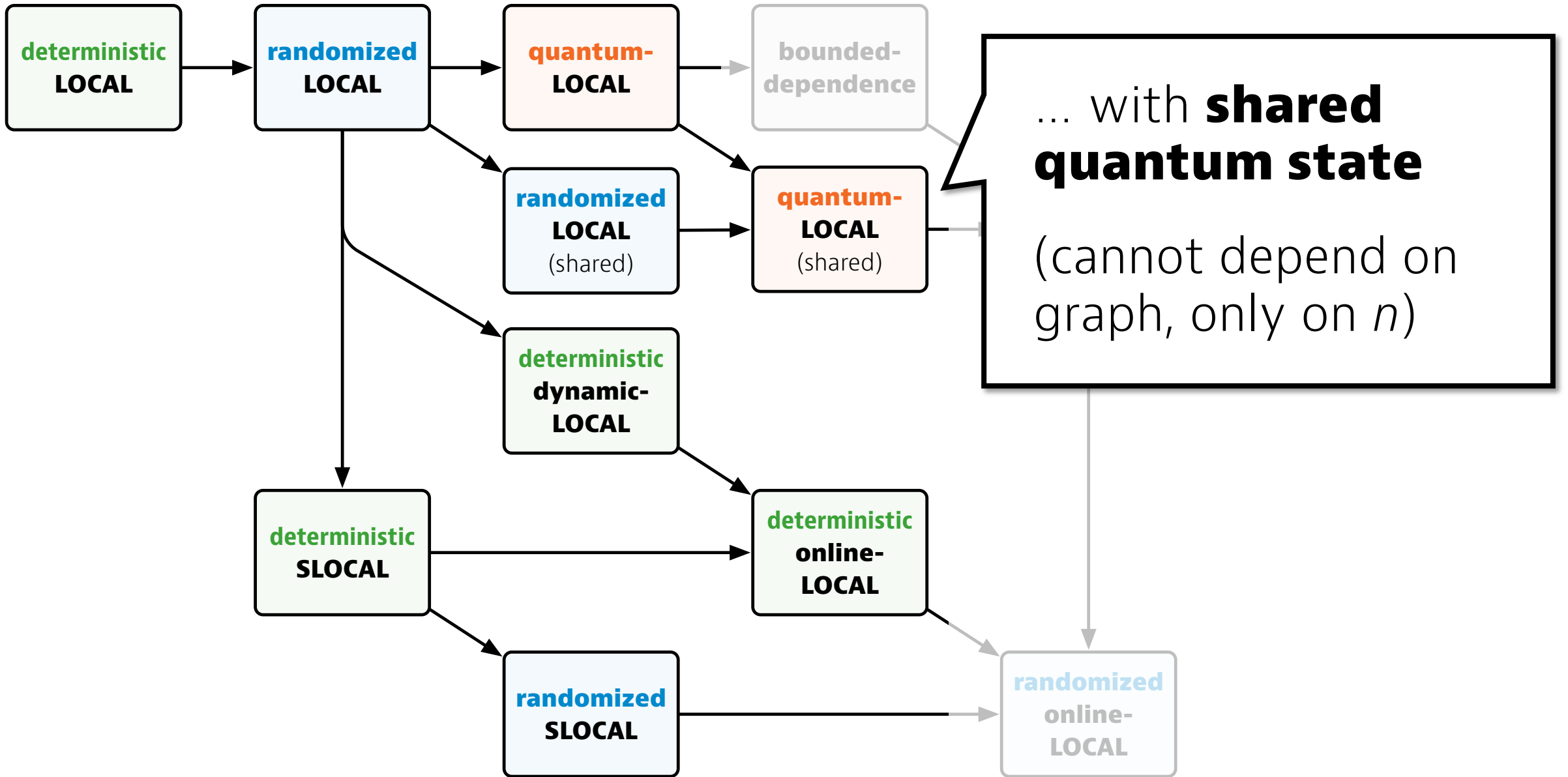
Again, we can also **derandomize** local algorithms!

