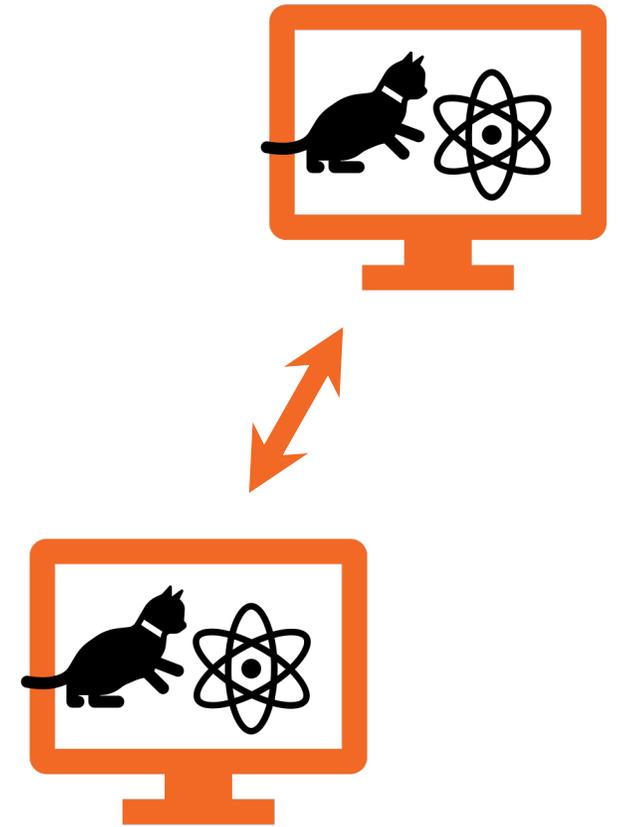


# *Distributed Quantum Advantage*



**Jukka Suomela**  
Aalto University



# **SOFSEM 2006**

Merin, Czech Republic

# Distributed algorithms

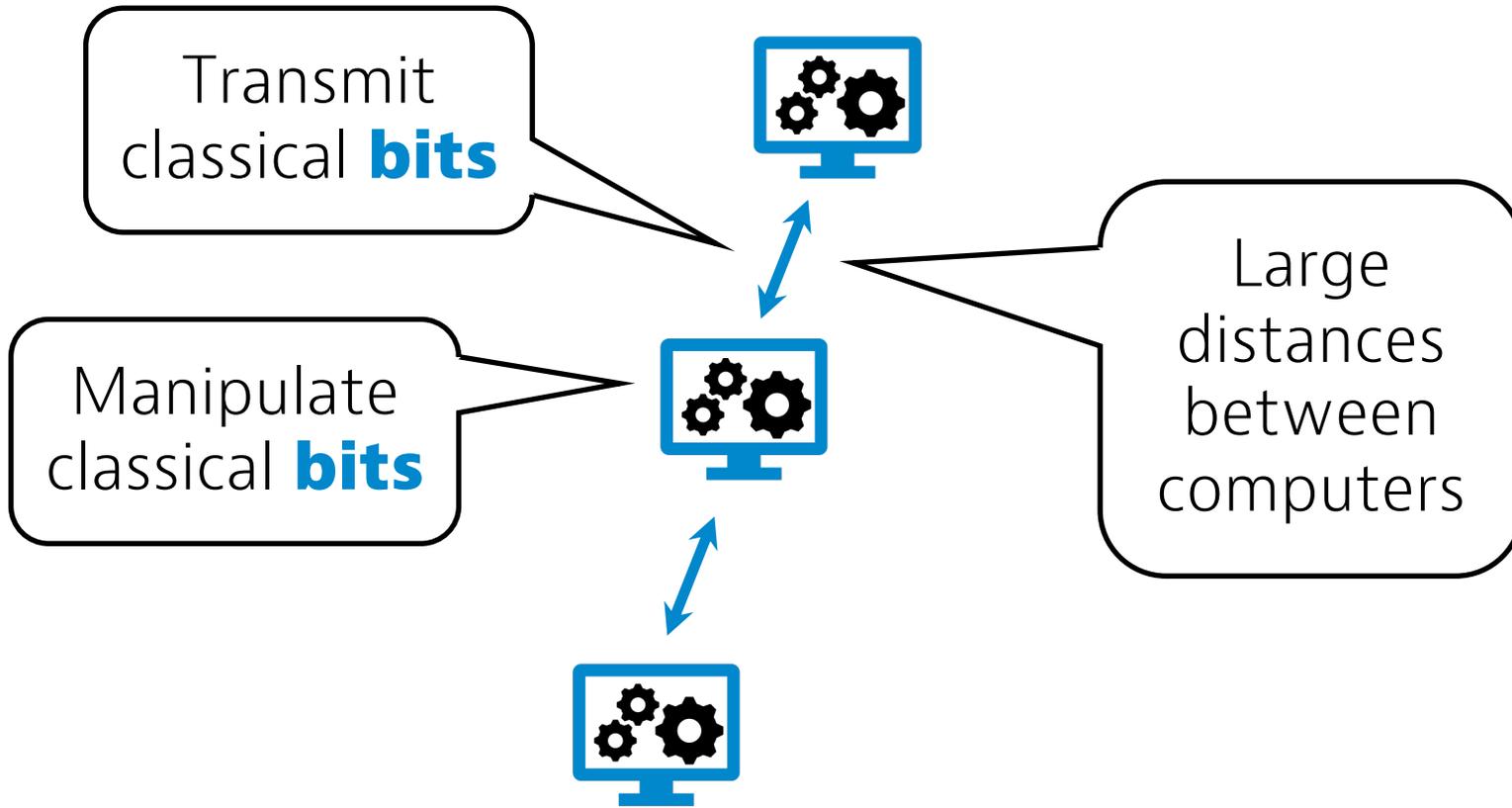
- What tasks can be solved efficiently in a very large computer network?
- What are the fundamental limits?
- Does **quantum** change any of this?

