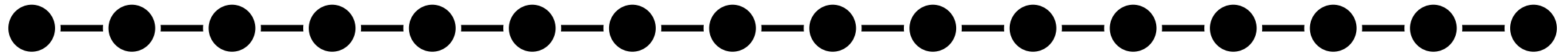


Jukka Suomela

Aalto University

Limits of Distributed Computing

Computer network



Computer network



- exchange information with your neighbors

Computer network



- exchange information with your neighbors
- update your own state

Computer network



- exchange information with your neighbors
- update your own state
- exchange information with your neighbors

Computer network



- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state

Computer network



- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- exchange information with your neighbors

Computer network

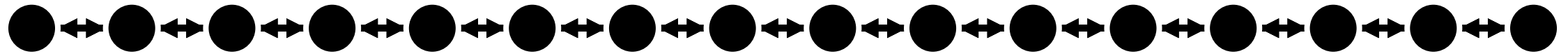


- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state

Computer network



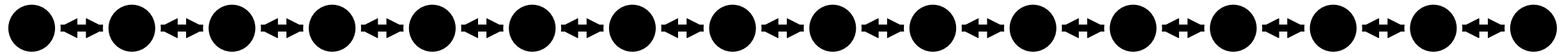
- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- produce your final output (e.g. color)



- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- produce your final output (e.g. color)