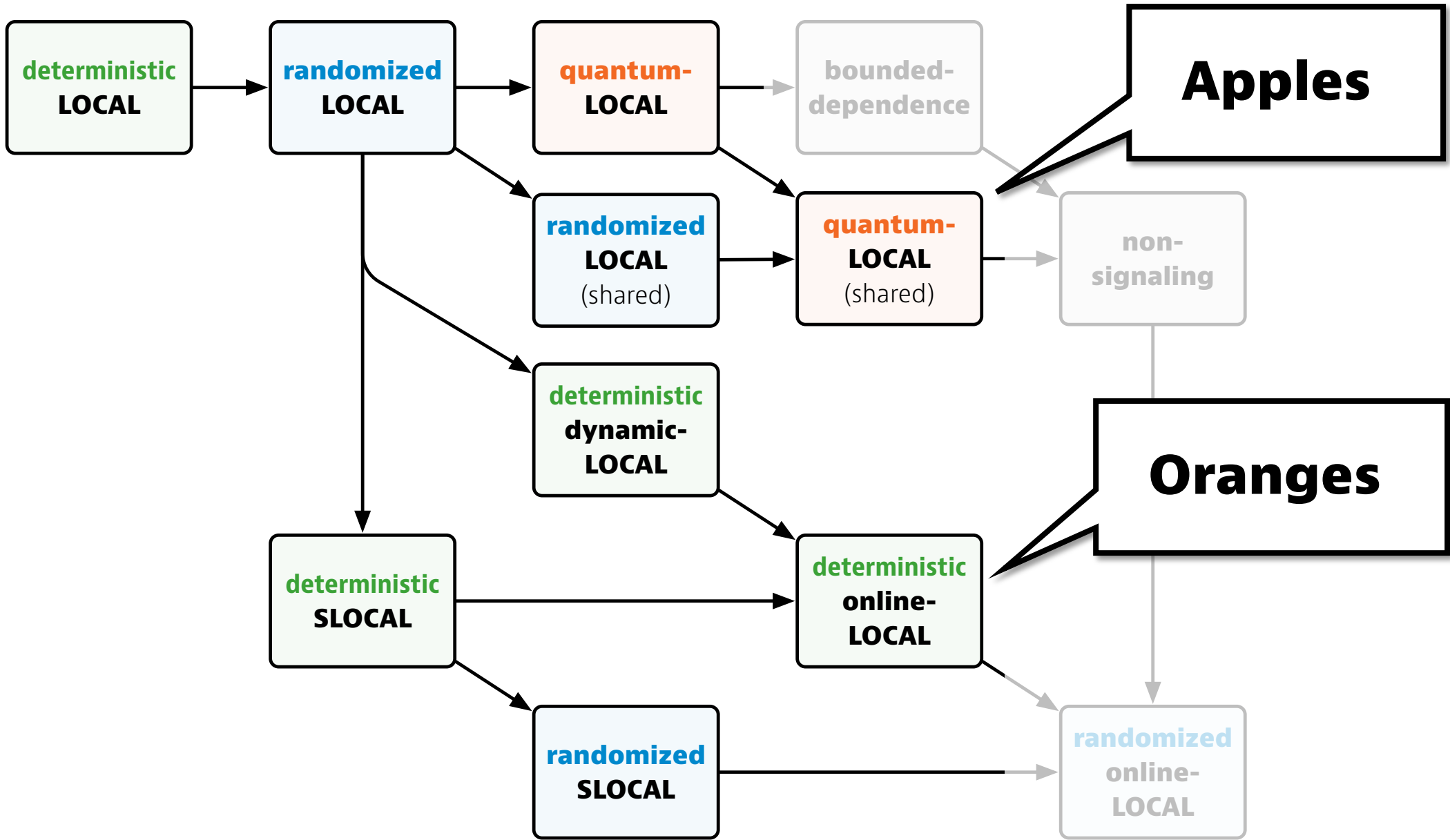


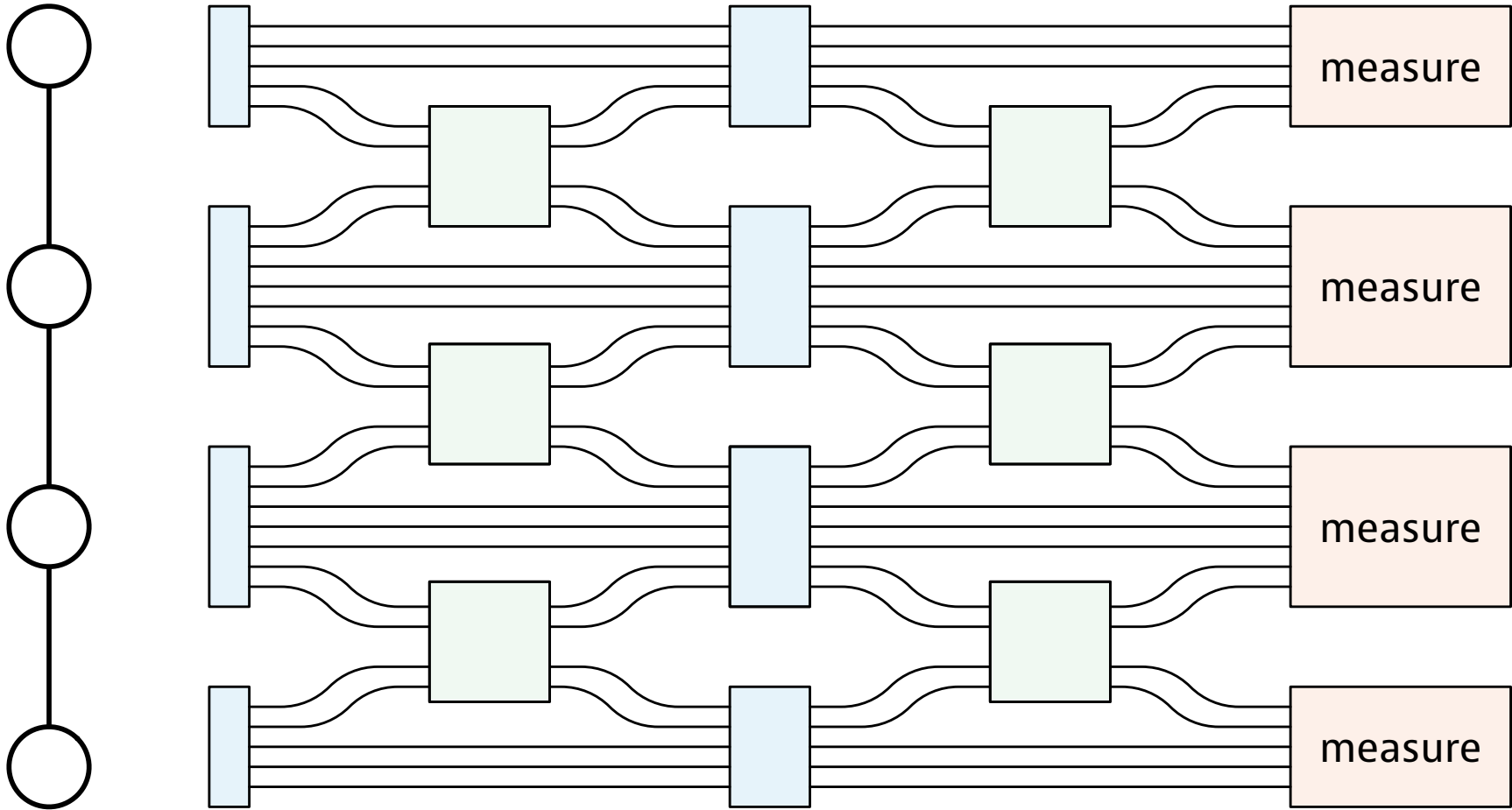


Equivalently:
SLOCAL with
global memory







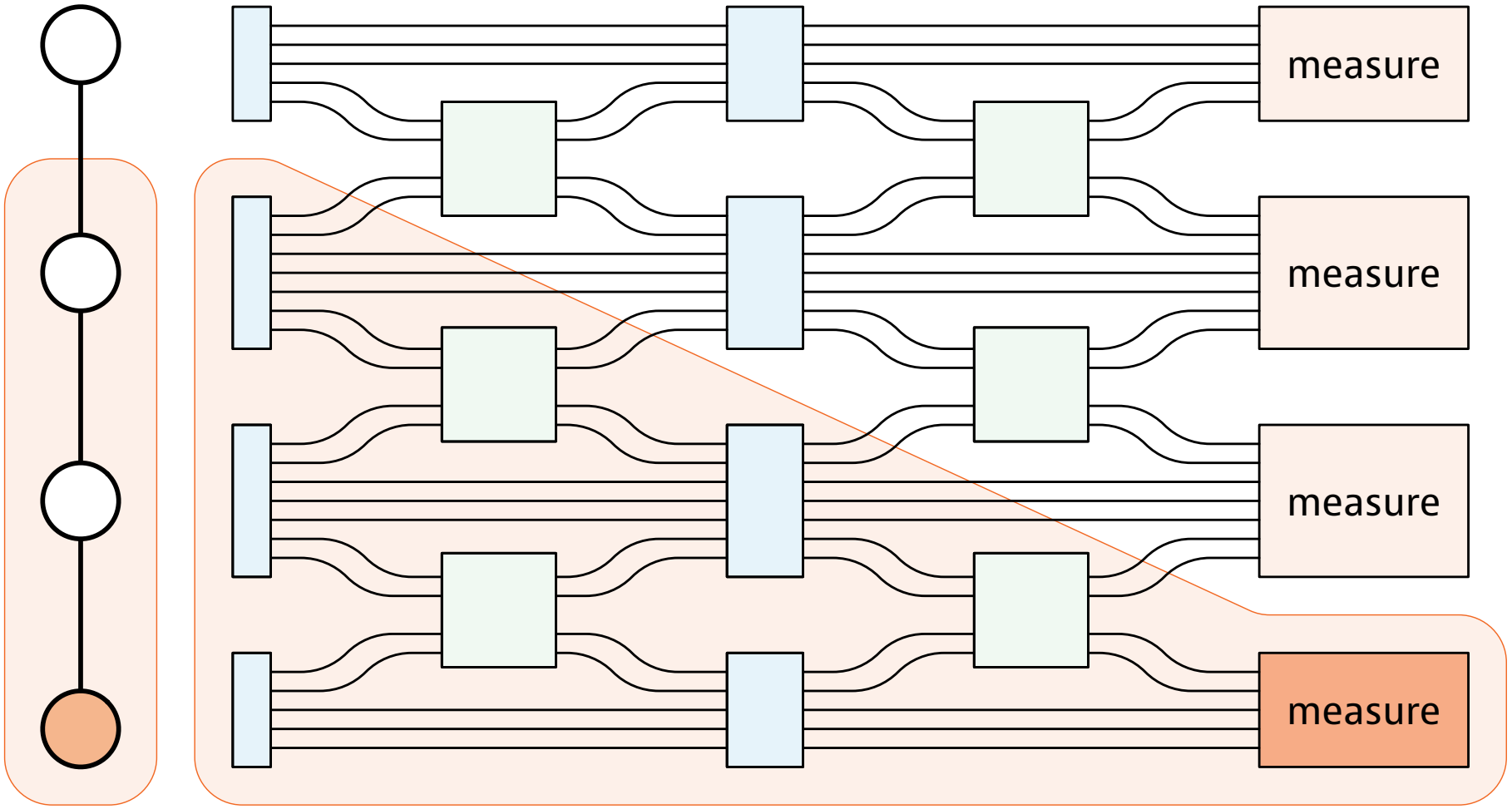


2 rounds

communication

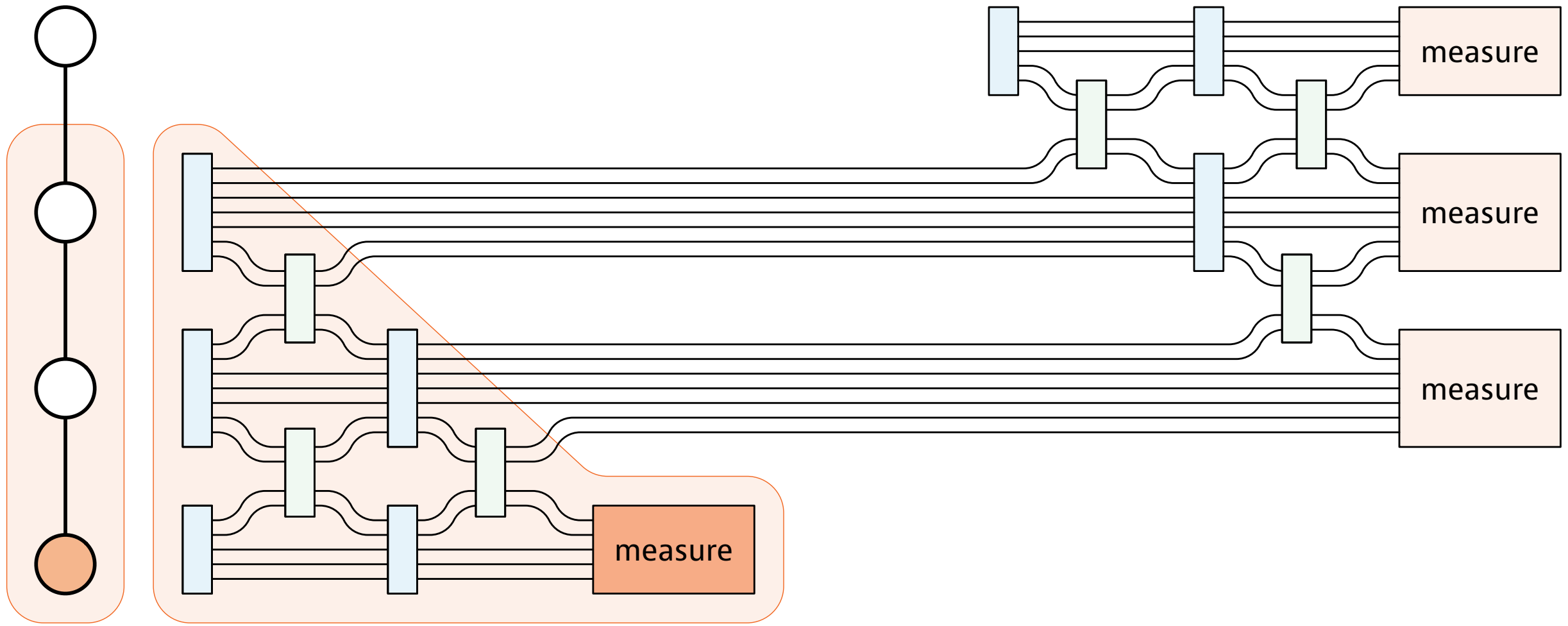
*local
computation*

communication



light cone

2 rounds



2 rounds

light cone

Non-signaling model

- Quantum LOCAL can't violate causality
- Key idea: **define** a model so that it can do **anything** except violating causality

Non-signaling model

Definition (*non-signaling distribution*):

- fix any **set of nodes X** ...