**Distributed  
quantum  
computing?**

Manipulate  
classical **bits**



Classical  
computer

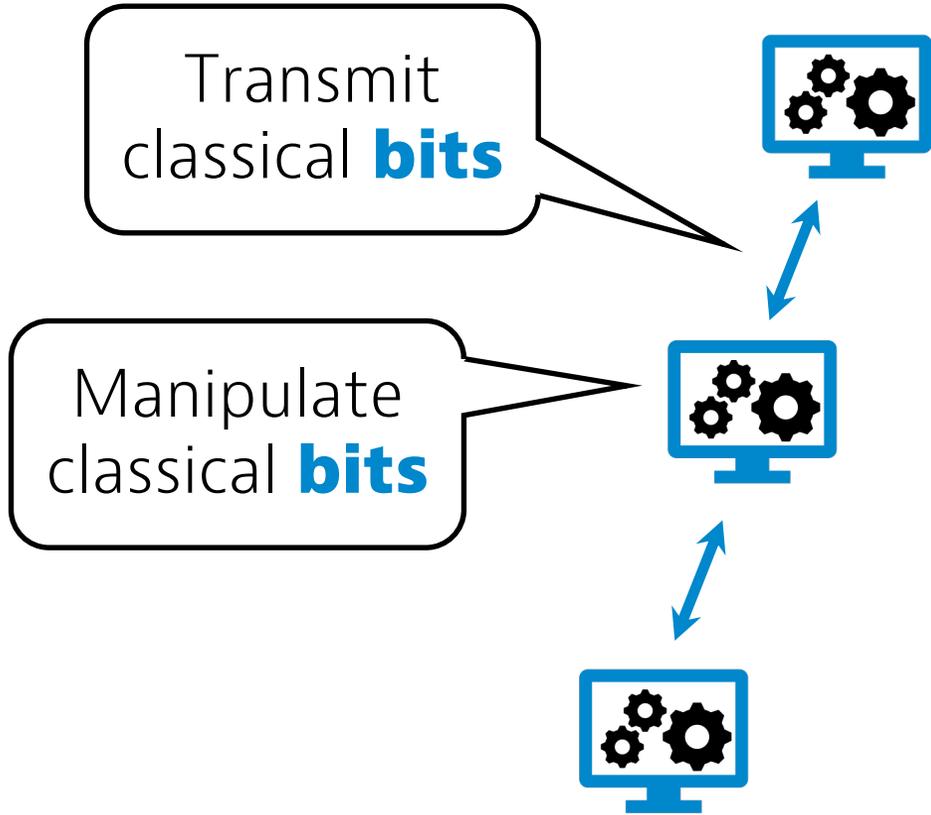


Transmit  
classical **bits**

Manipulate  
classical **bits**

Large  
distances  
between  
computers

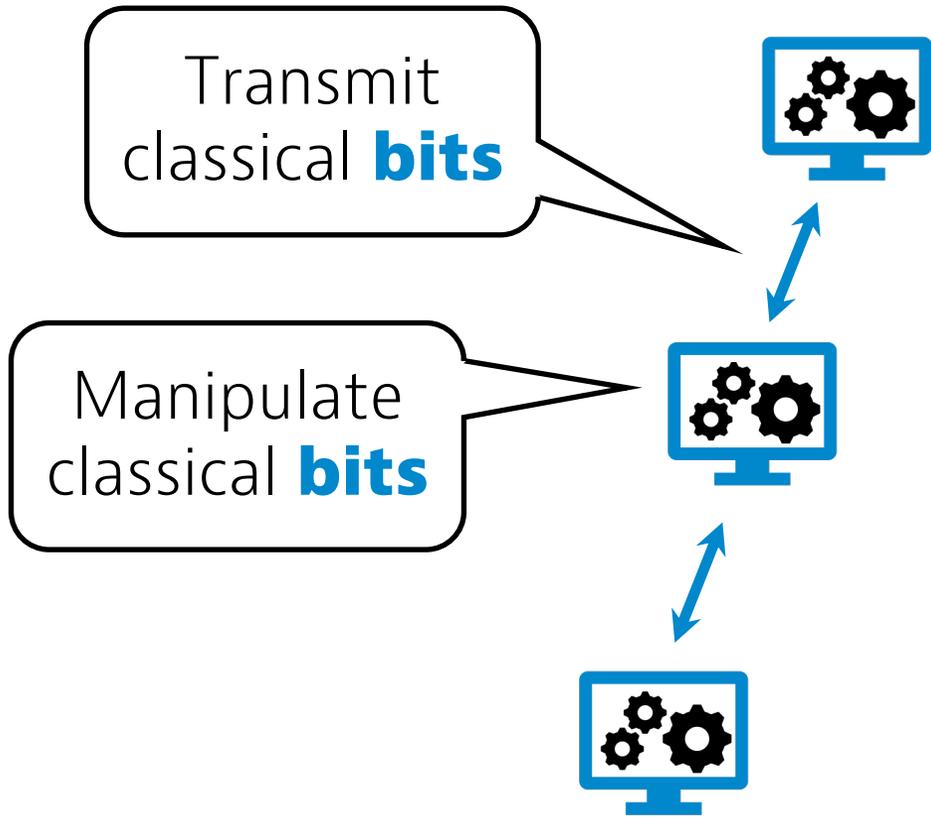
Classical  
computer  
network



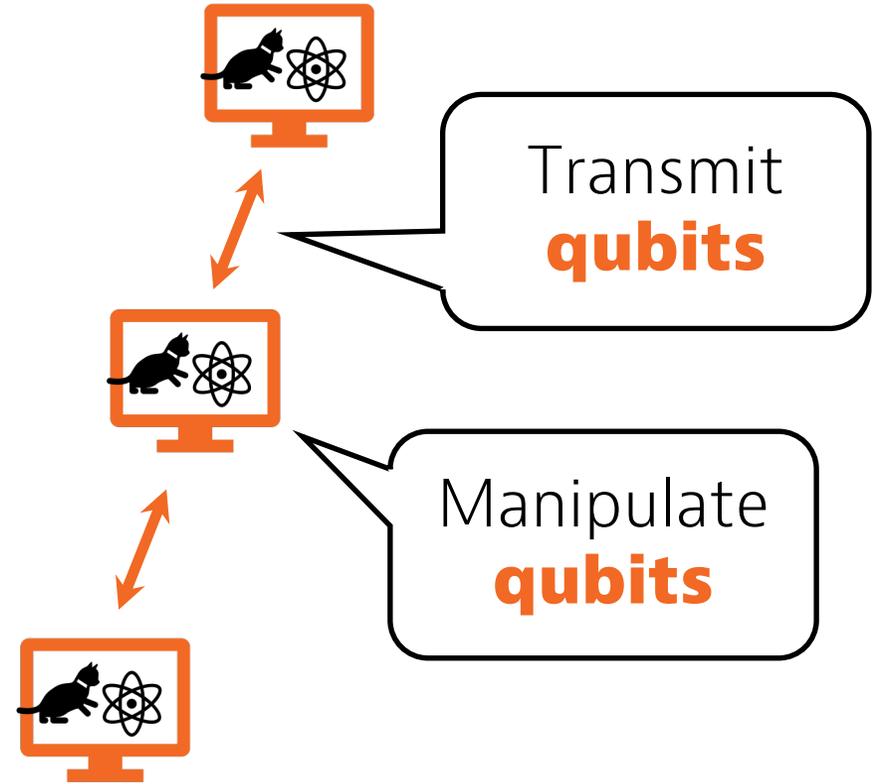
Classical  
computer  
network



Quantum  
computer



Classical  
computer  
network

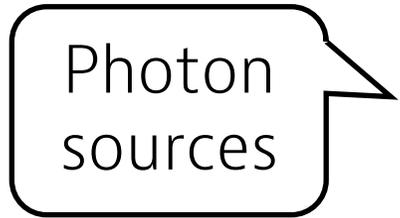


Quantum  
computer  
network

# Quantum circuits

“Real-world” example: **qubit = photon**

Photon  
sources



e.g. always  
known  
polarization



# "Real-world" example: **qubit = photon**

Photon sources

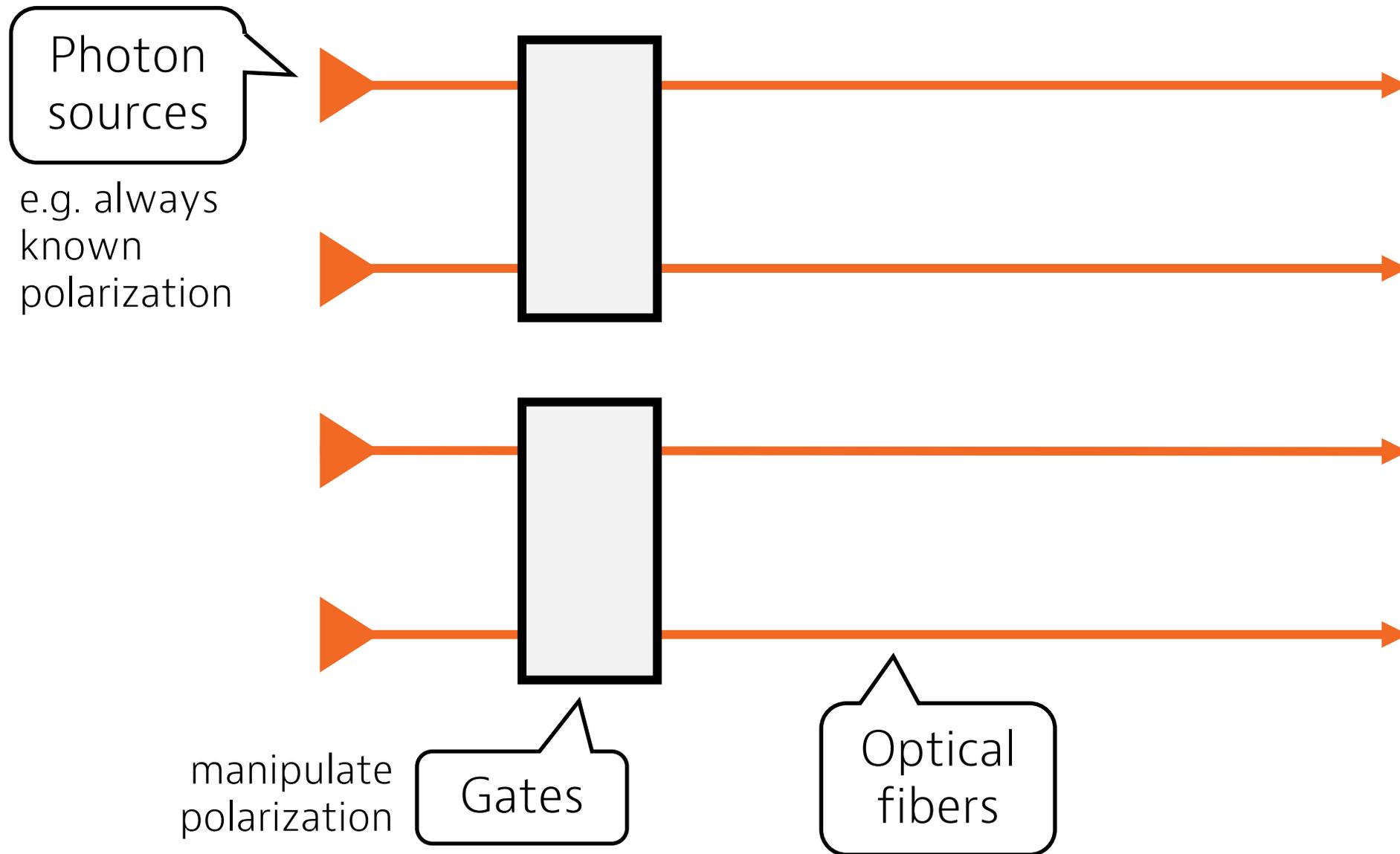
e.g. always known polarization



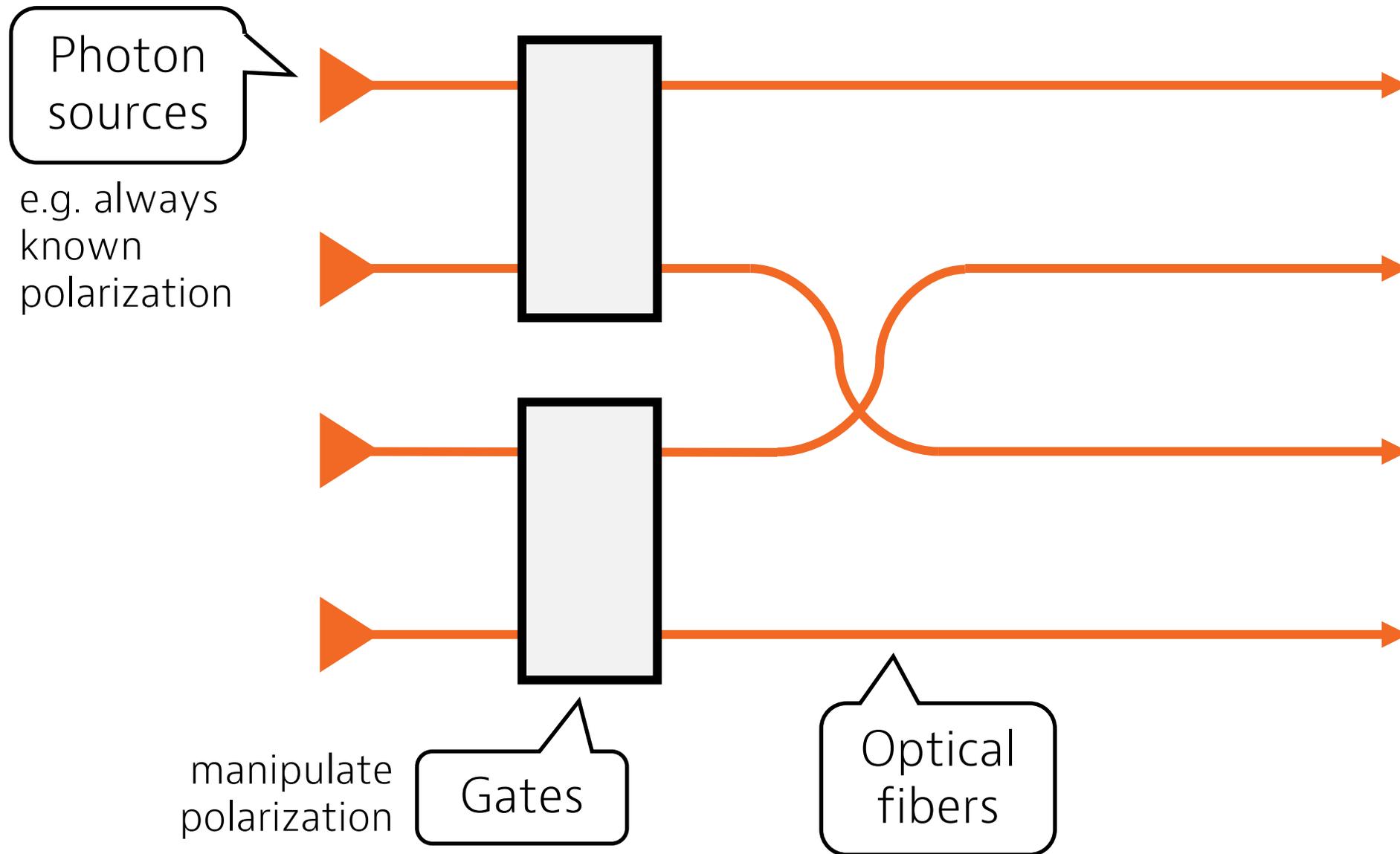
Optical fibers

do not change photons

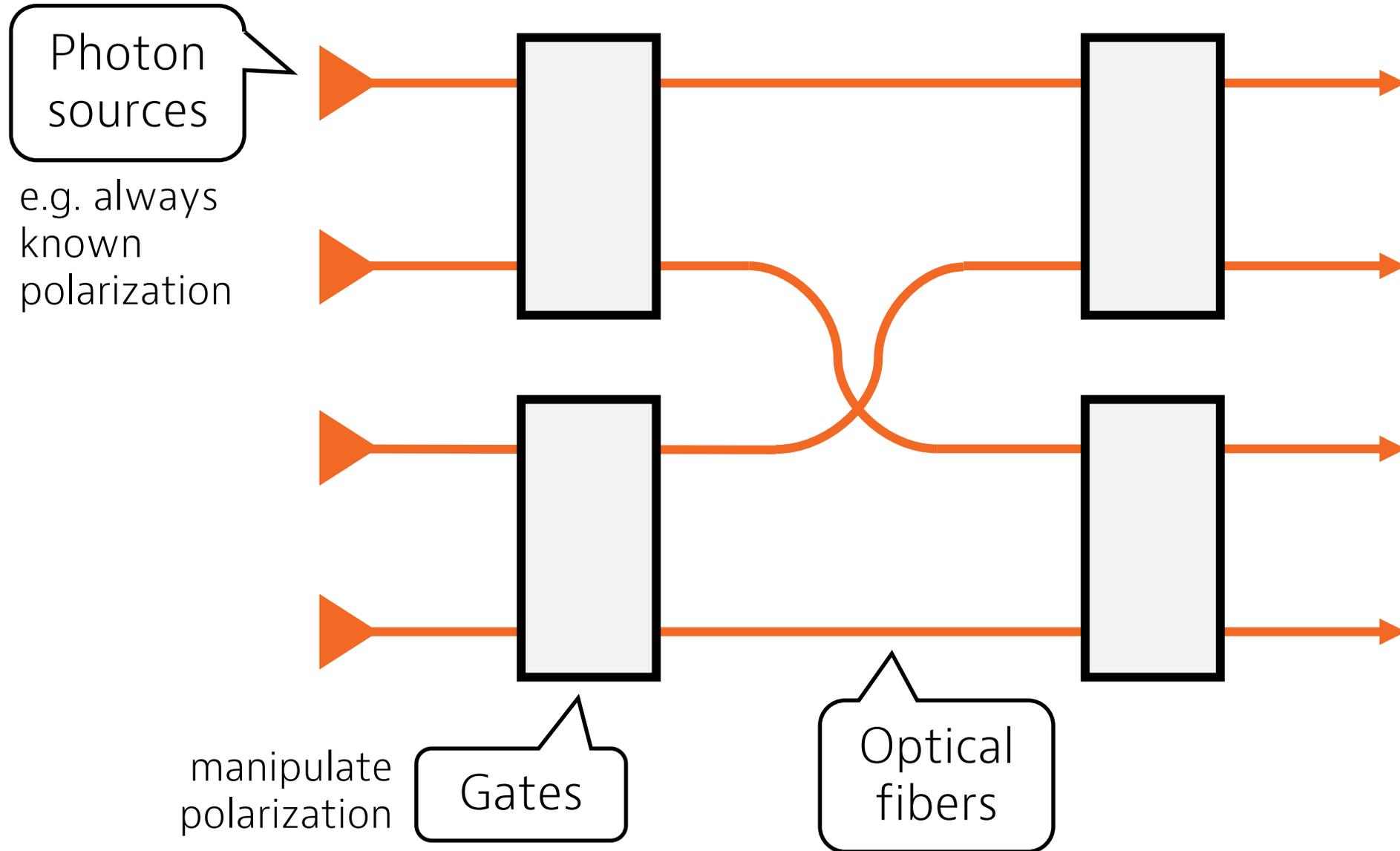
# "Real-world" example: **qubit = photon**



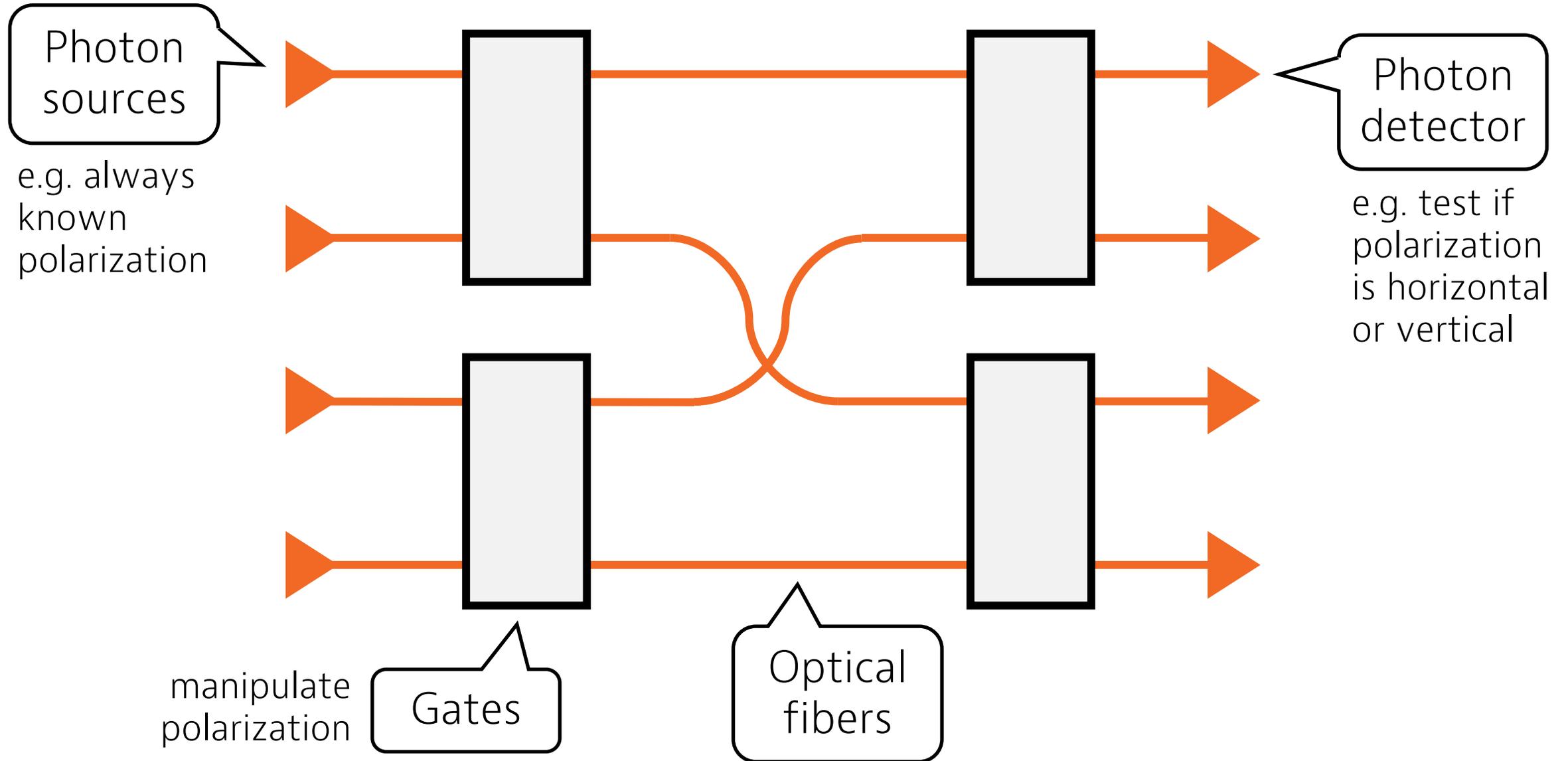
# "Real-world" example: **qubit = photon**