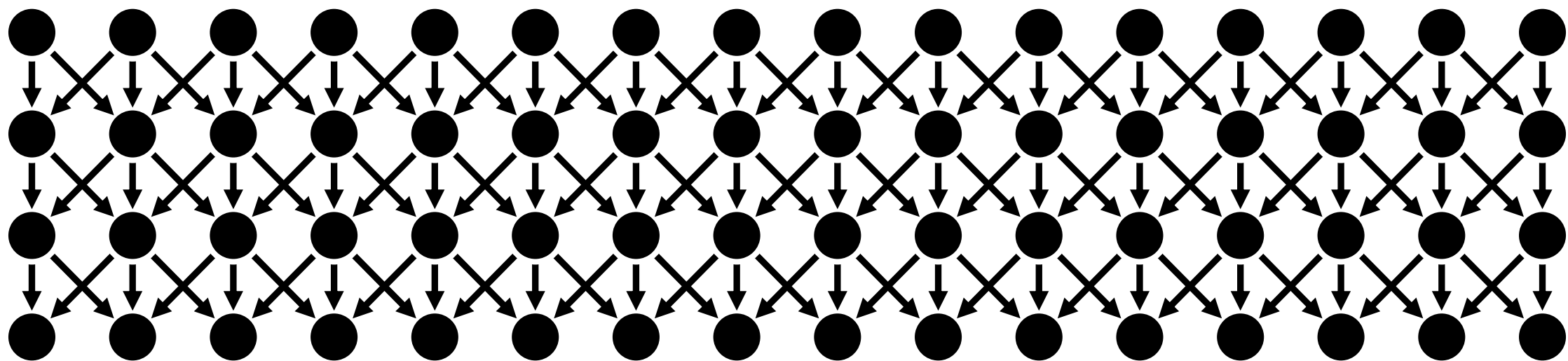
**$T = 3$
rounds**

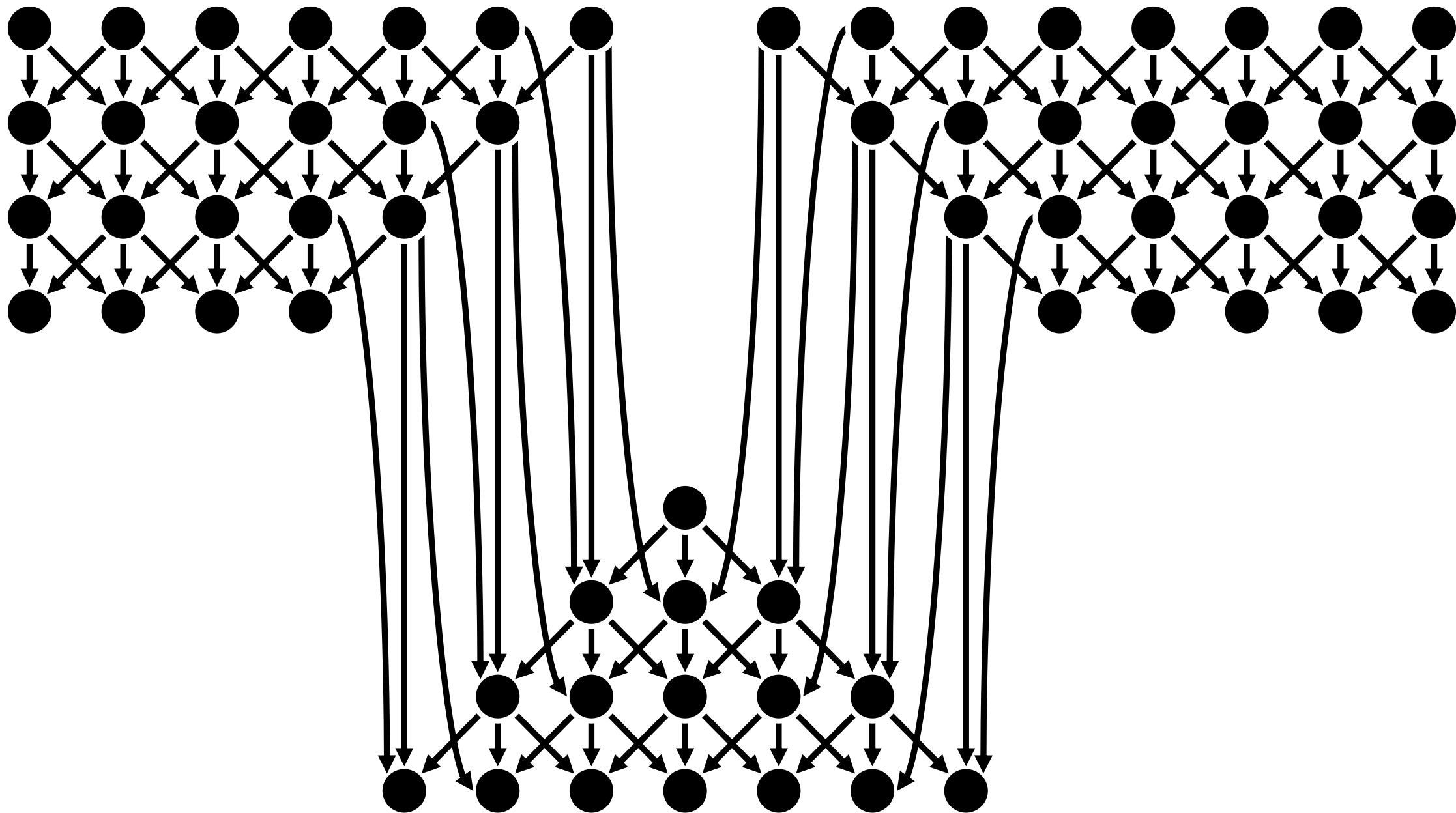
What can we say about final outputs after T rounds?