Gavoille, Kosowski, Markiewicz 2009
Arfaoui, Fraigniaud 2014



Non-signaling model

Definition (*non-signaling distribution*):

- fix any **set of nodes X**
- changes in the input **more than T hops away** from **X** do not influence the output distribution of **X**

Gavoille, Kosowski, Markiewicz 2009
Arfaoui, Fraigniaud 2014



**Classical
probability
theory**

Classical (randomized) distributed algorithms



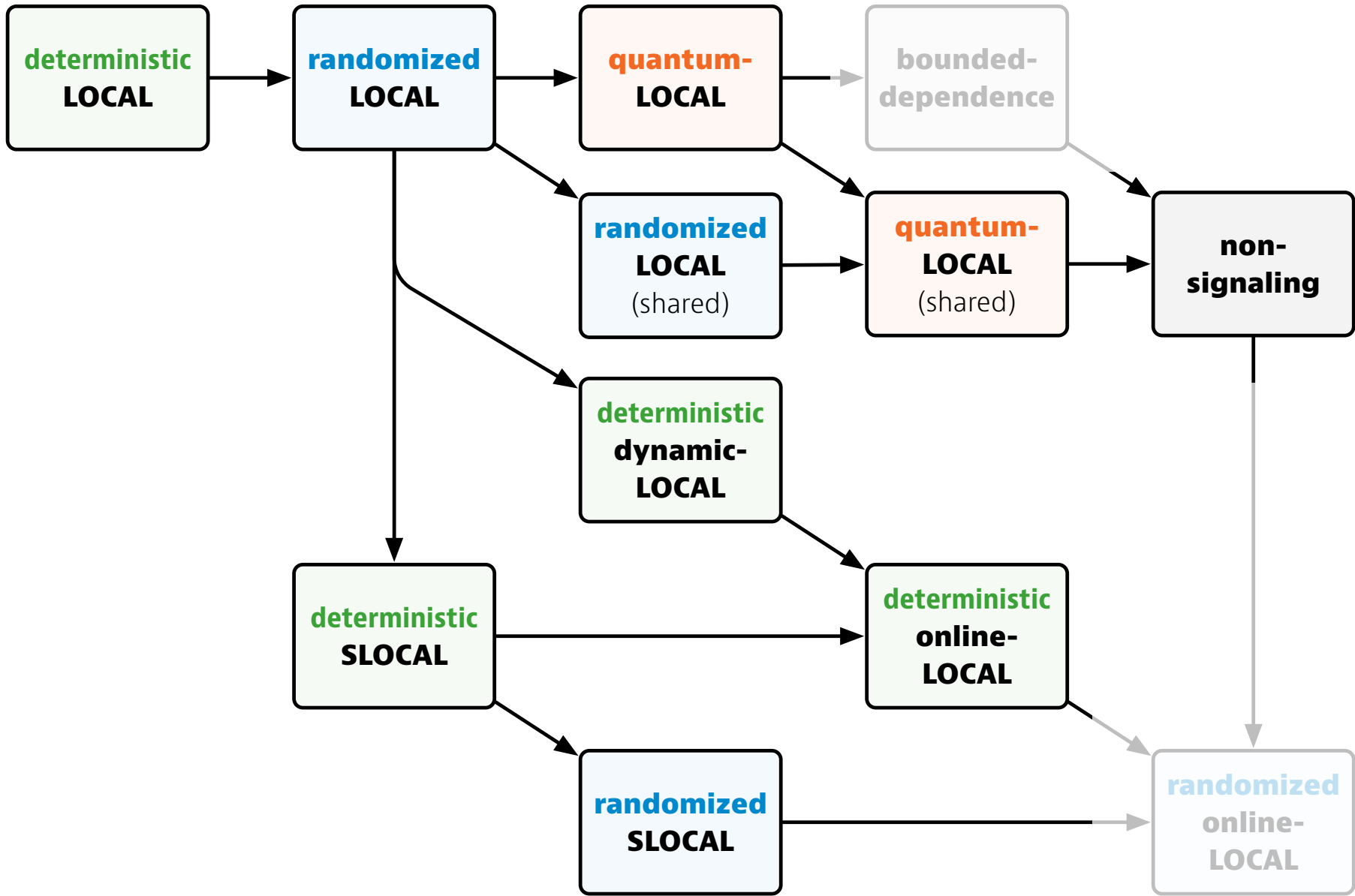
Quantum distributed algorithms

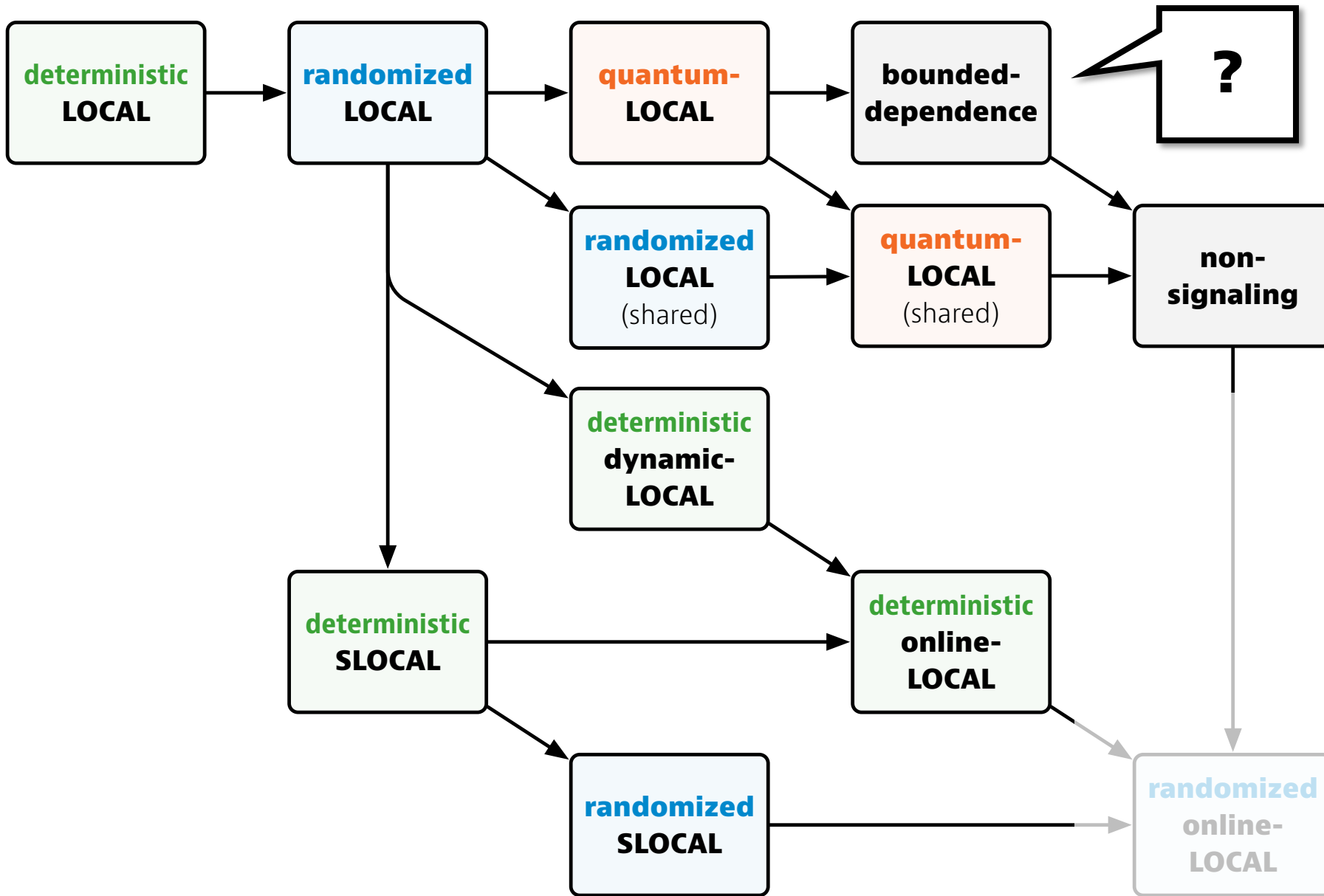
**Weird
quantum
things**



Non-signaling "algorithms"

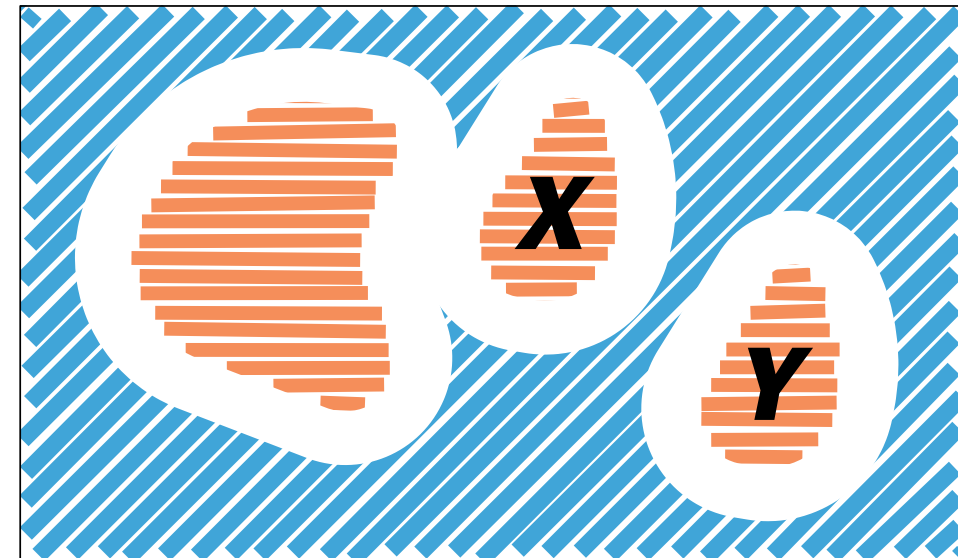
**Classical
probability
theory**





Bounded dependence

- **Finely dependent distributions**
 - X and Y far from each others \rightarrow independent
 - usually "far" = some constant
- For clarity, we call it here **bounded dependence model** when "far" = $T(n)$