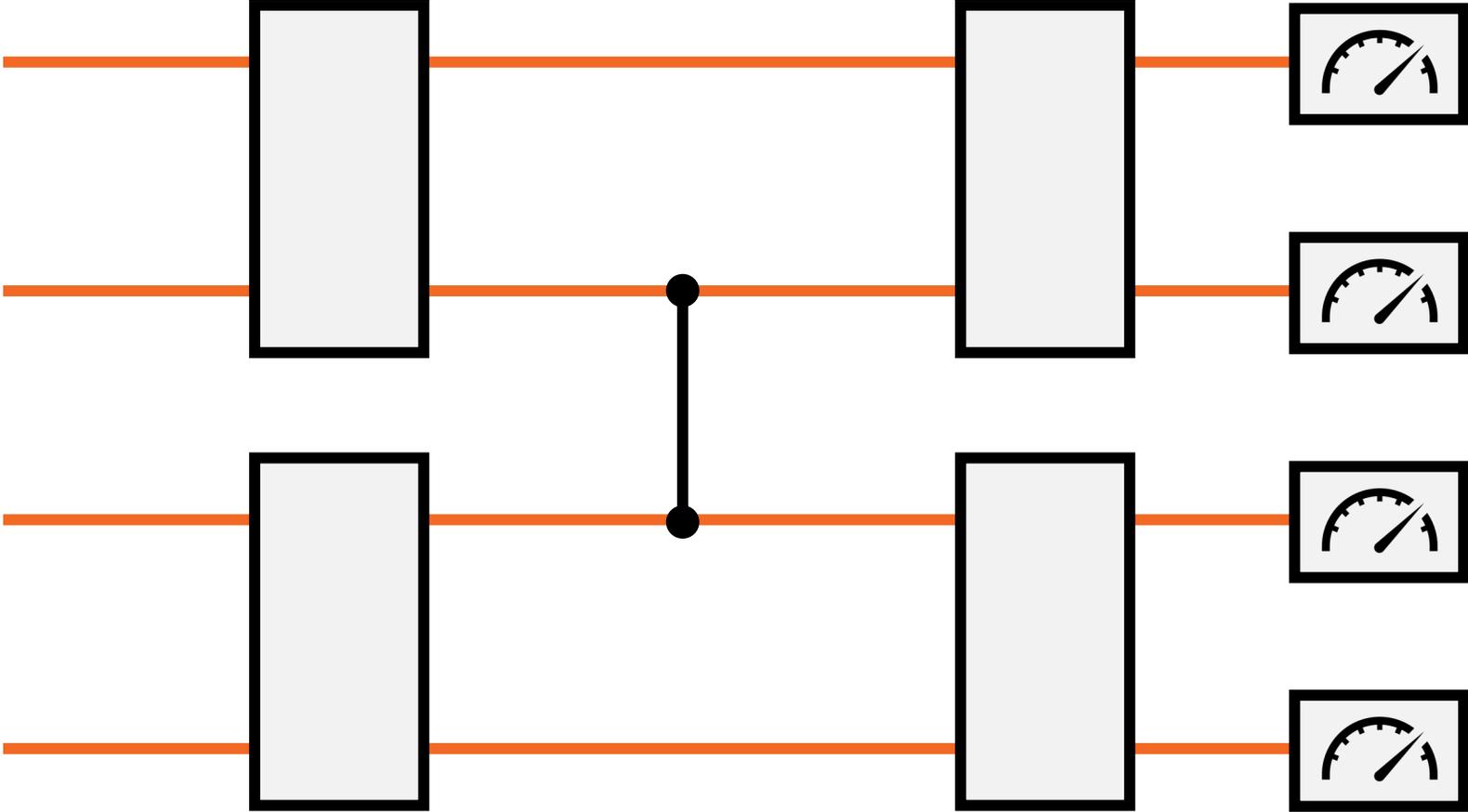
# "Real-world" example: **qubit = photon**