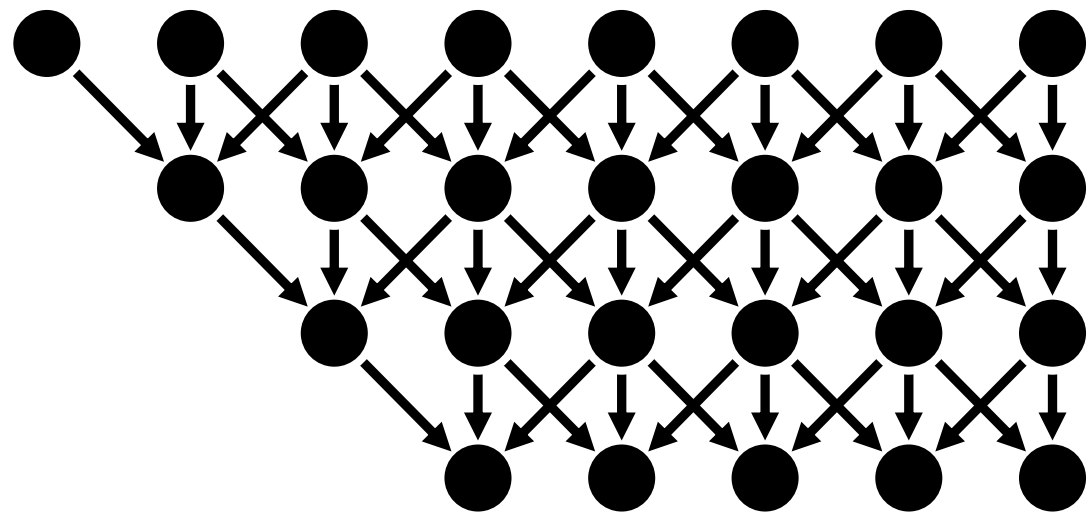
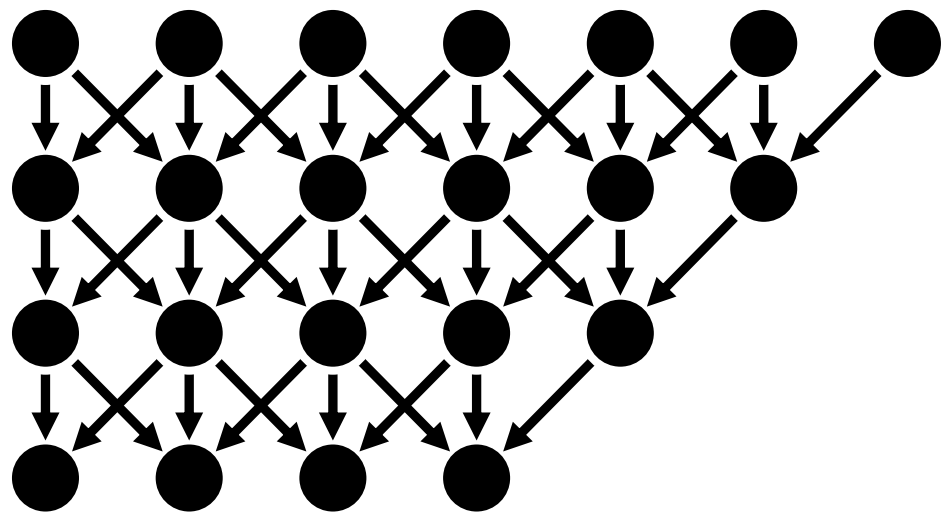


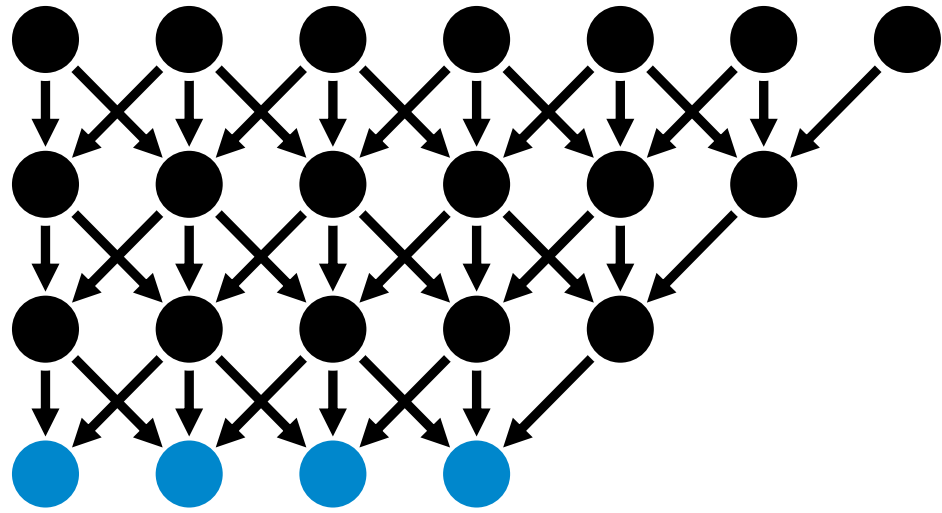
- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- exchange information with your neighbors
- update your own state
- produce your final output (e.g. color)