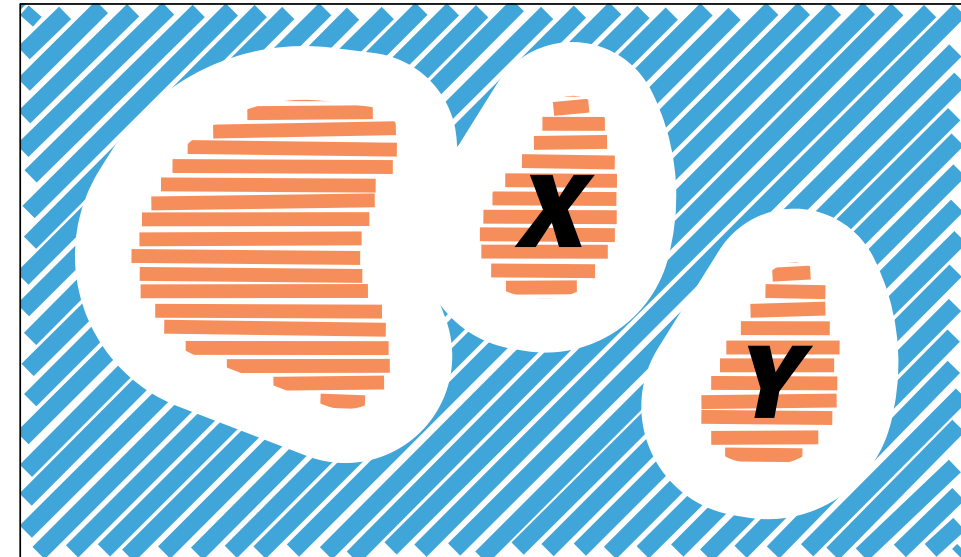
Bounded dependence

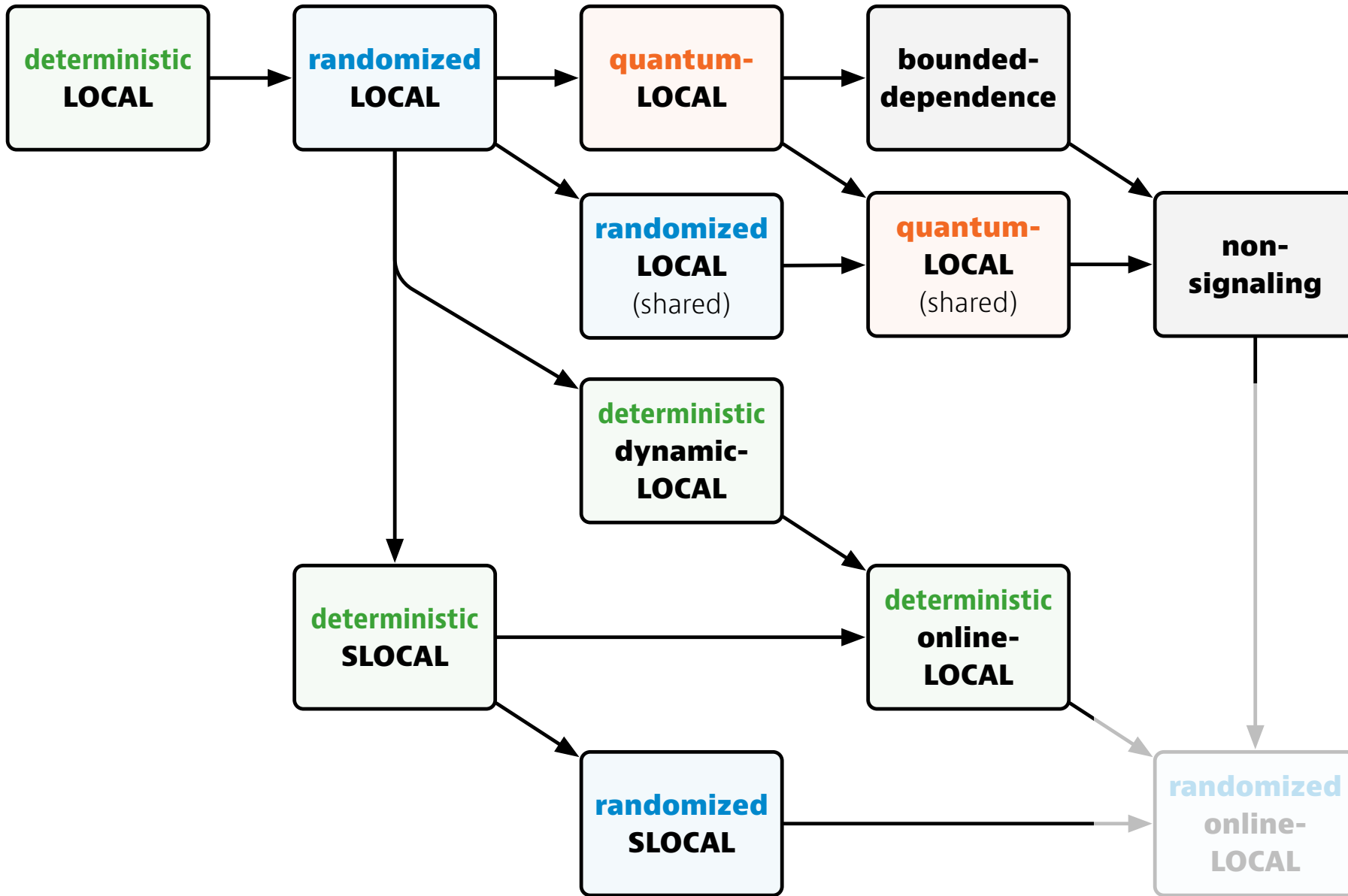
Quantum-LOCAL without shared quantum state:

Output of $T(n)$ -round algorithm



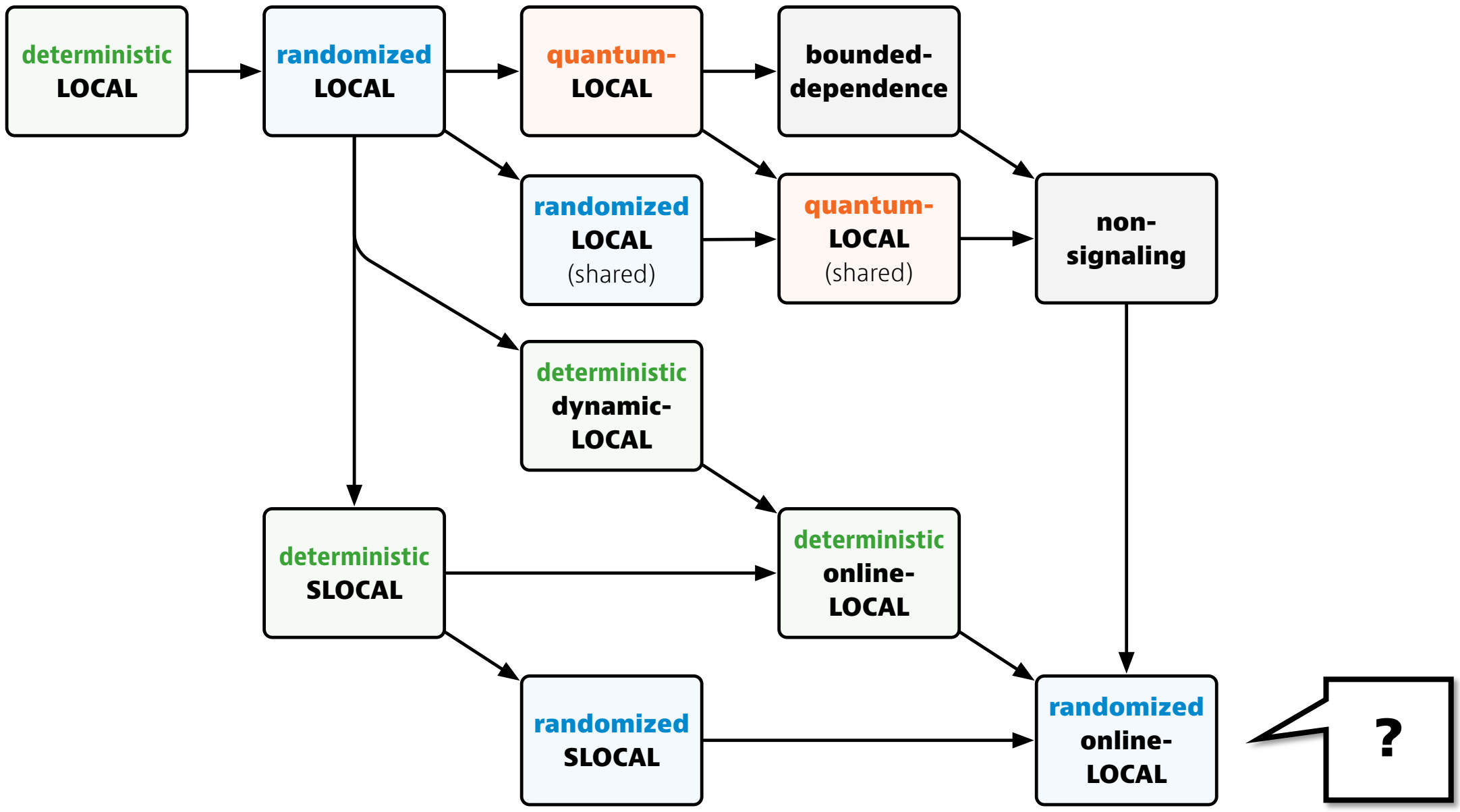
Not just non-signaling but also bounded dependence for distance $\approx T(n)$





Constant-round quantum-LOCAL algorithms output finitely dependent distributions!

Useful tool for creating interesting finitely dependent distributions?