# "Real-world" example: **qubit = photon**

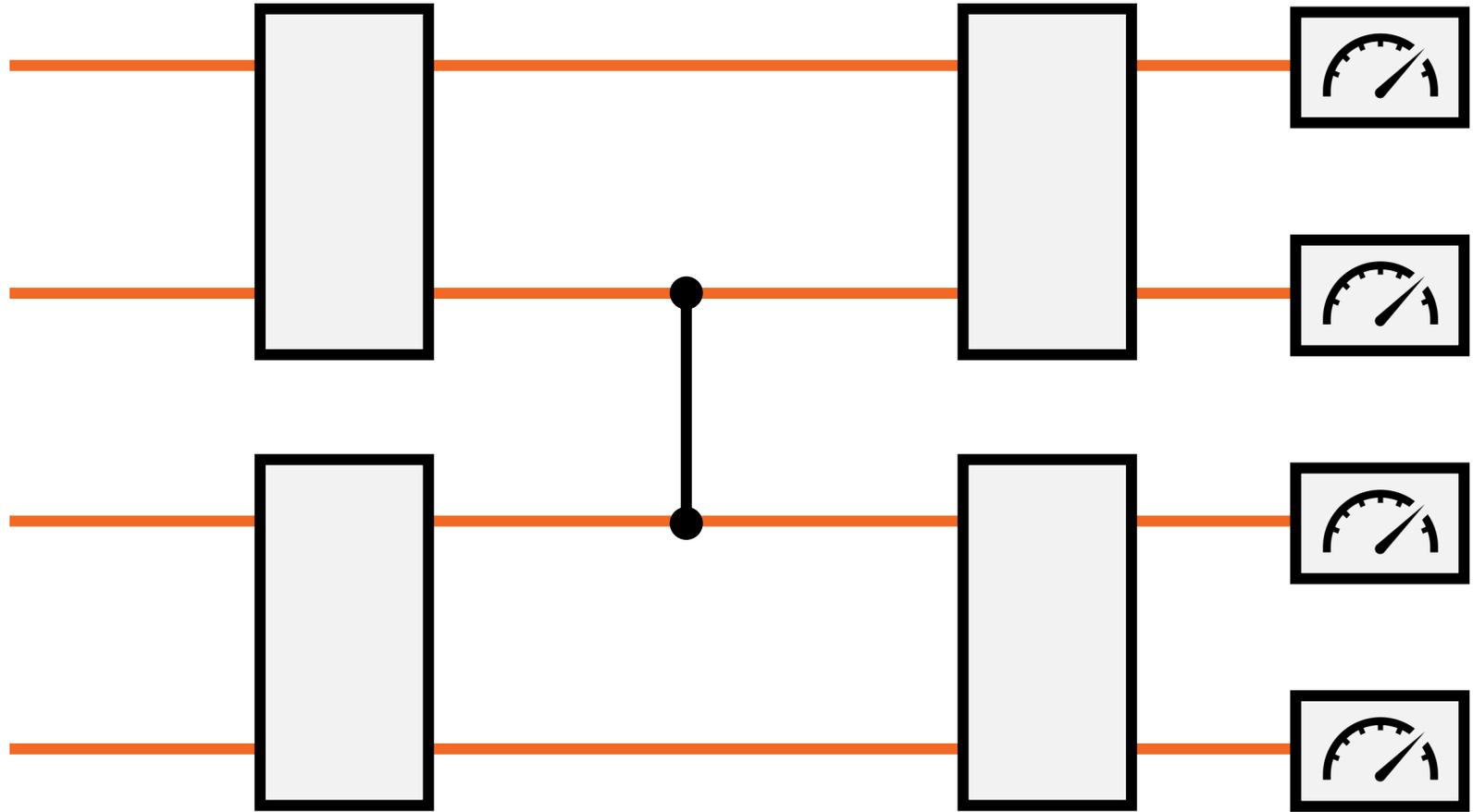


# Quantum circuit



# "Shut up and calculate"

State vector:  
 $2^4 = 16$   
complex amplitudes

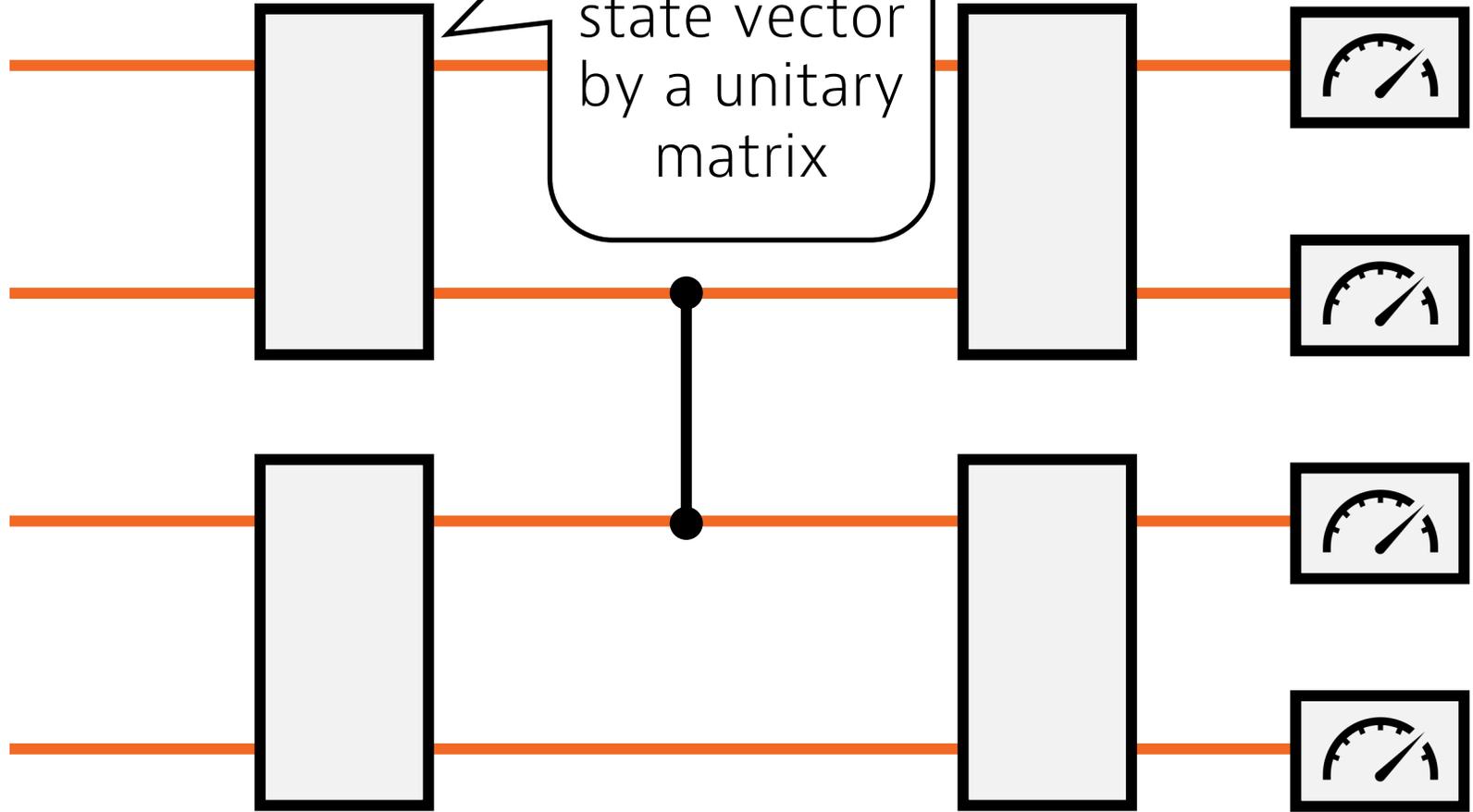


- $c_{0000}$
- $c_{0001}$
- $c_{0010}$
- ...
- $c_{1111}$

# "Shut up and calculate"

unitary matrix: inverse  
= conjugate transpose

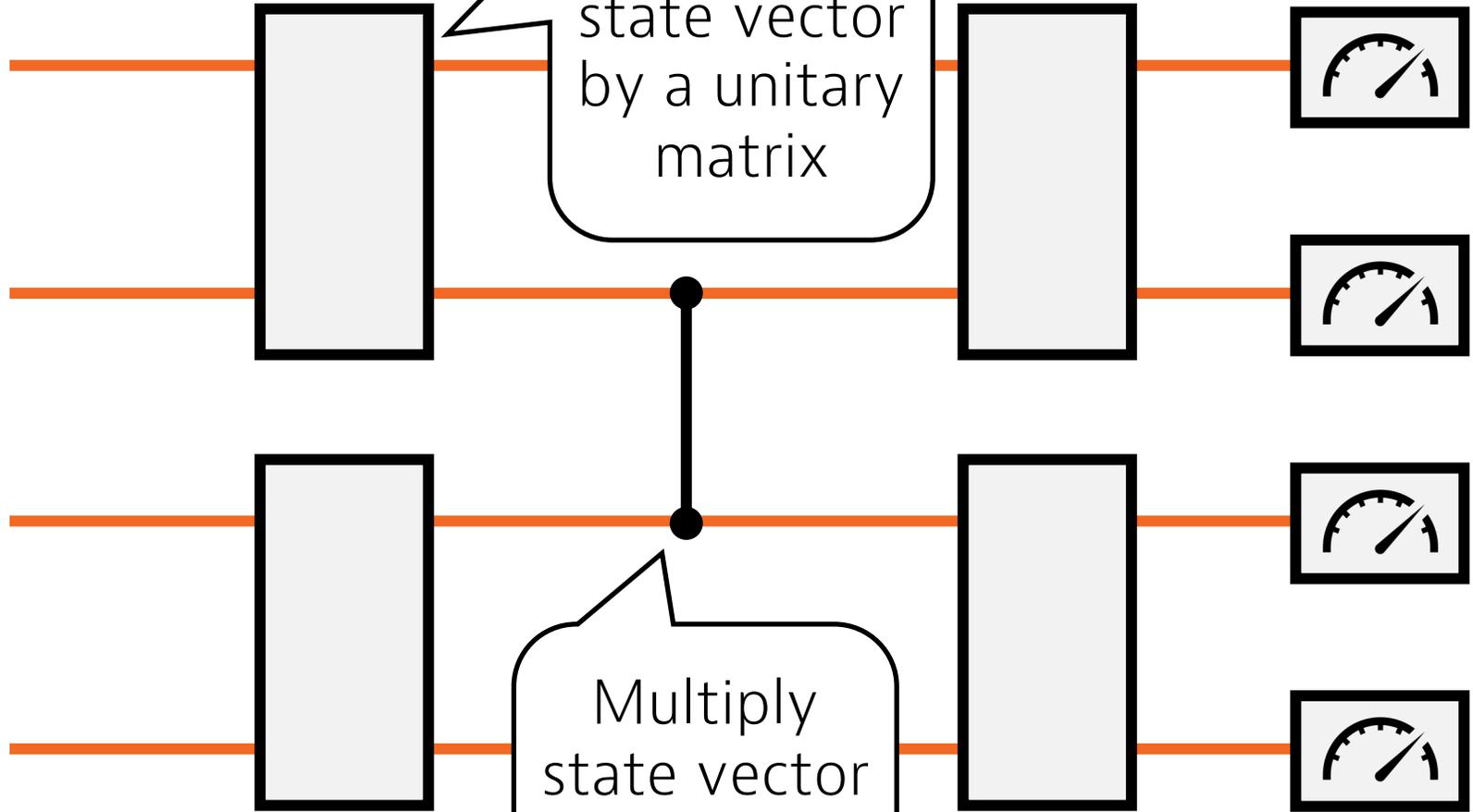
State vector:  
 $2^4 = 16$   
complex amplitudes



- $c_{0000}$
- $c_{0001}$
- $c_{0010}$
- ...
- $c_{1111}$

# "Shut up and calculate"

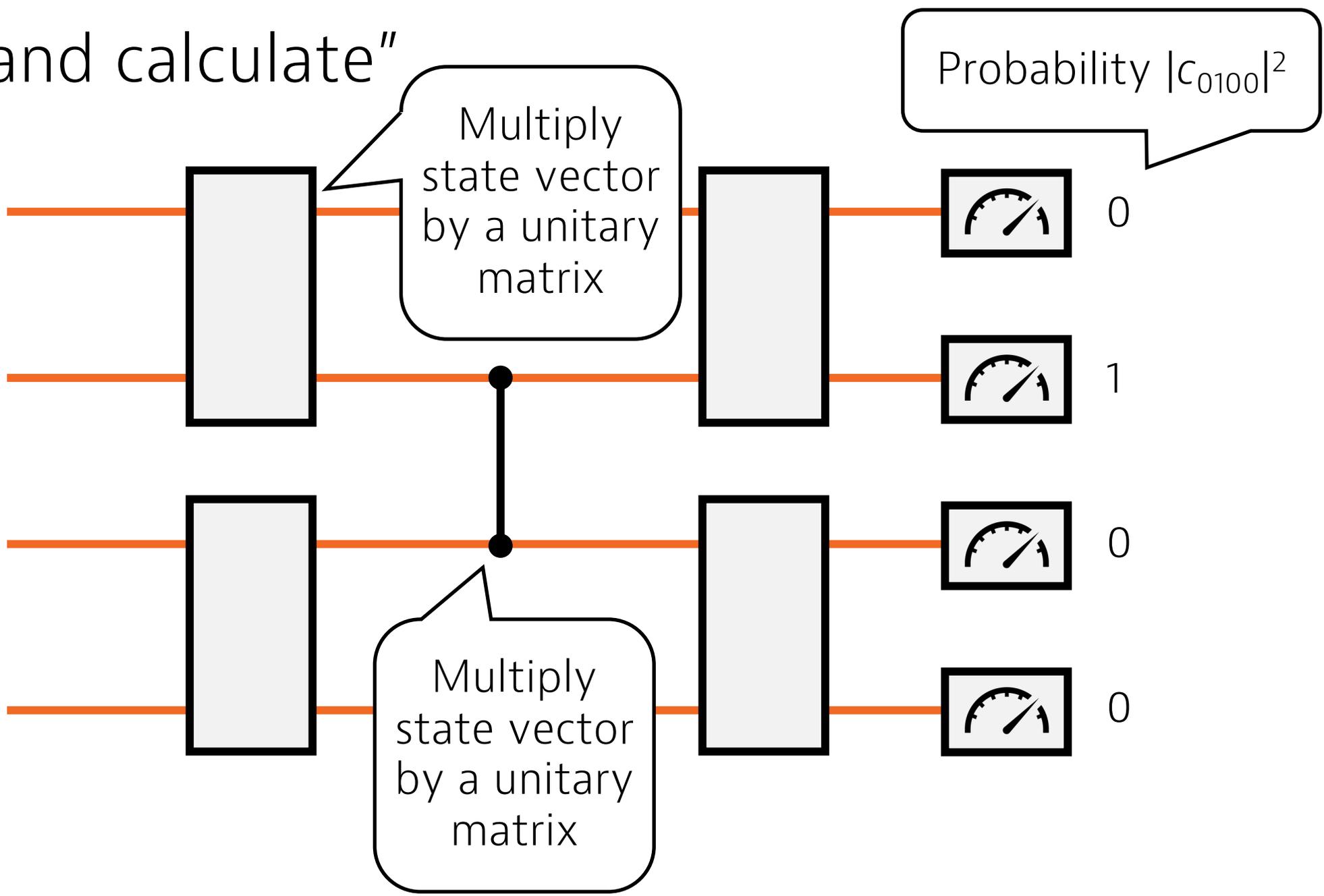
State vector:  
 $2^4 = 16$   
complex amplitudes



- $c_{0000}$
- $c_{0001}$
- $c_{0010}$
- ...
- $c_{1111}$

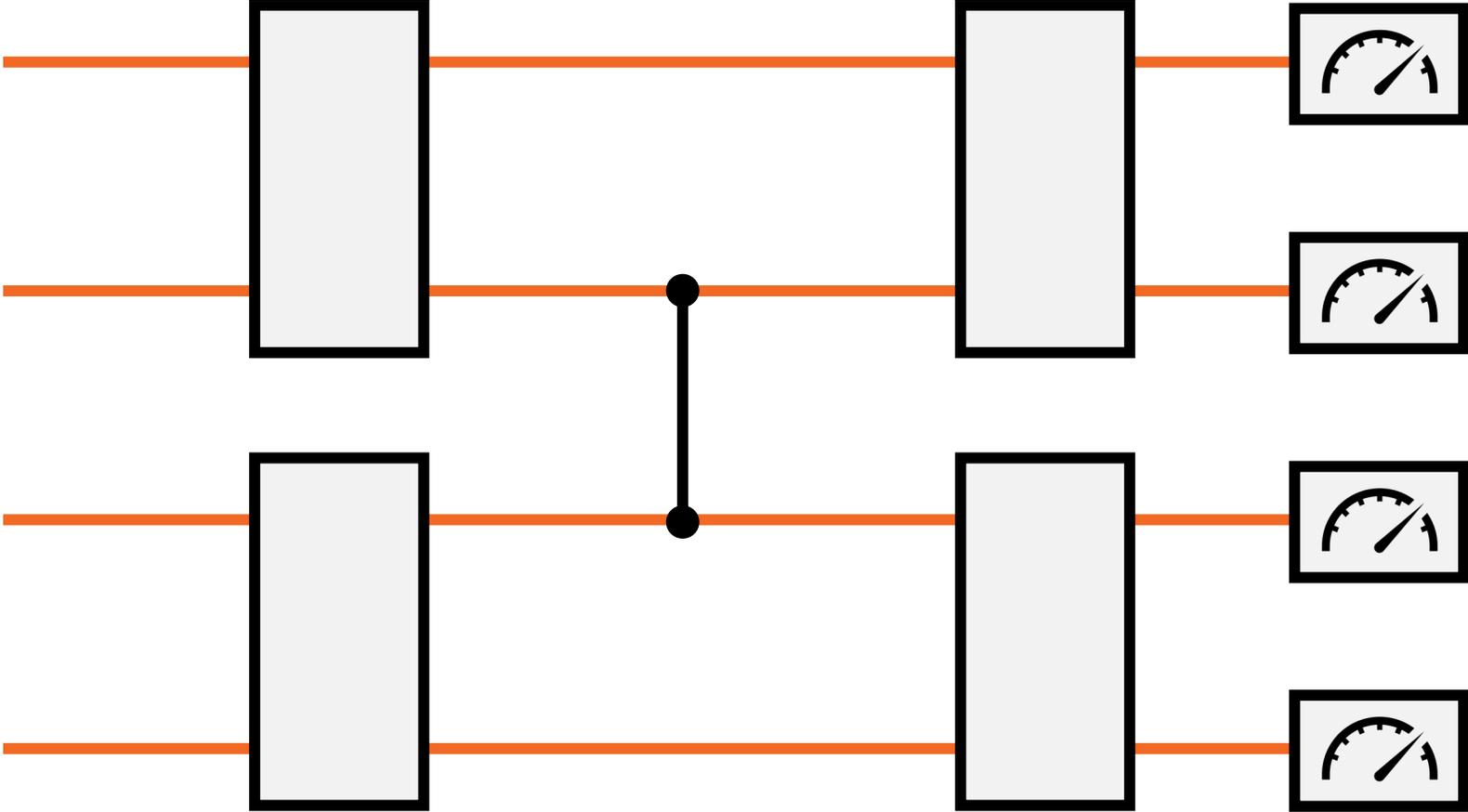
# "Shut up and calculate"

State vector:  
 $2^4 = 16$   
complex amplitudes



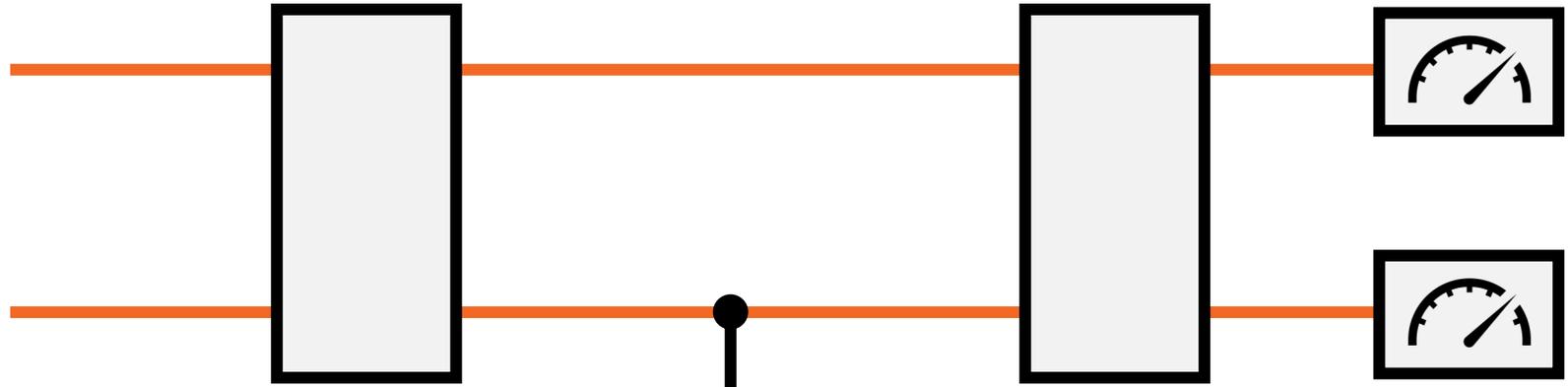
- $c_{0000}$
- $c_{0001}$
- $c_{0010}$
- ...
- $c_{1111}$