**$T = 3$
rounds**

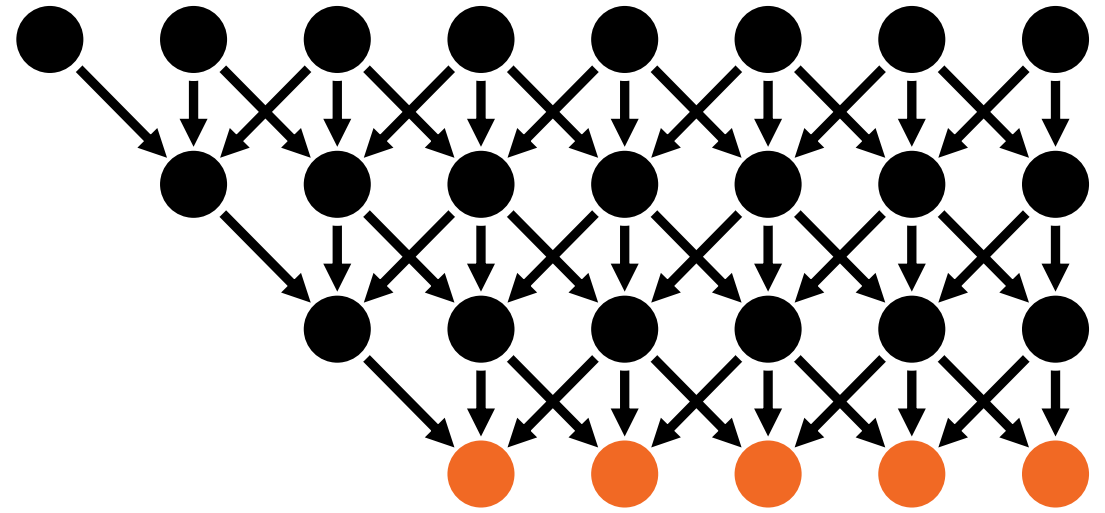








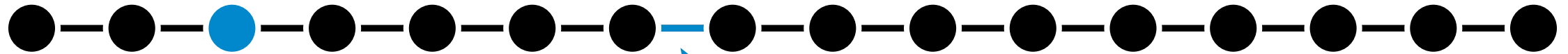
Whatever
these nodes
output ...