Rand. online-LOCAL

- Adversary fixes a **graph** + **order** in which nodes are revealed
- For each node v
 - algorithm sees radius- T neighborhood of v
 - algorithm must choose the label of v
- Algorithm can **remember** everything, algorithm can use **randomness**

Oblivious adversary

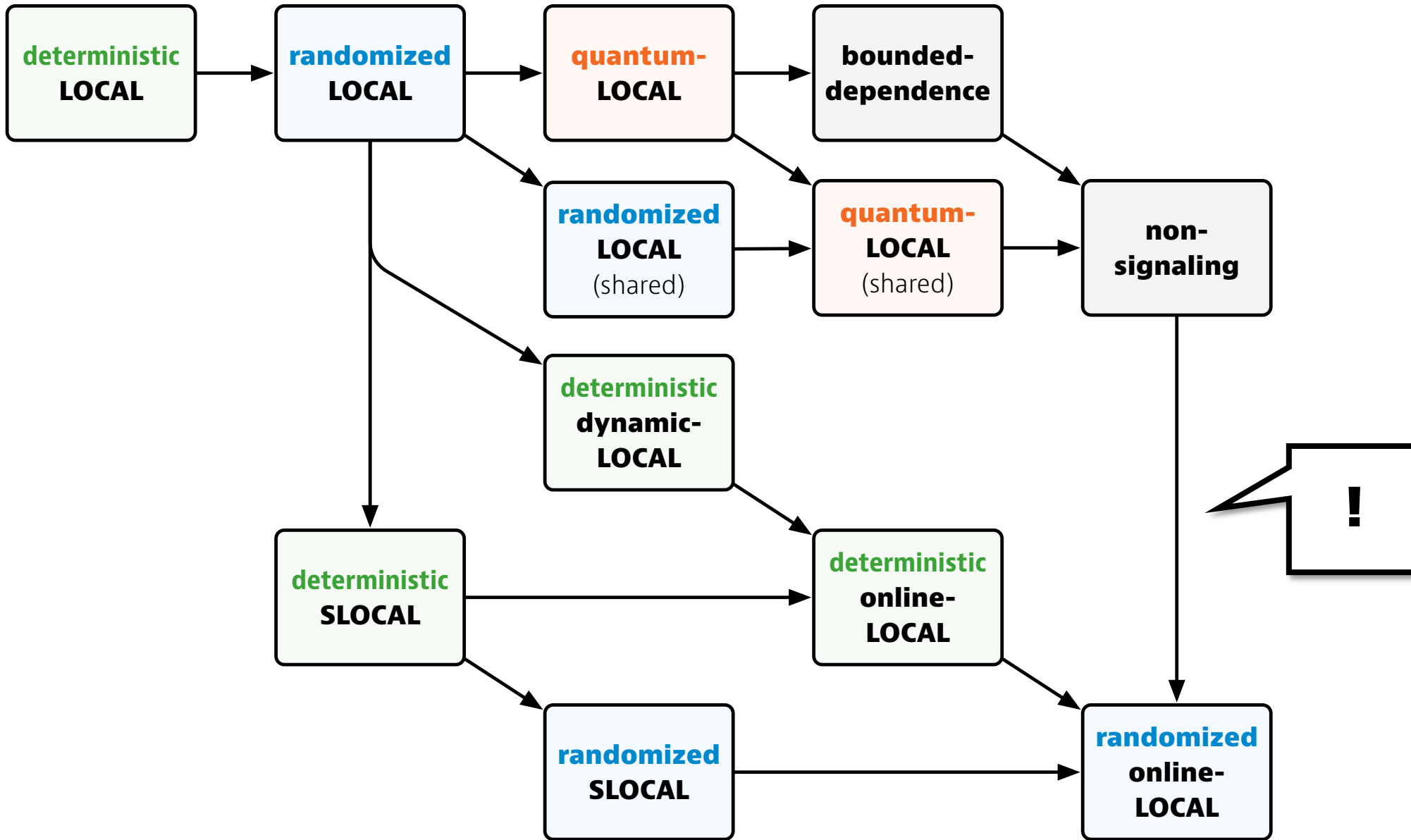
Rand. online-LOCAL

- **Trivial:**

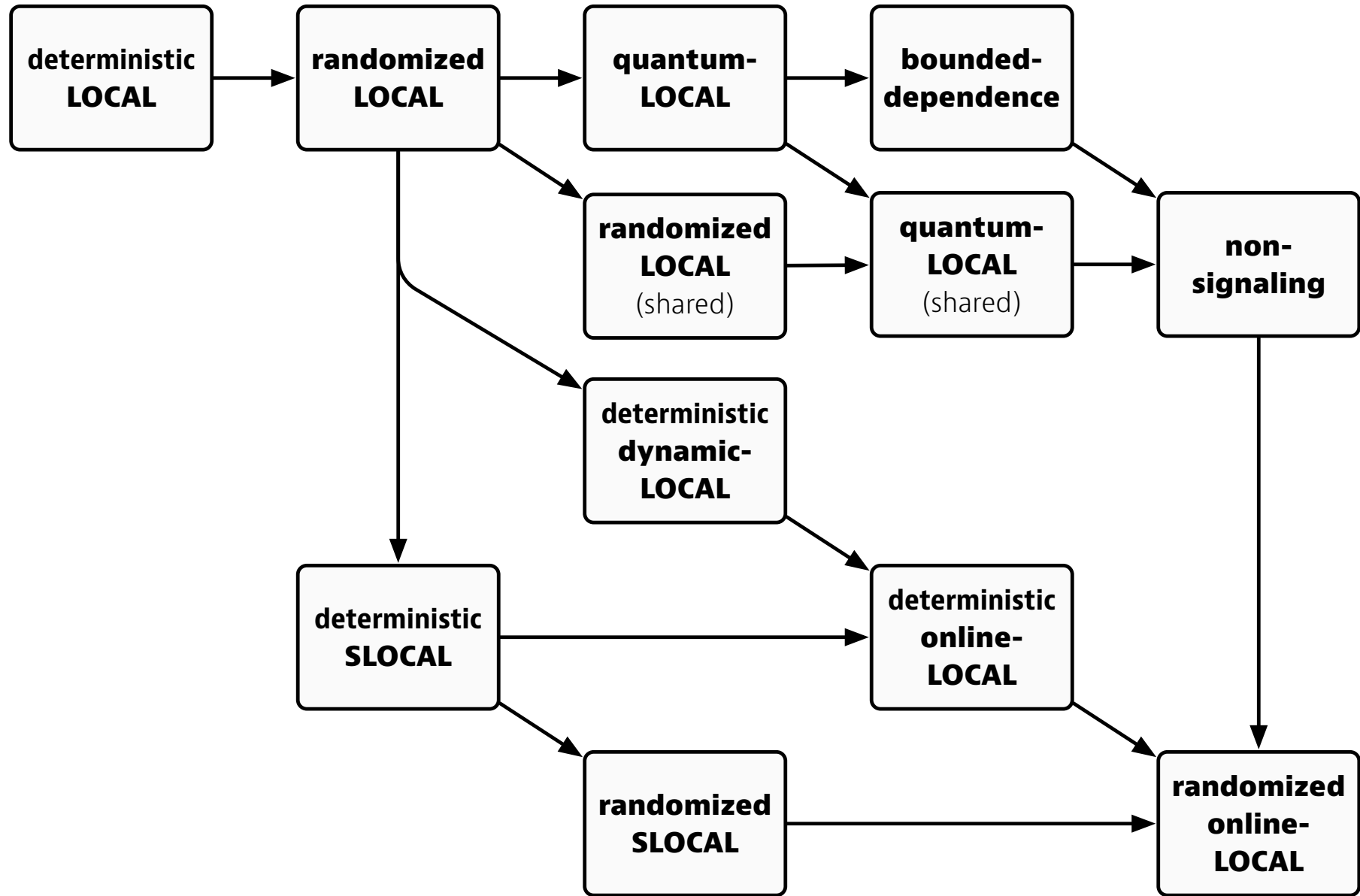
- randomize online-LOCAL can simulate deterministic online-LOCAL