# Quantum circuit

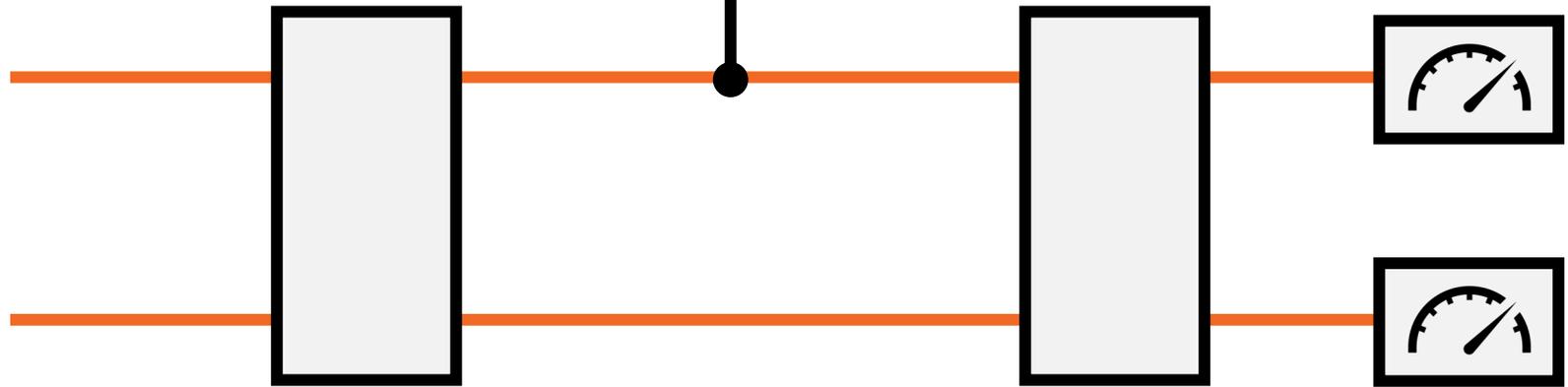


# Quantum circuit "stretched" between two labs

**Alice**

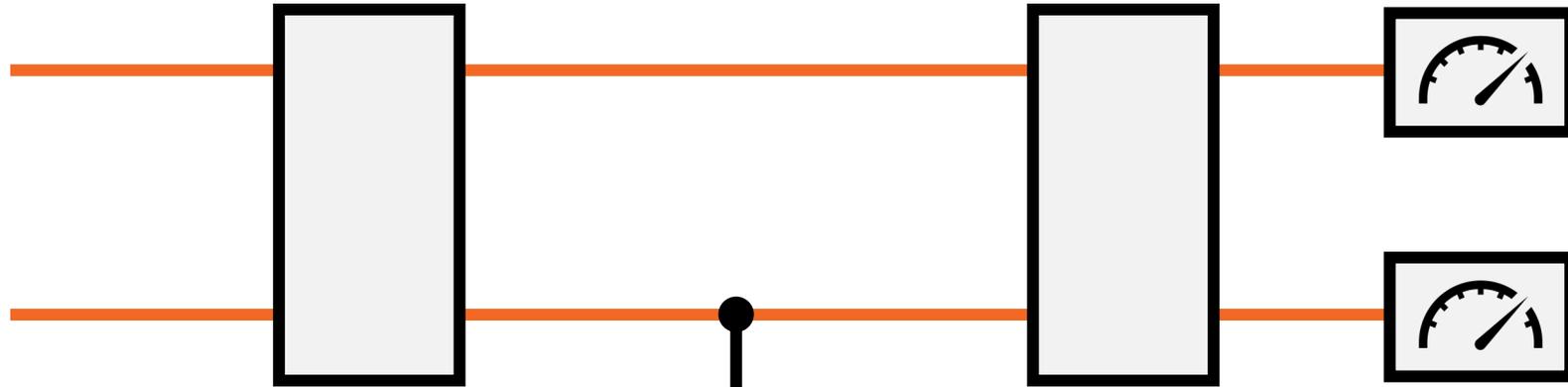


**Bob**

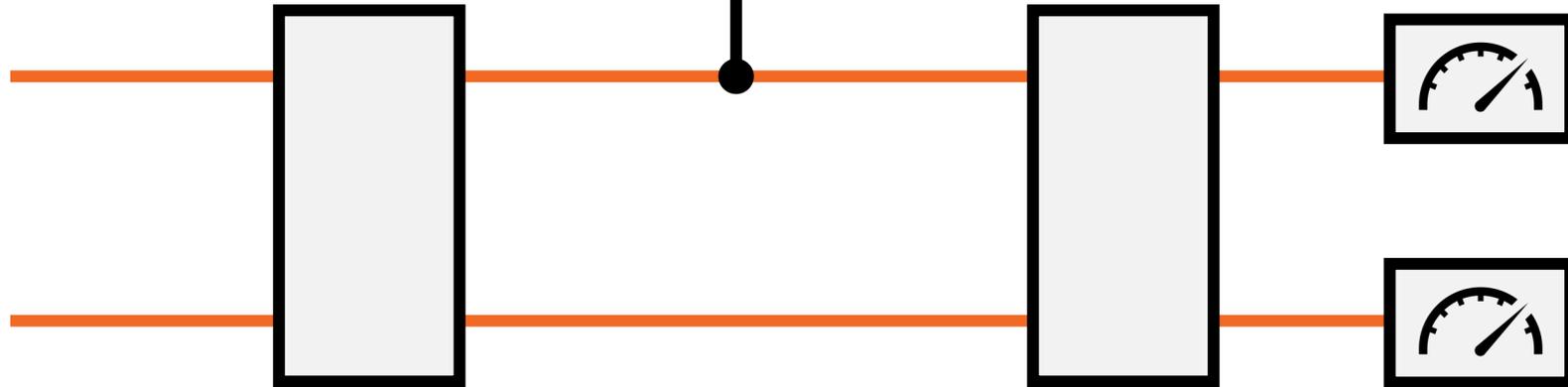


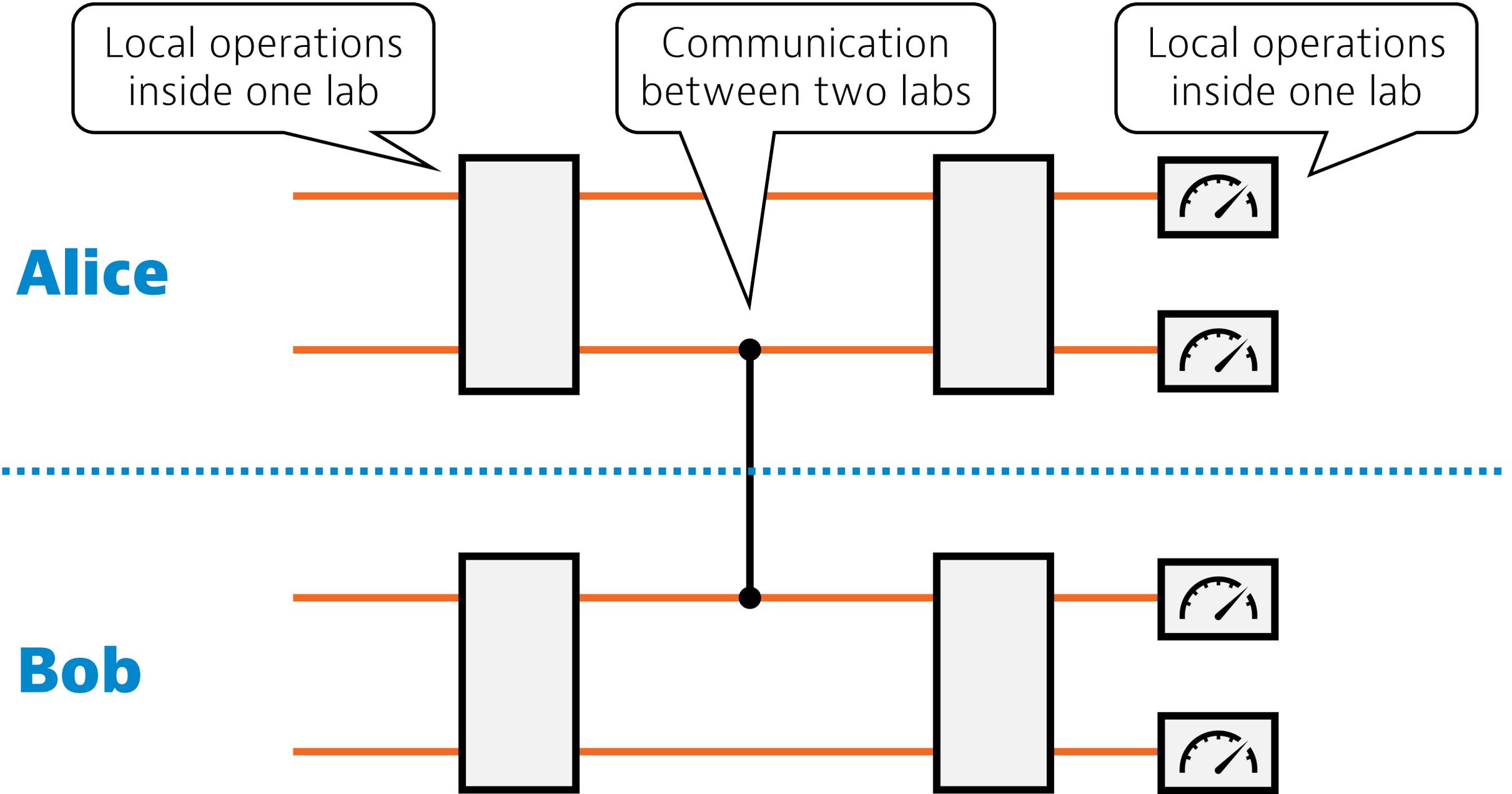
# Quantum circuit "stretched" between two labs

**Alice**



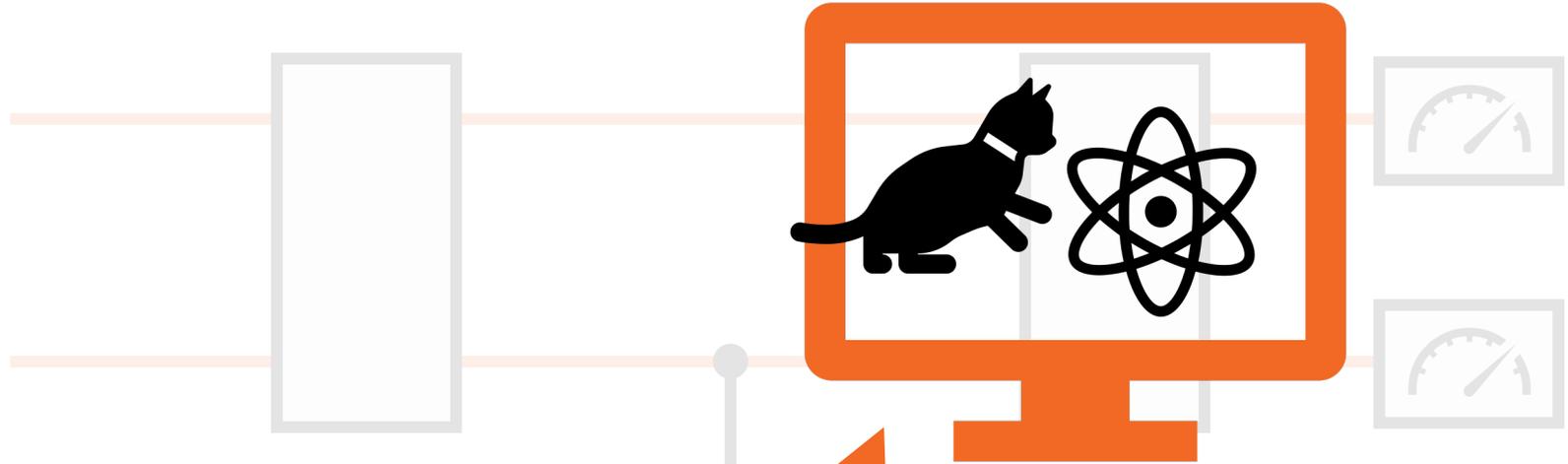
**Bob**



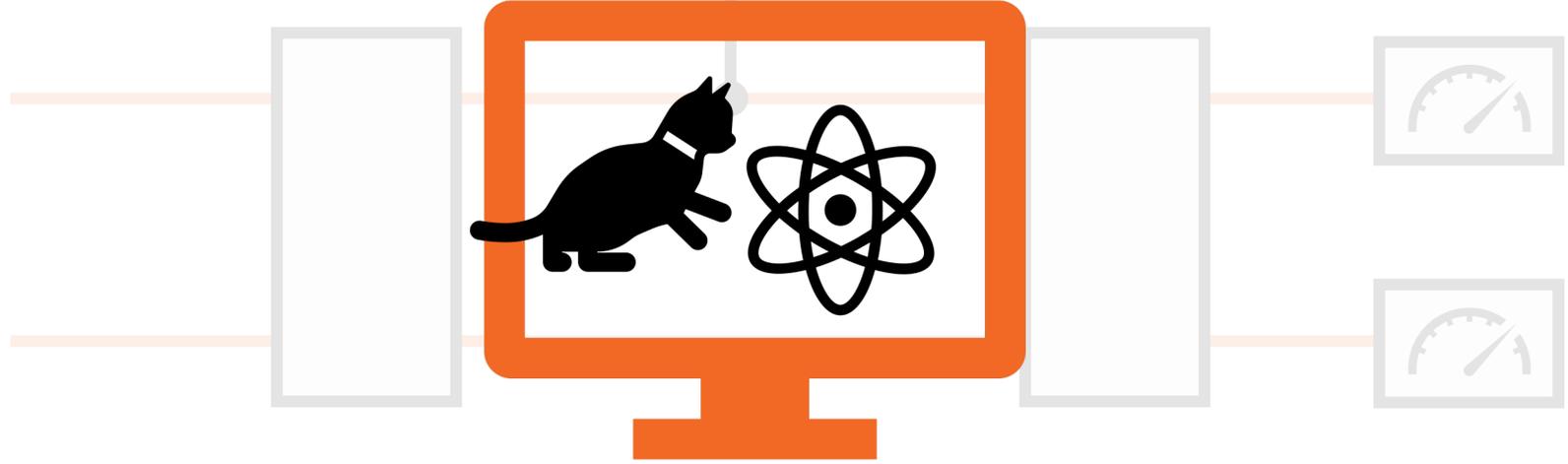


# Quantum computer network

**Alice**



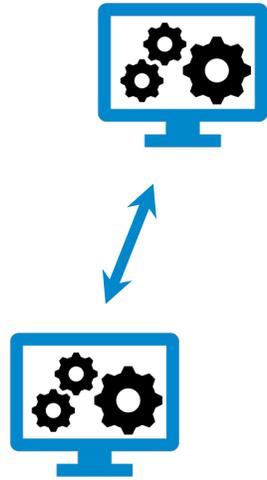
**Bob**



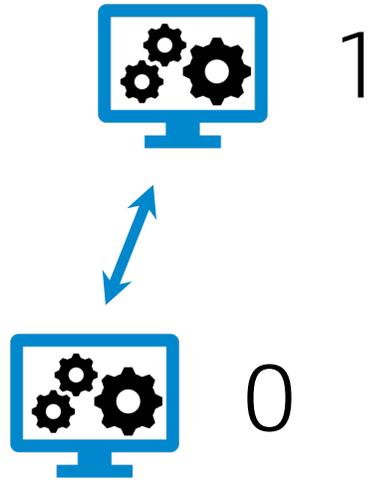
**Not really  
practical yet...**

**... but what  
could we do if  
this becomes  
reality?**

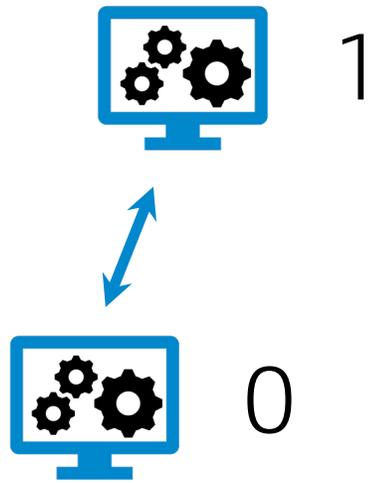
Two **identical** computers...



Two **identical** computers,  
exactly one must output 1



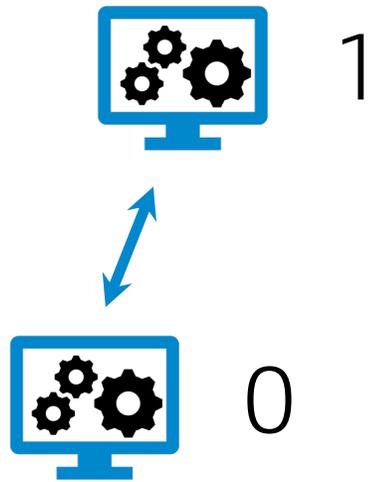
Two **identical** computers,  
exactly one must output 1



**Classical** algorithm:

- produce 100 random bits
- send to the other party
- larger value  $\rightarrow$  output 1

Two **identical** computers,  
exactly one must output 1



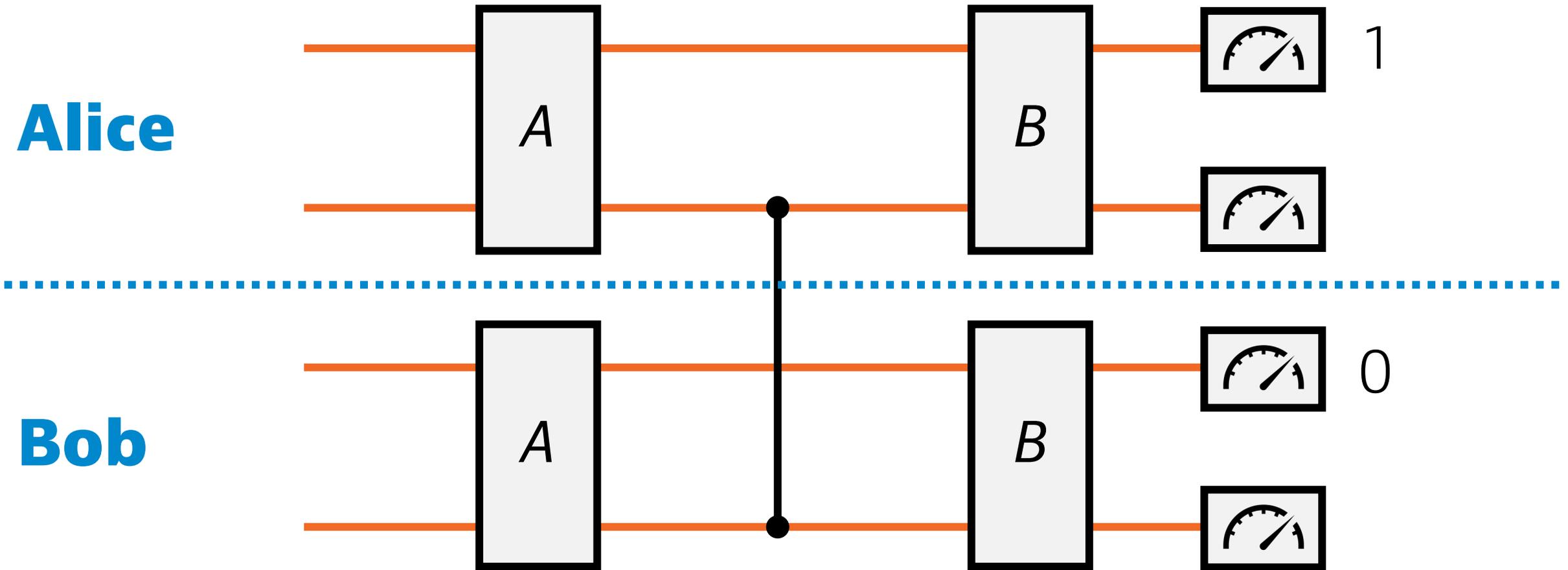
**Classical** algorithm:

- produce 100 random bits
- send to the other party
- larger value  $\rightarrow$  output 1

Non-zero **error probability**,  
unavoidable in classical 1-round  
protocols with finite messages!

Two **identical** computers,  
exactly one must output 1

Succeeds **with probability 1**



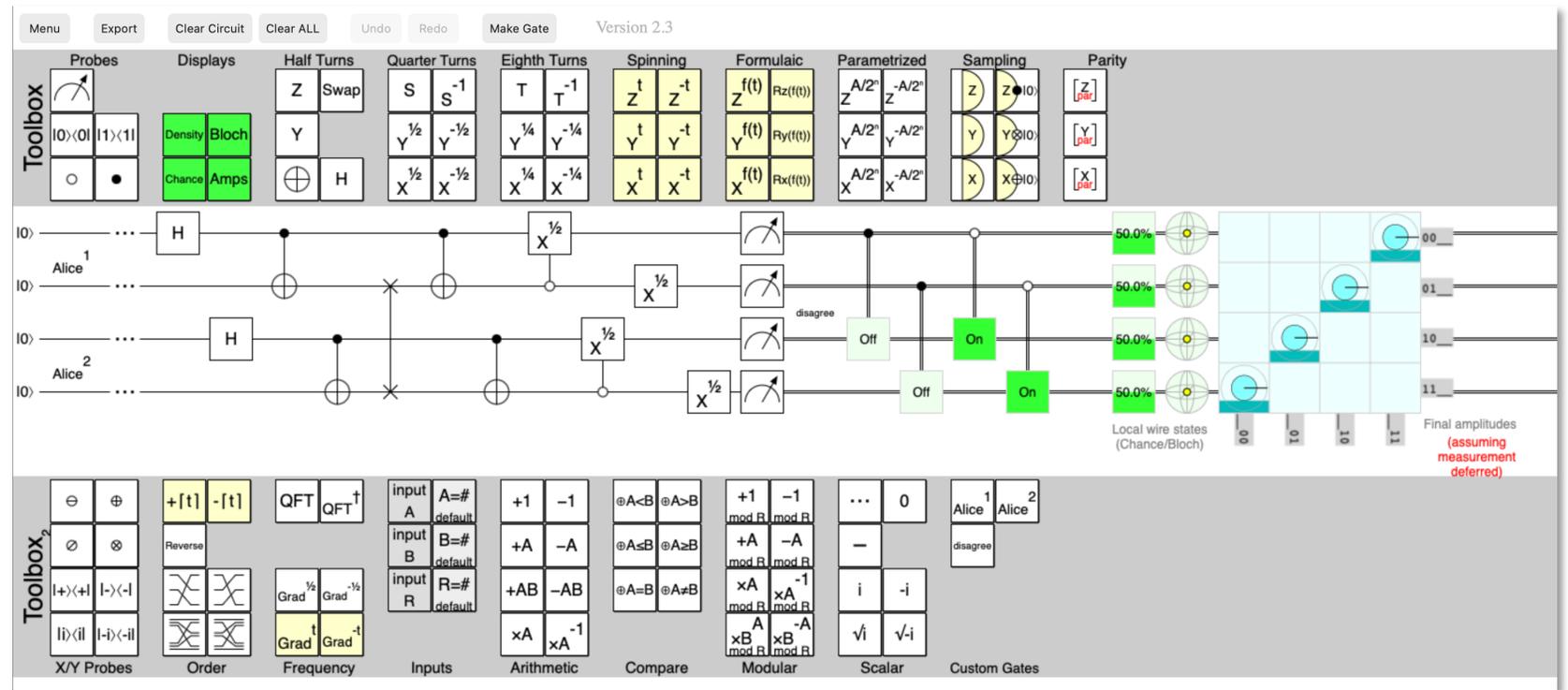
Two **identical** computers,  
exactly one must output 1

Succeeds **with probability 1**

Read more in Craig  
Gidney's 2014 blog  
post and play with  
the circuit in his  
**Quirk** simulator

[algassert.com/quirk](http://algassert.com/quirk)

This is one of the  
sample circuits!



**Confusion:  
4 different  
perspectives**

# Physicists:

testing physics

- *e.g. quantum nonlocality  
(local hidden variables do  
not explain what we see)*

## **Physicists:**

testing physics

- *e.g. quantum nonlocality  
(local hidden variables do  
not explain what we see)*

## **Crypto:**

quantum key distribution

- *classical authenticated  
channel + quantum  
→ confidential channel*

## **Physicists:**

testing physics

- *e.g. quantum nonlocality  
(local hidden variables do  
not explain what we see)*

## **Engineers:**

building big quantum  
computers is hard

- *run quantum algorithms in  
a network of many small  
quantum computers*

## **Crypto:**

quantum key distribution

- *classical authenticated  
channel + quantum  
→ confidential channel*

## **Physicists:**

testing physics

- *e.g. quantum nonlocality  
(local hidden variables do  
not explain what we see)*

## **Engineers:**

building big quantum computers is hard

- *run quantum algorithms in  
a network of many small  
quantum computers*

## **Crypto:**

quantum key distribution

- *classical authenticated  
channel + quantum  
→ confidential channel*

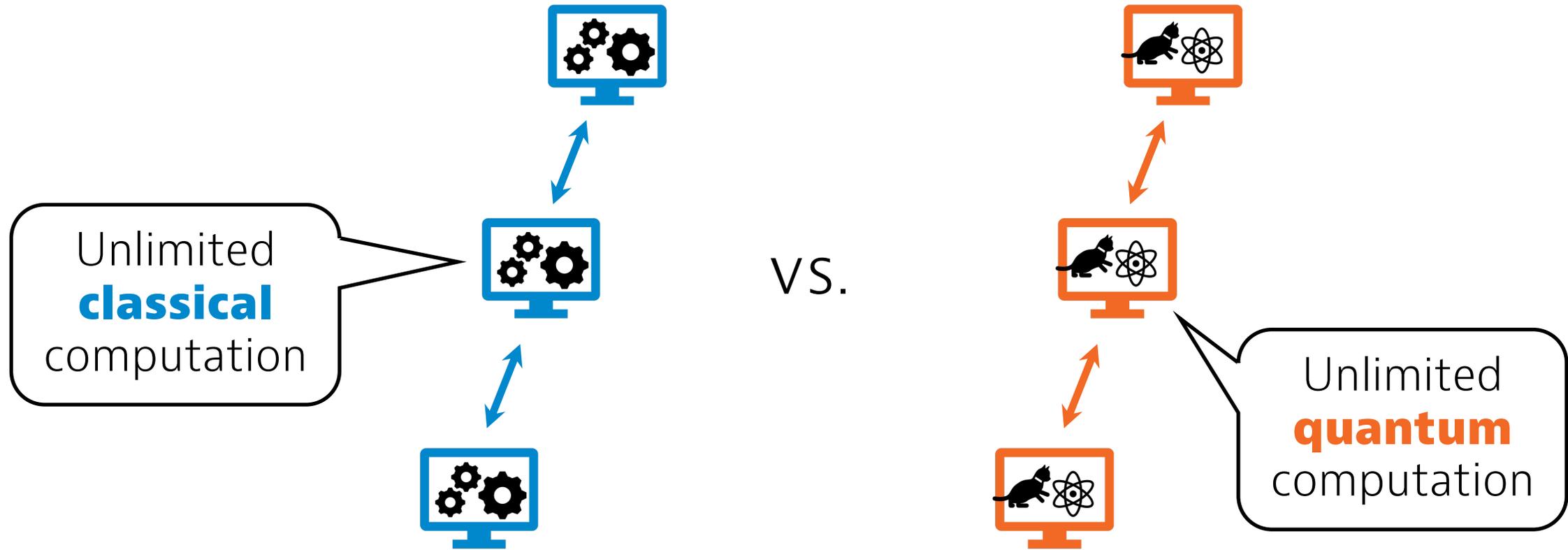
## **Distributed**

**computing theory:**

fundamental

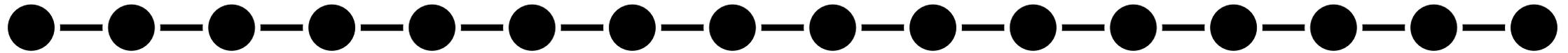
distributed quantum

advantage?