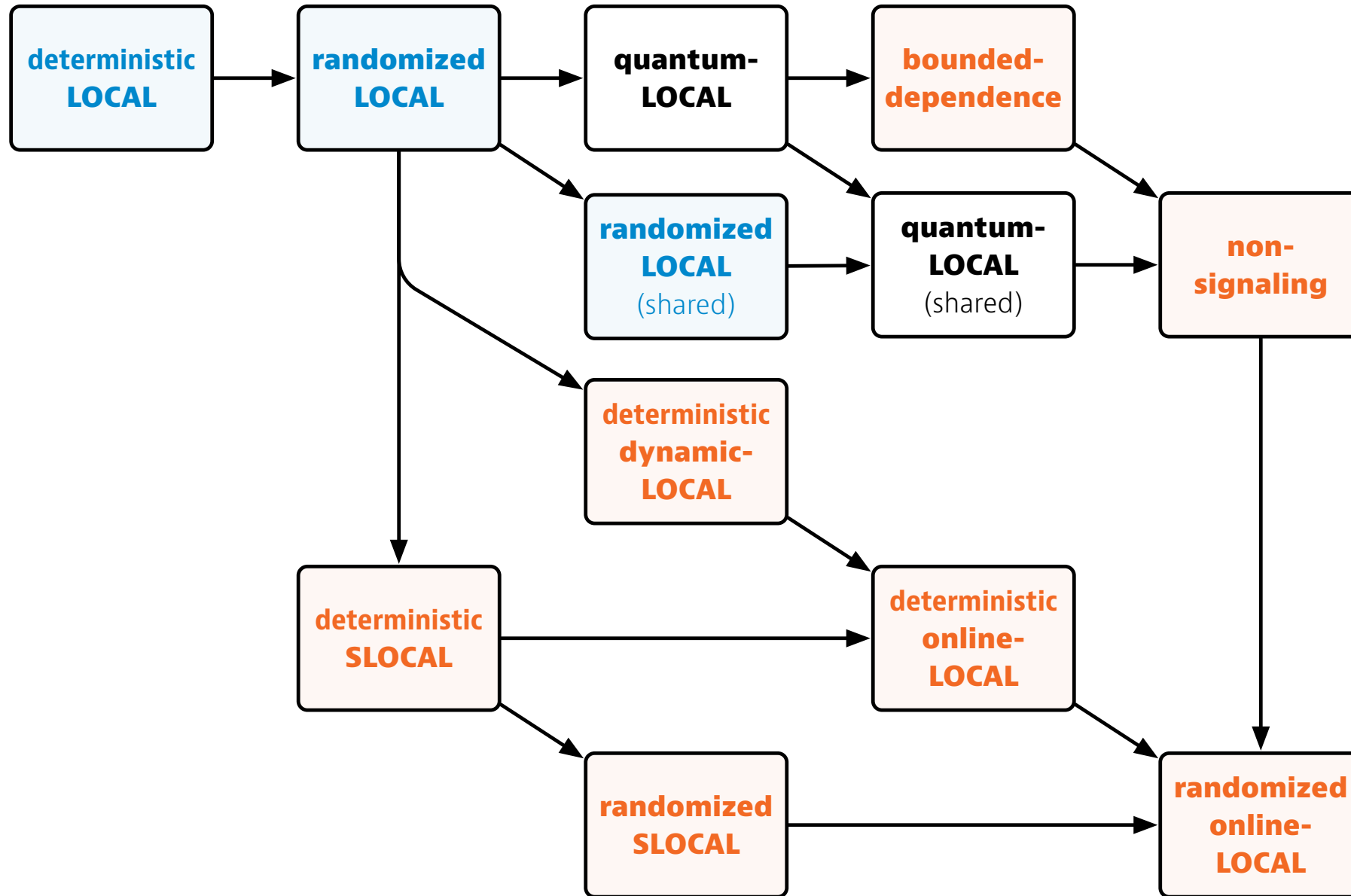
- **Suprise:**

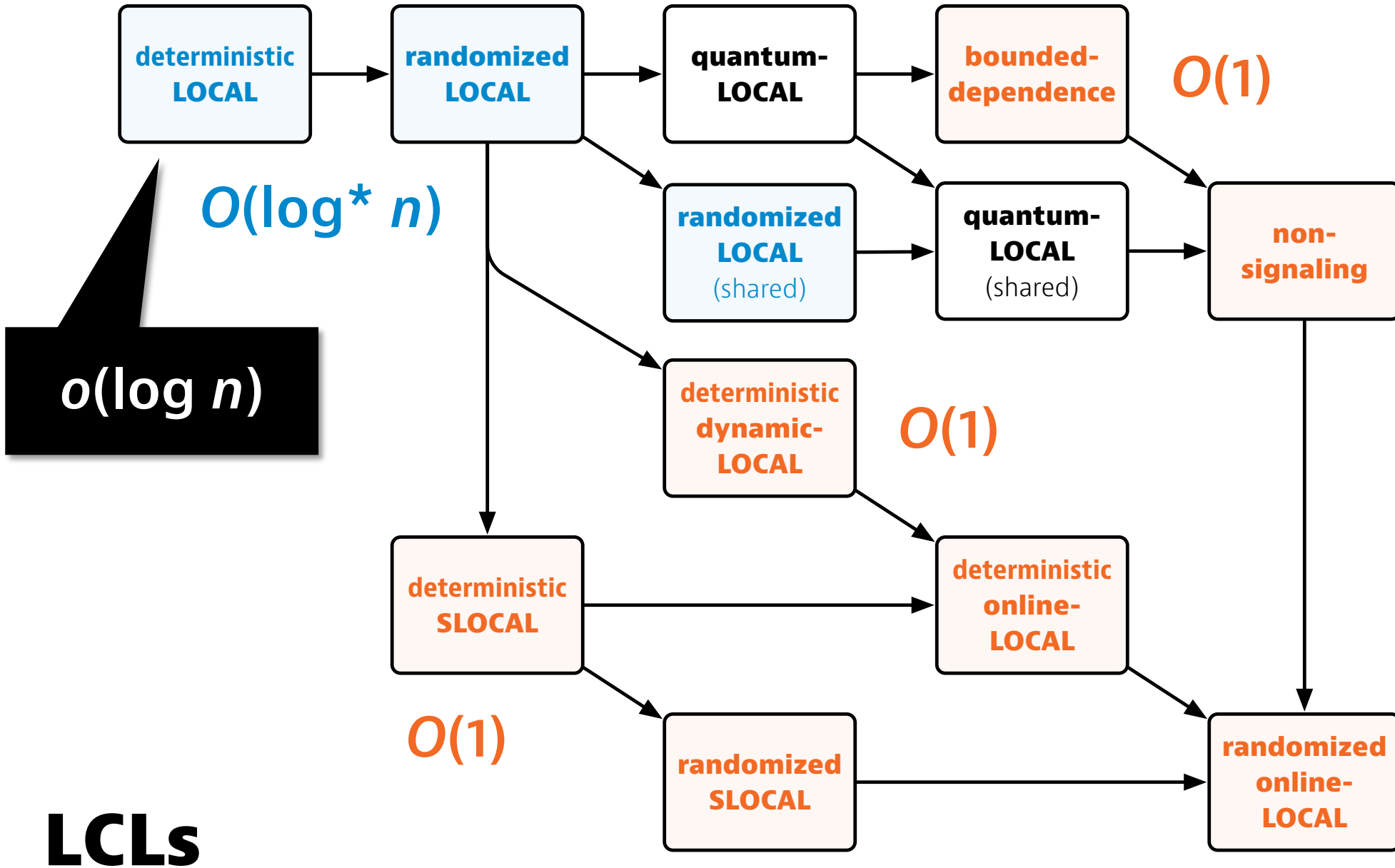
- randomized online-LOCAL *can simulate any non-signaling distribution* (with the same asymptotic locality)



Some results...

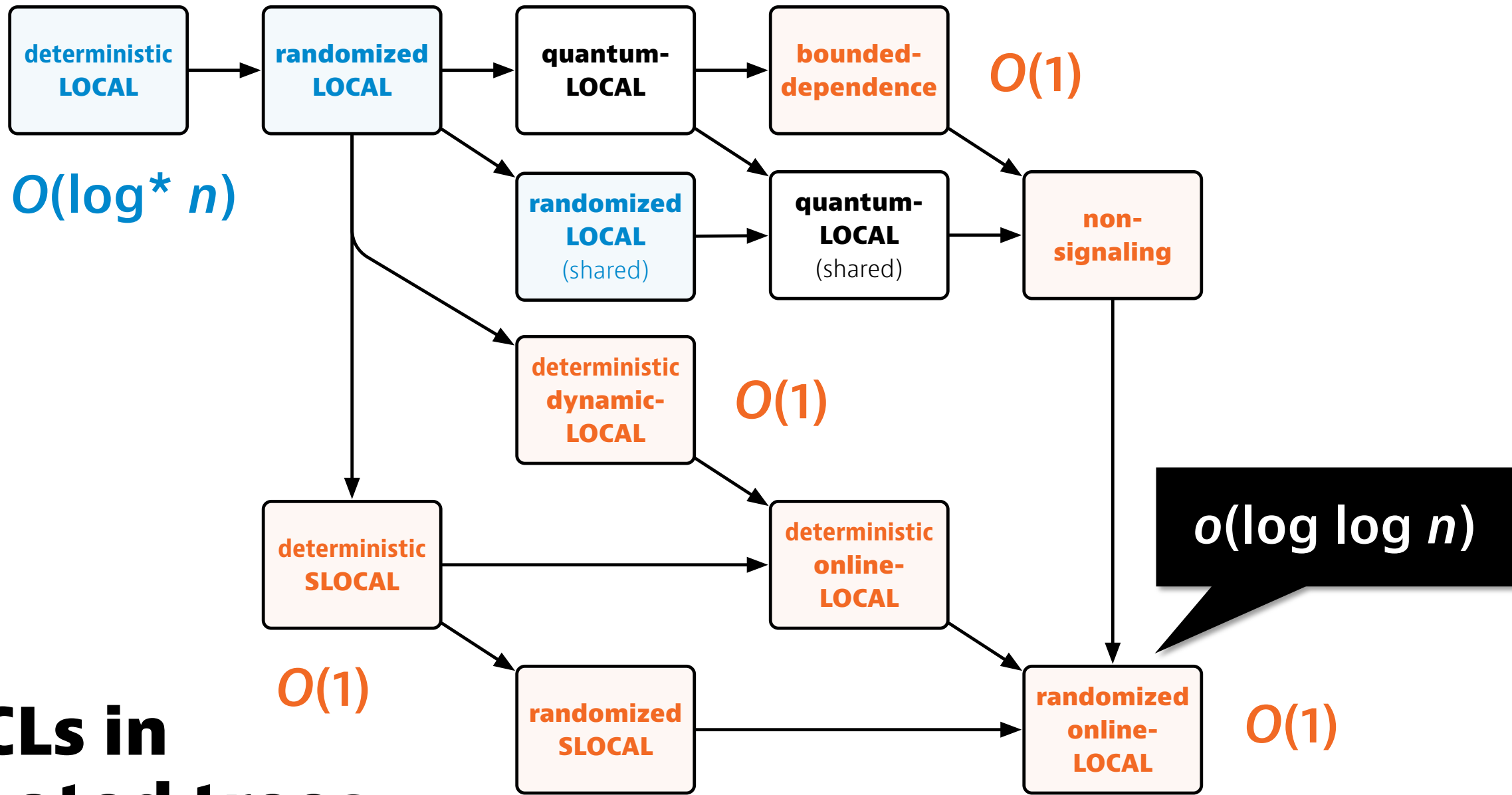






LCLs

- Finely dependent coloring in paths
- rooted trees
 - graphs
 - power graphs
 - fake IDs
 - solve any LCL