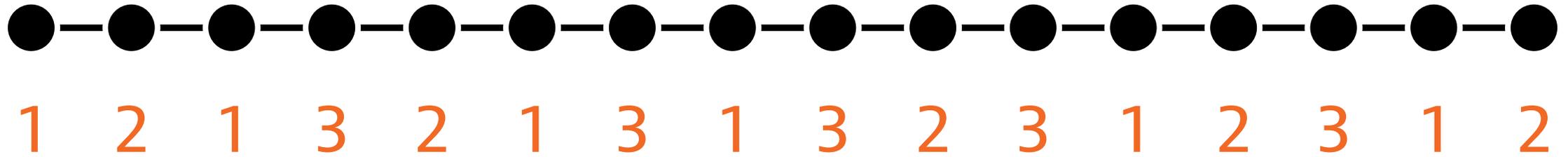
How many **communication rounds** are needed until all computers stop and announce their local outputs?

**Toy examples**

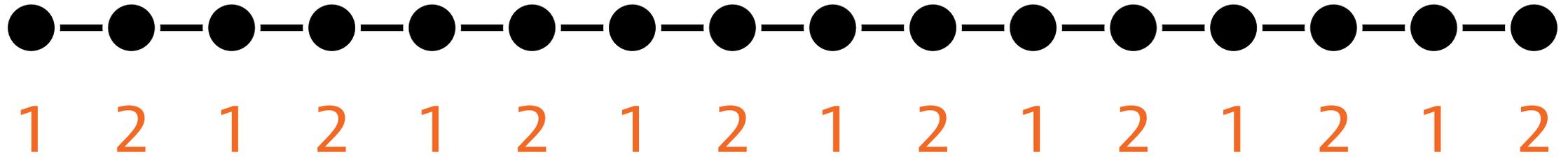


## **Computer network: path**

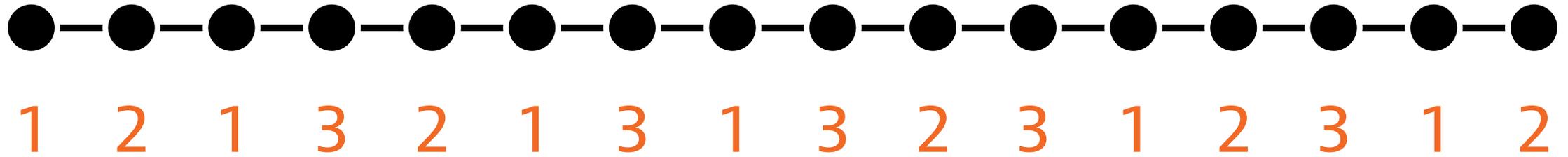
*n identical computers*



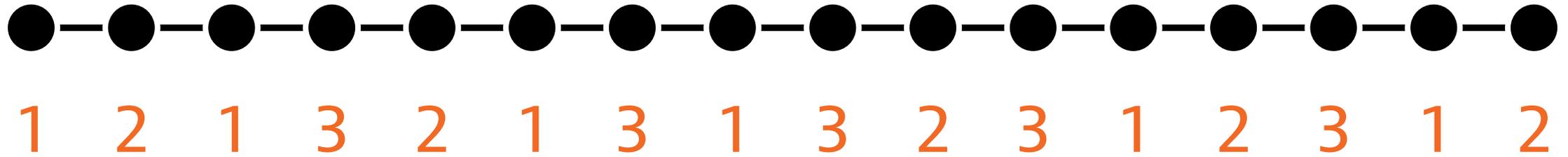
**Coloring:** *eventually each node has to stop and output its own color*



**2-coloring:** *fundamentally global:  
cannot solve in  $o(n)$  rounds*



**3-coloring:** *much easier!*

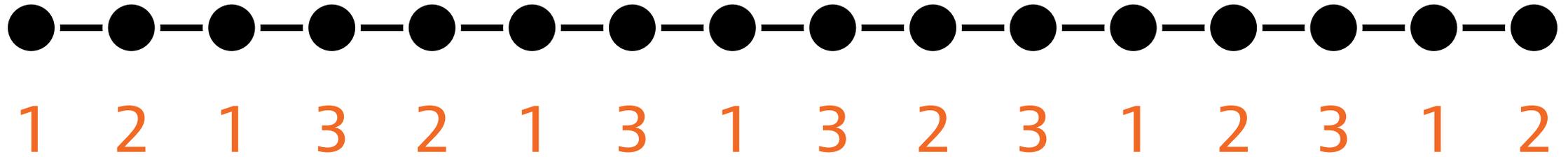


**3-coloring:** *can be solved in  $O(\log^* n)$  rounds,  
and this is tight*

Cole, Vishkin (1986),  
Linial (1992), Naor (1991)

≈ inverse of power tower

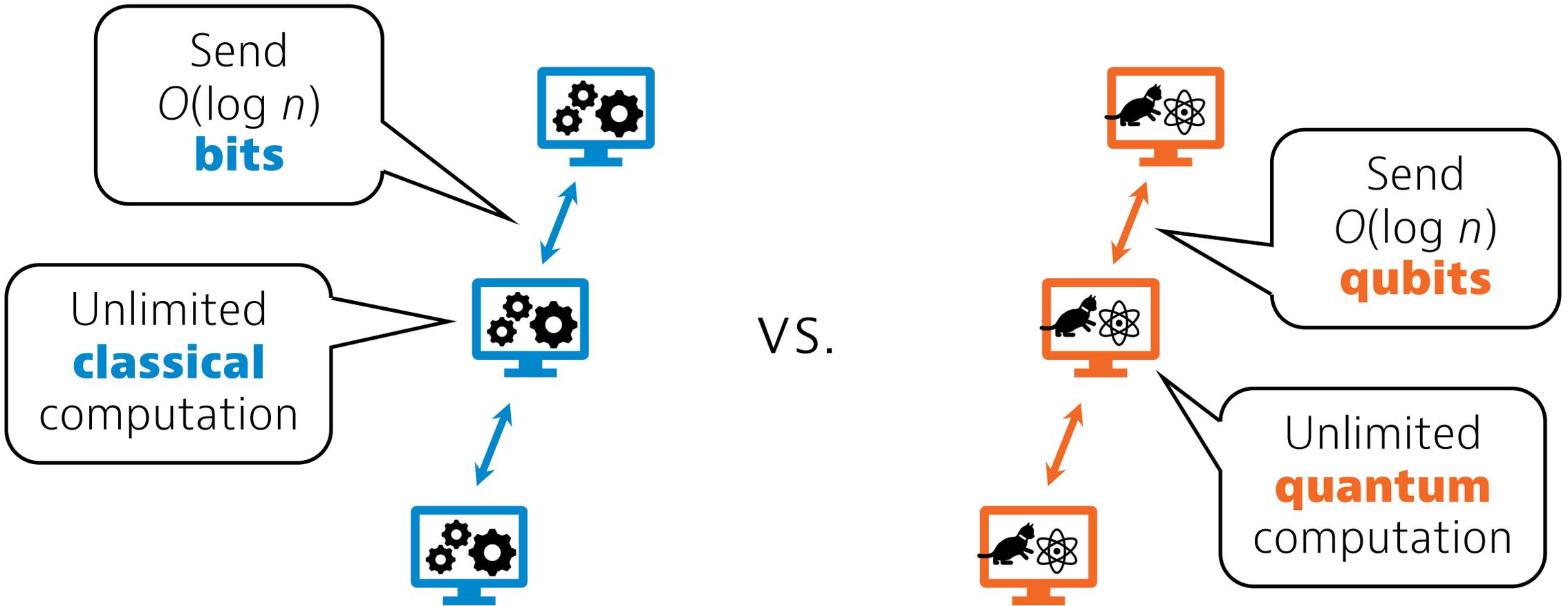
$\log^* n \leq 5$  for pretty much any  
value of  $n$  you will ever encounter



**3-coloring:** *can be solved in  $O(\log^* n)$  rounds, and this is tight for **classical** algorithms*

Does **quantum** help?

**CONGEST**



**CONGEST** vs. **quantum-CONGEST**:  
bandwidth-limited setting

# Example: diameter

- E.g. is the diameter 10 or 11?
- Quick and easy without bandwidth constraints:
  - build shortest-path trees in parallel, starting at all nodes
- Slow but correct: check all nodes one by one
  - this is close to optimal for classical CONGEST

Frischknecht, Holzer, Wattenhofer (2012)

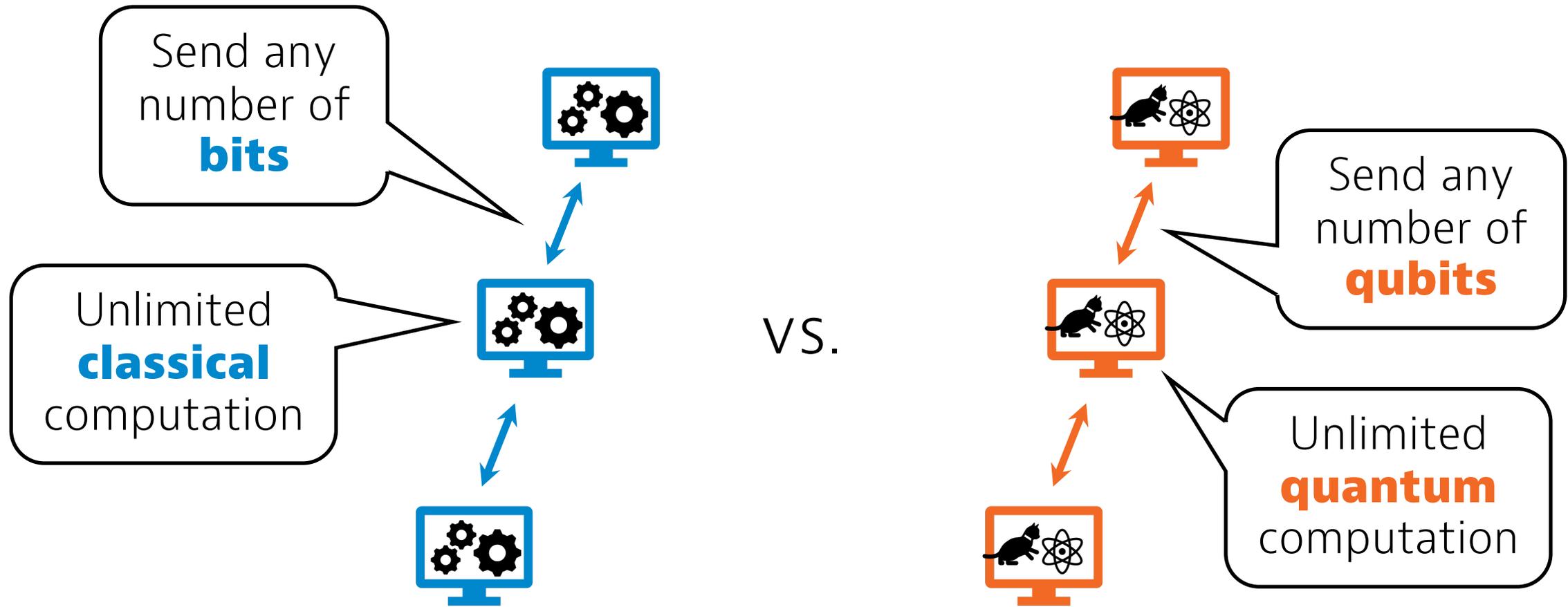
# Example: diameter

- **Centralized Grover search**
  - check  $n$  candidates in  $\approx n^{1/2}$  time
- **Distributed Grover search**
  - can replace  $n$  classical search operations with  $n^{1/2}$  quantum search operations
  - here “search operation” for node  $v$ : can we find a shortest path of length 11 starting from  $v$ ?

# Limitations

- Distributed Grover search only gives moderate savings in bandwidth
- E.g. diameter in  $\tilde{O}(n^{1/2})$  rounds:
  - distributed Grover search:  $O(\log n)$  qubits/message
  - trivial classical algorithm:  $O(n^{1/2})$  bits/message
- If we have infrastructure for sending qubits, can't we also send large classical messages?

**LOCAL**



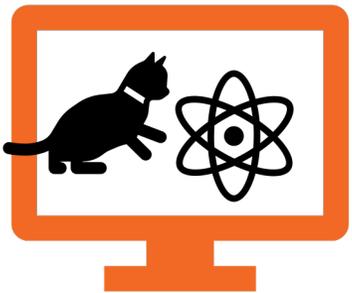
**LOCAL** vs. **quantum-LOCAL**:  
advantage beyond bandwidth savings?

Le Gall, Nishimura,  
Rosmanis (2019)

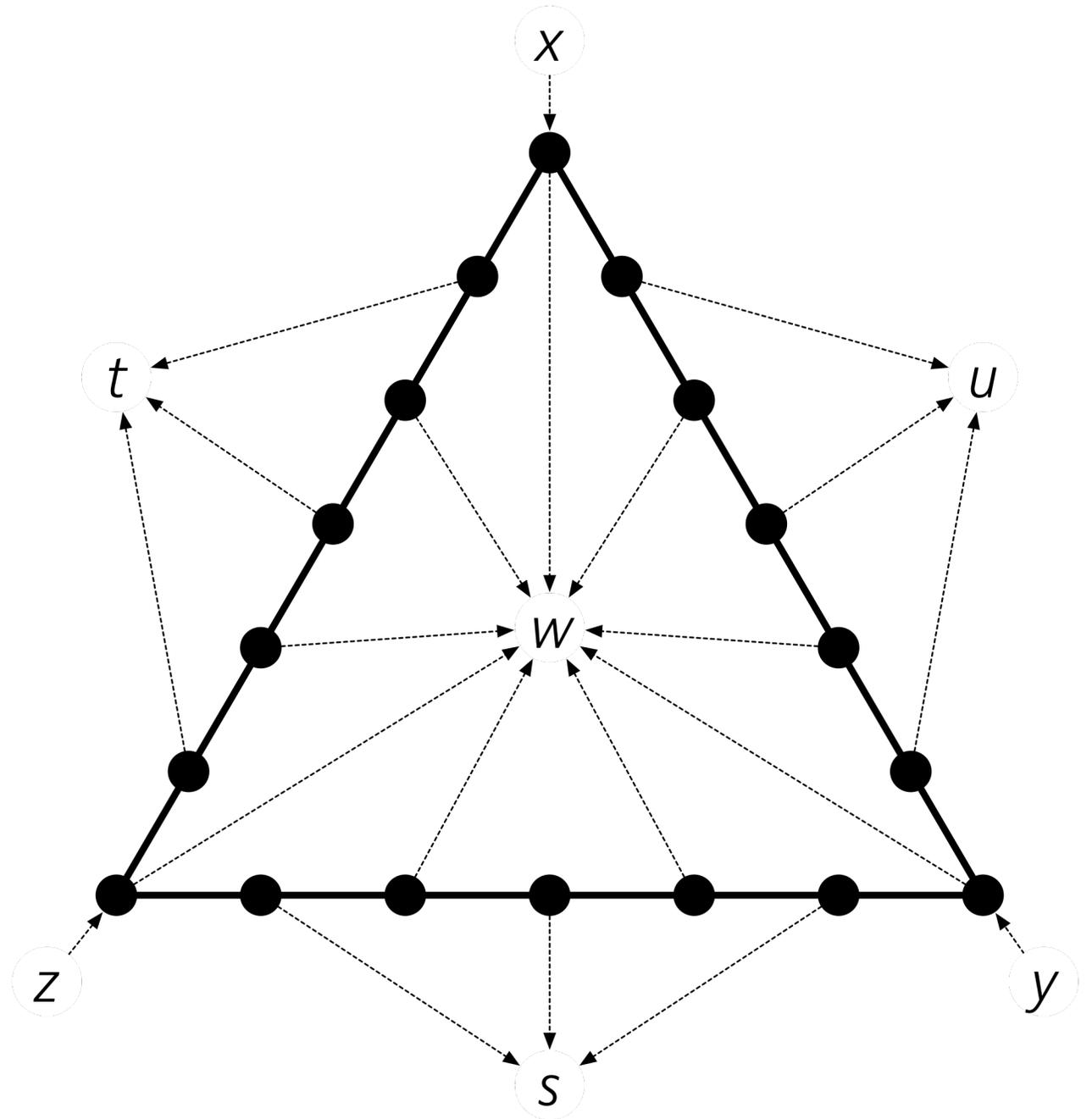
# Quantum Advantage for the LOCAL Model in Distributed Computing



$n/6$  rounds

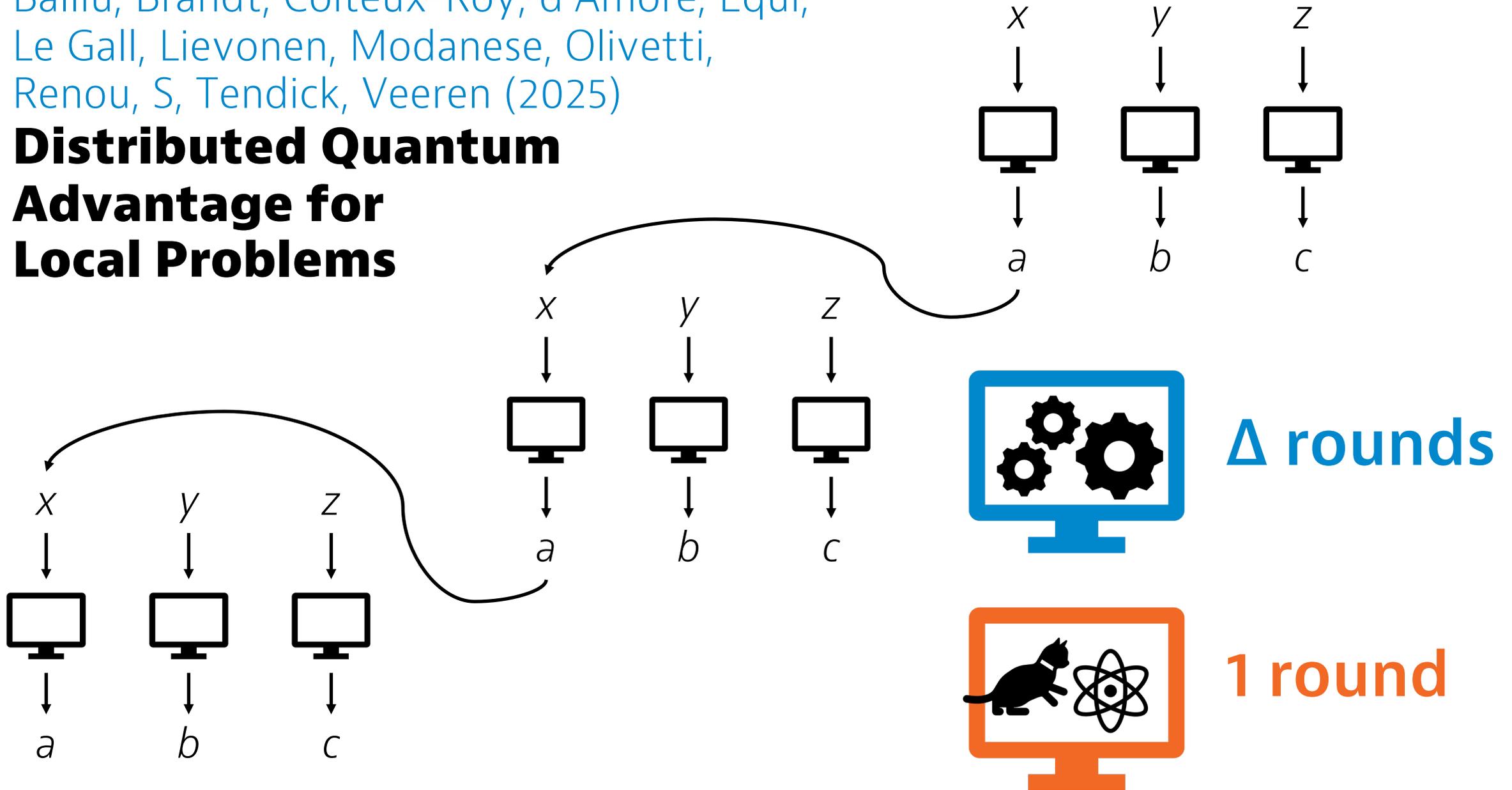


2 rounds



Balliu, Brandt, Coiteux-Roy, d'Amore, Equi,  
Le Gall, Lievonen, Modanese, Olivetti,  
Renou, S, Tendick, Veeren (2025)

## Distributed Quantum Advantage for Local Problems

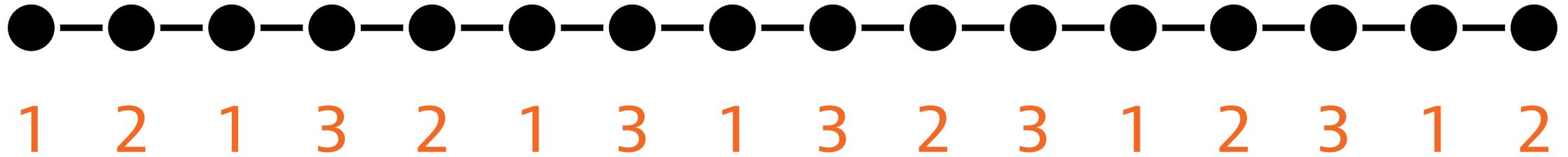


# Quantum-LOCAL

- We have **artificial** problems that separate LOCAL and quantum-LOCAL
- Does quantum help with any **practically relevant** problem?

# Quantum-LOCAL

- We have **artificial** problems that separate LOCAL and quantum-LOCAL
- Does quantum help with any **practically relevant** problem?
- ***We do not know!***



**3-coloring:** *can be solved in  $O(\log^* n)$  rounds, and this is tight for **classical** algorithms*

Does **quantum** help? **We do not know!**

# Obstacles

Standard technique for showing lack of quantum advantage:

**LOCAL**  $\leq$  **quantum-LOCAL**  $\leq$  **non-signaling**

# Obstacles

Standard technique for showing lack of quantum advantage:

**LOCAL**  $\leq$  **quantum-LOCAL**  $\leq$  **non-signaling**

Prove an upper bound here...

... and a matching lower bound here

# Obstacles

Standard technique for showing lack of quantum advantage:

**LOCAL**  $\leq$  **quantum-LOCAL**  $\leq$  **non-signaling**

Prove an upper bound here...

No need to touch quantum!

... and a matching lower bound here

# Obstacles

Standard technique for showing lack of quantum advantage:

**LOCAL**  $\leq$  **quantum-LOCAL**  $\leq$  **non-signaling**

# Obstacles

Standard technique for showing lack of quantum advantage:

**LOCAL**  $\leq$  **quantum-LOCAL**  $\leq$  **non-signaling**

3-coloring  
not possible in  
 $O(1)$  rounds

3-coloring  
is possible in  
 $O(1)$  "rounds"

# Obstacles

Standard technique for showing lack of quantum advantage:

**LOCAL**  $\leq$  **quantum-LOCAL**  $\leq$  **non-signaling**

3-coloring  
not possible in  
 $O(1)$  rounds

We must get  
our hands dirty  
with quantum

3-coloring  
is possible in  
 $O(1)$  "rounds"

# Summary

# Distributed quantum advantage

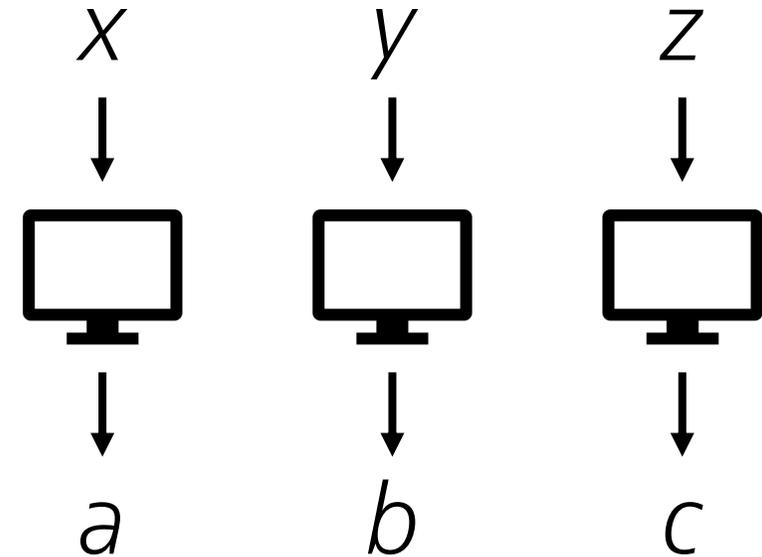
- **Quantum computer network:** *"quantum circuit stretched between multiple sites"*
- **CONGEST:** advantage for "interesting" tasks
- **LOCAL:** advantage for artificial tasks so far
- **Open:** 3-coloring paths in  $O(1)$  rounds?

**Bonus**

# GHZ game

Greenberger,  
Horne, Zeilinger

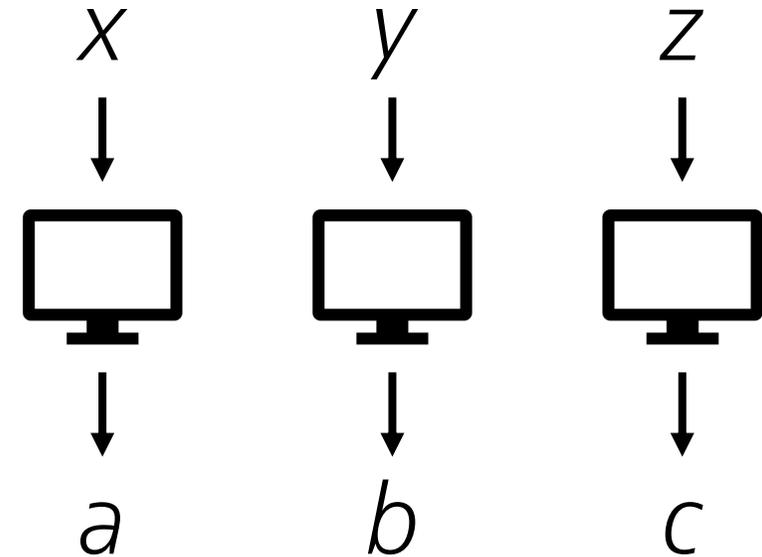
| $x + y + z$ | $a + b + c$<br>mod 2 |
|-------------|----------------------|
| 0           | 0                    |
| 1           | (forbidden)          |
| 2           | 1                    |
| 3           | (forbidden)          |



# GHZ game

Greenberger,  
Horne, Zeilinger

| $x + y + z$ | $a + b + c$<br>mod 2 |
|-------------|----------------------|
| 0           | 0                    |
| 1           | 0 or 1               |
| 2           | 1                    |
| 3           | 0 or 1               |



| $x + y + z$ | $a + b + c \pmod 2$ |
|-------------|---------------------|
| 0           | 0                   |
| 1           | 0 or 1              |
| 2           | 1                   |
| 3           | 0 or 1              |