LCLs in rooted trees

$O(\log^* n)$

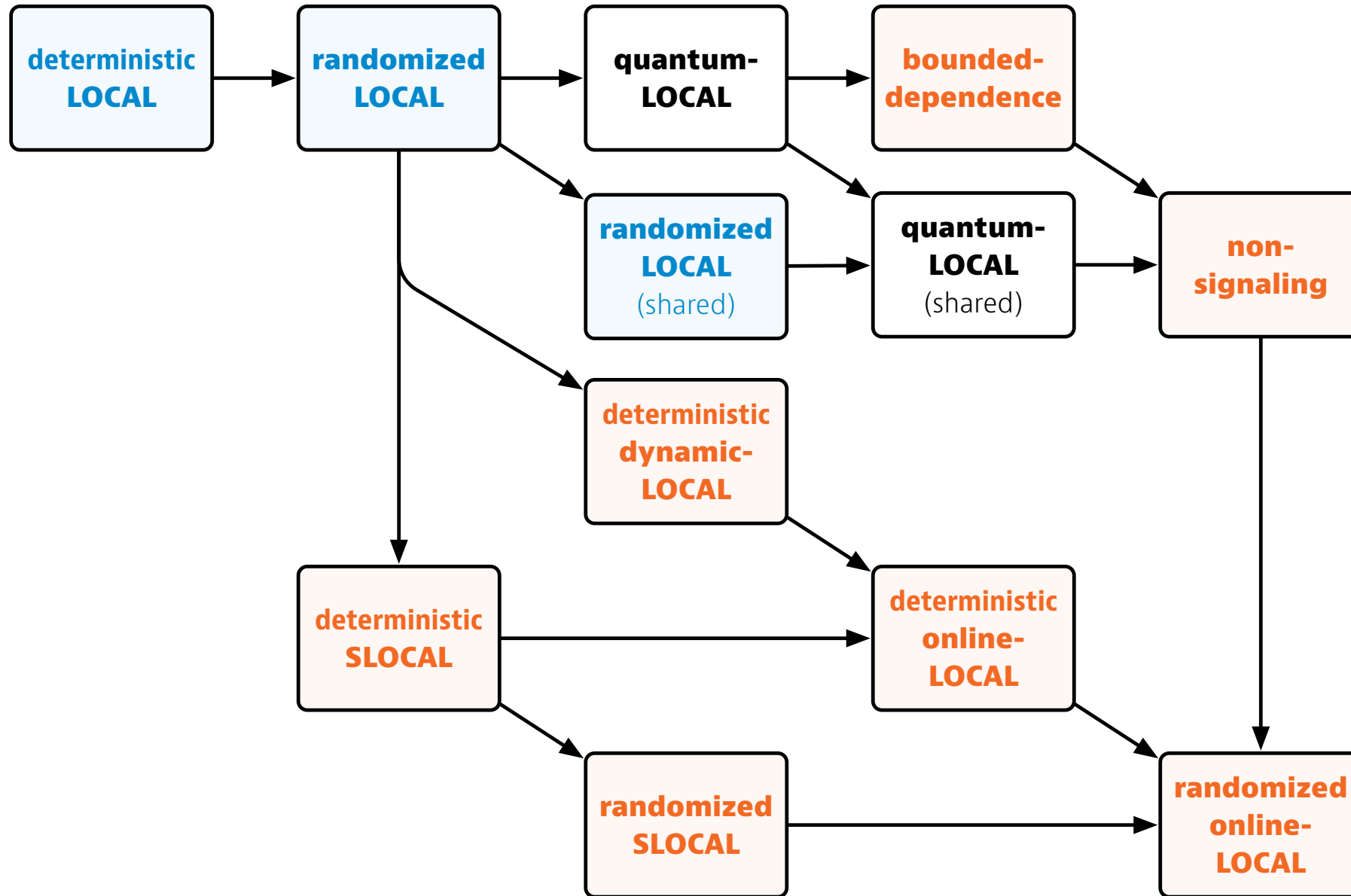
$O(1)$

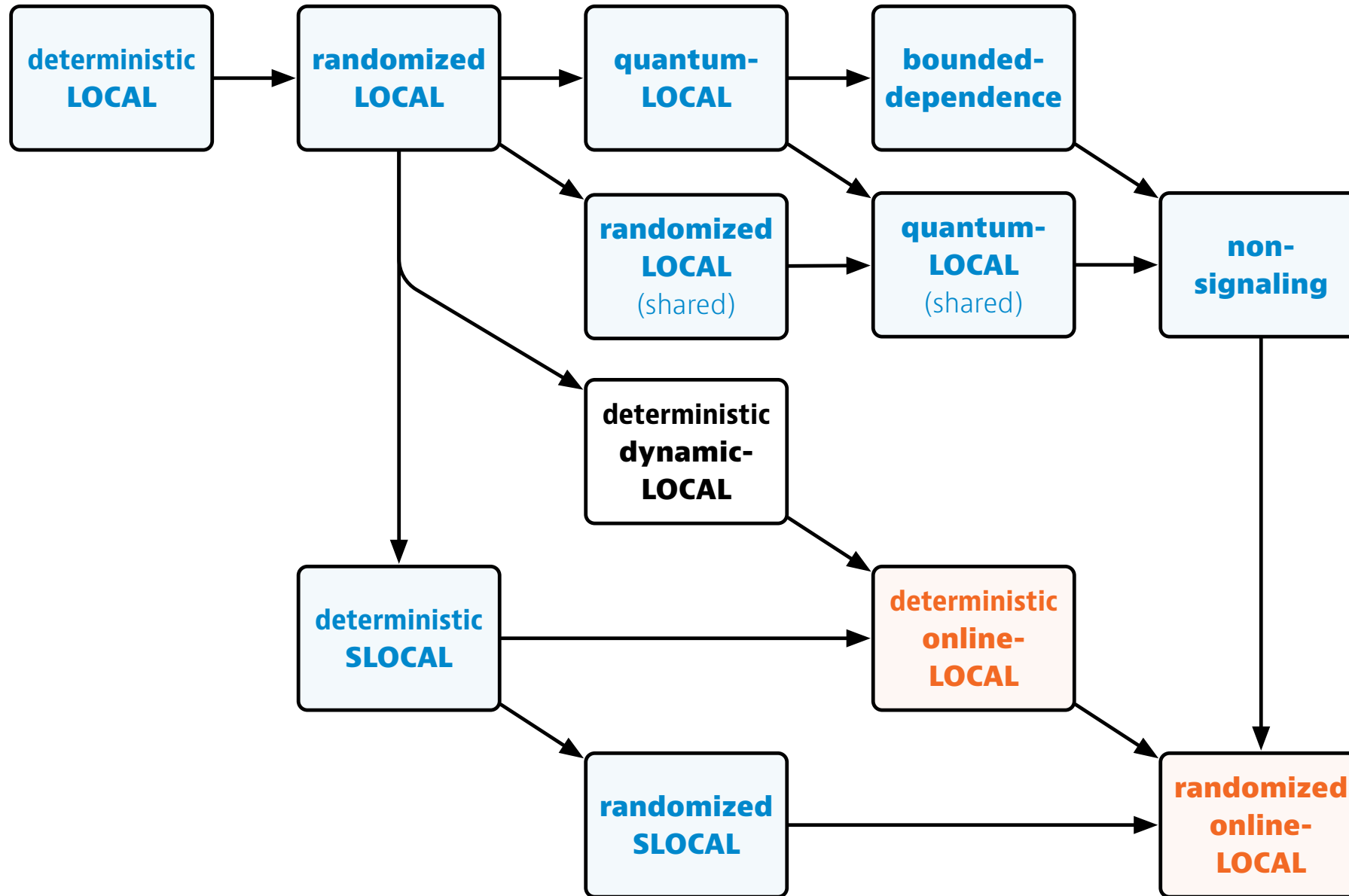
$O(1)$

$O(1)$

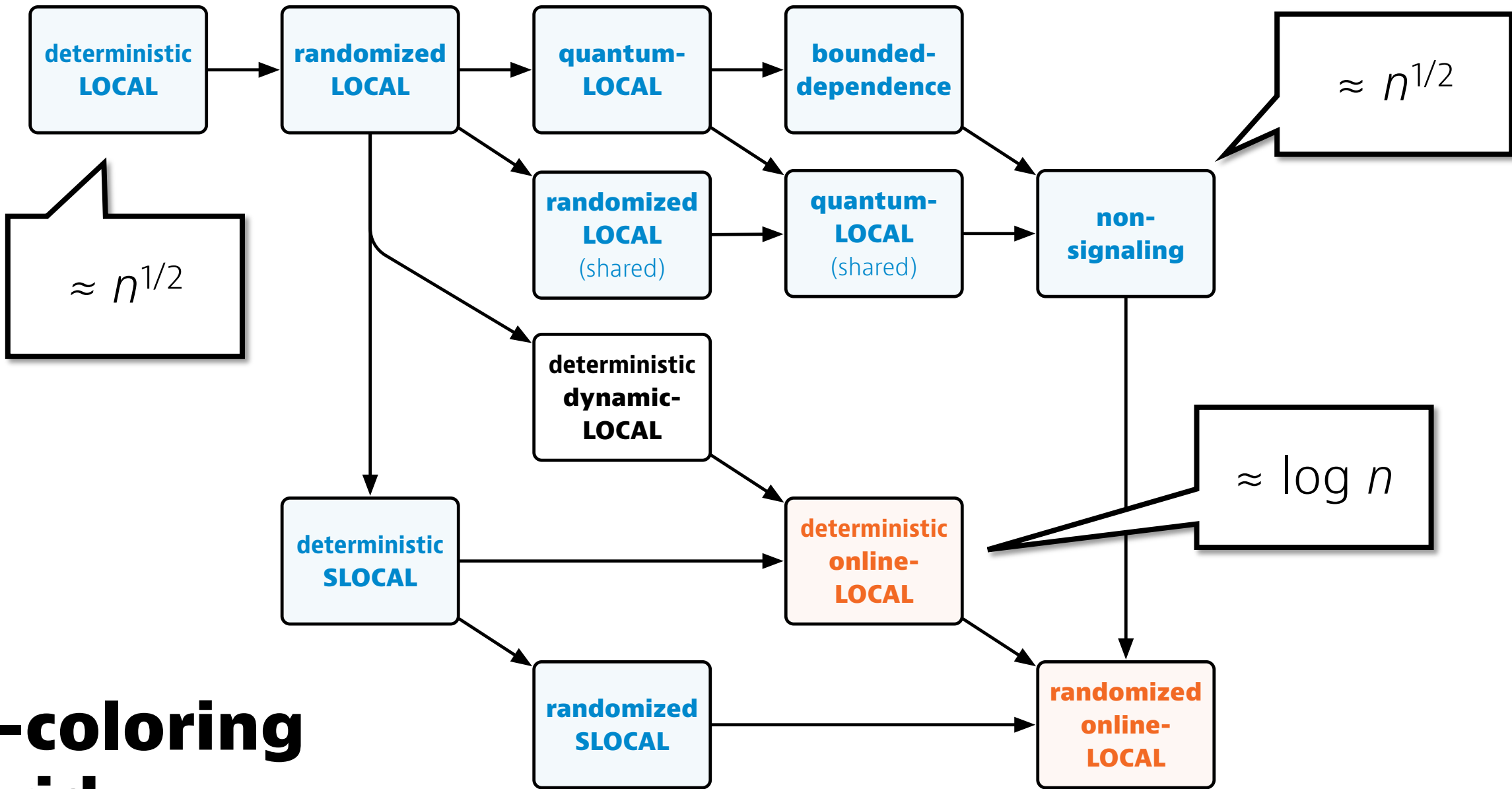
$o(\log \log n)$

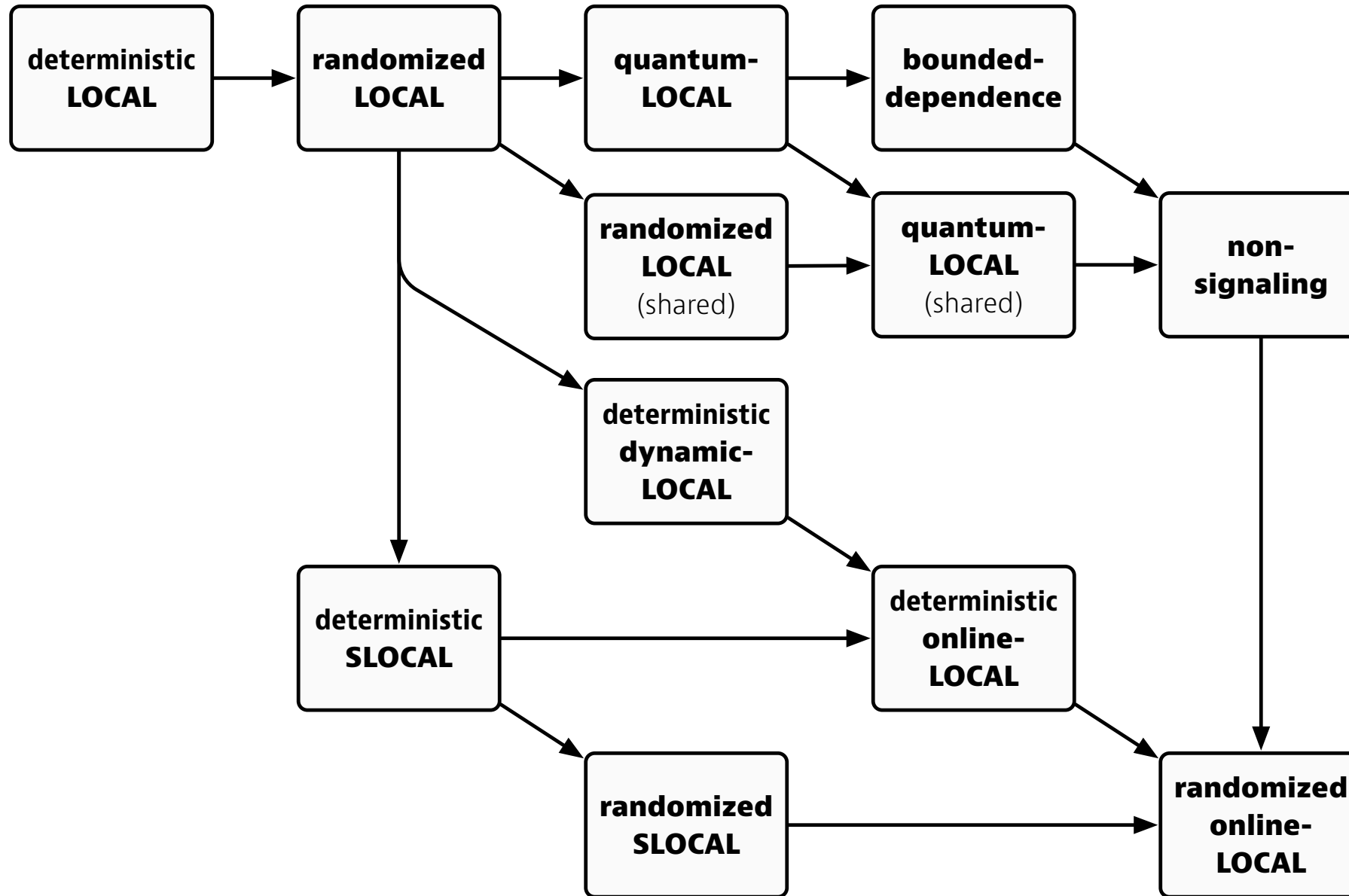
$O(1)$

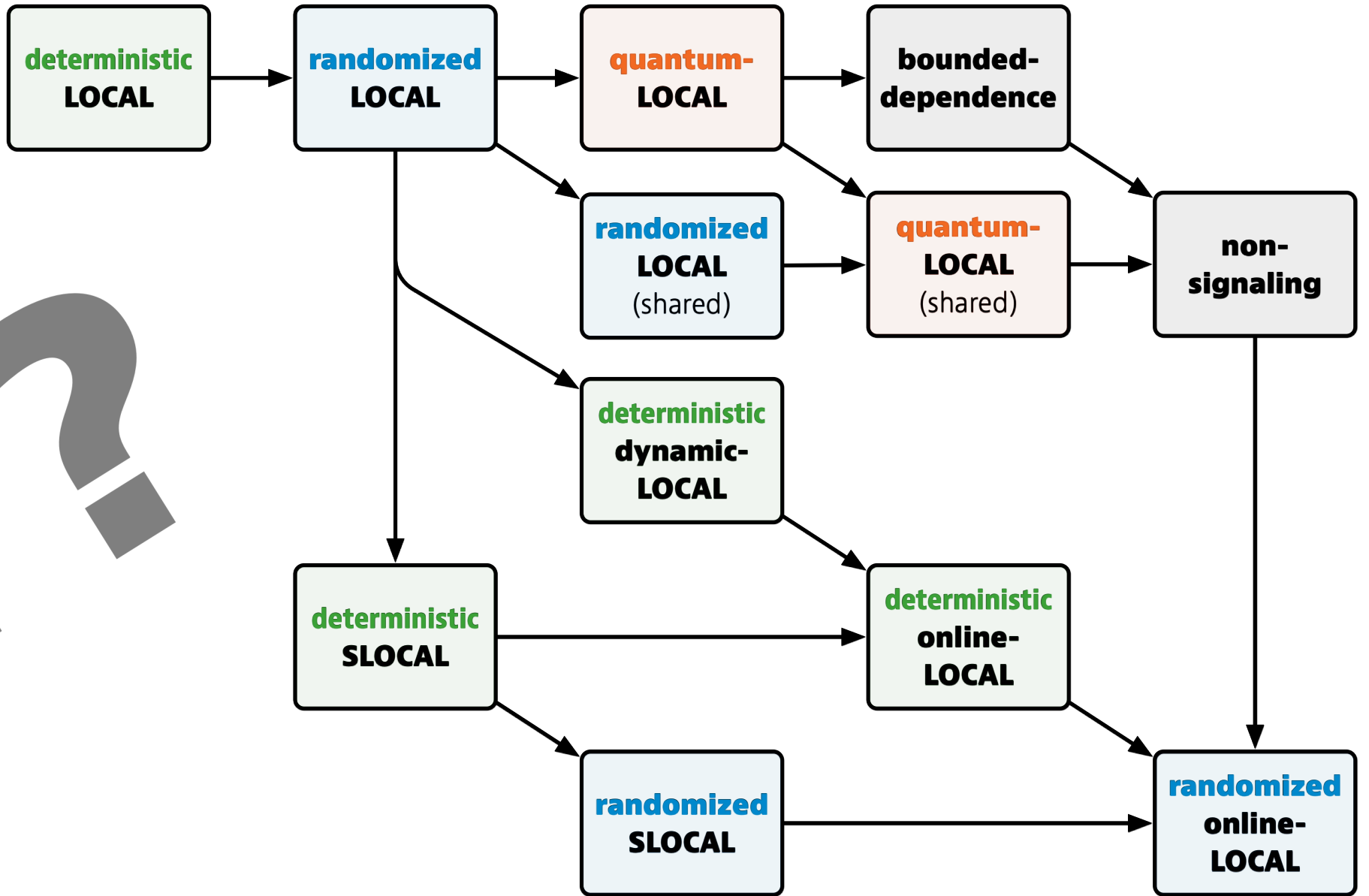




3-coloring grids







Open questions

- **Is there any distributed quantum advantage for LCLs?**
 - example: can you *color cycles* in $O(1)$ rounds in *quantum-LOCAL*?
- **Does global memory ever help in trees?**
 - could we *simulate online-LOCAL* in LOCAL model in unrooted trees?
 - or can we separate these for some LCL in trees?

Open questions

- **Lower bounds for sinkless orientation**
 - deterministic LOCAL: $\approx \log n$
 - randomized LOCAL: $\approx \log \log n$
 - deterministic SLOCAL: $\approx \log \log n$
 - randomized SLOCAL: $\approx \log \log \log n$
 - quantum-LOCAL: no lower bounds!
 - *we cannot exclude finitely dependent distributions for sinkless orientation!*

Open questions

- **New lower-bound proof techniques?**
 - *existential graph-theoretic arguments* and propagation arguments seem too weak to tackle e.g. sinkless orientation
 - *round elimination* does not work in quantum and beyond
 - how far can we stretch Marks' technique?
 - current best hope: new *simulation results*?

References & pointers

[arXiv:2109.06593](https://arxiv.org/abs/2109.06593) Locality in online, dynamic, sequential, and distributed graph algorithms (*ICALP 2023*)

[arXiv:2307.09444](https://arxiv.org/abs/2307.09444) No distributed quantum advantage for approximate graph coloring (*STOC 2024*)

[arXiv:2403.01903](https://arxiv.org/abs/2403.01903) Online locality meets distributed quantum computing

jukkasuomela.fi/talks