... is independent
of whatever
these nodes
output

**What if it's
quantum?**

Network of **quantum** computers



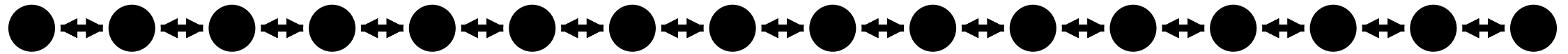
**quantum
computer**

any number
of qubits

**quantum
communication
channel**

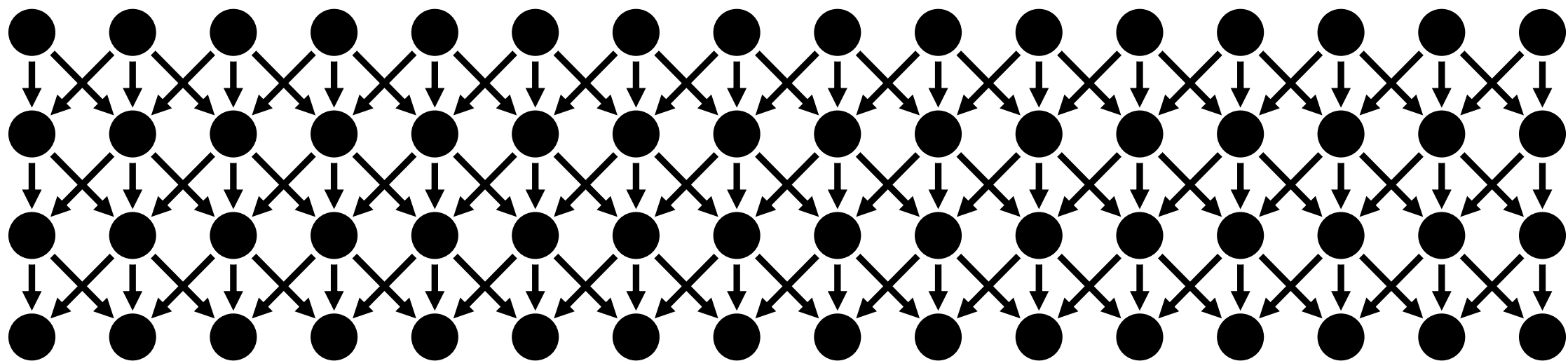
any number
of qubits

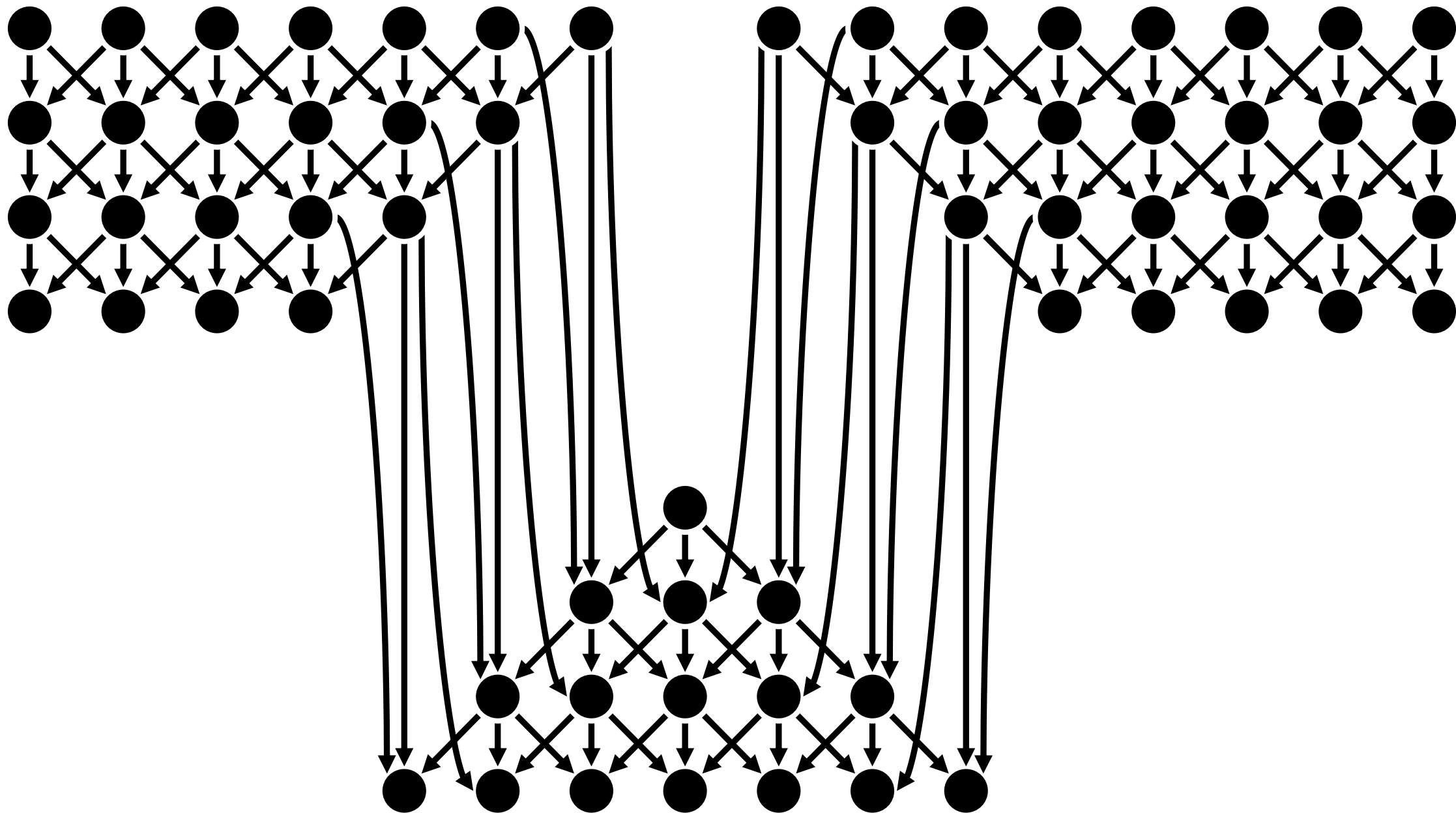
What can we say about final outputs after T rounds?

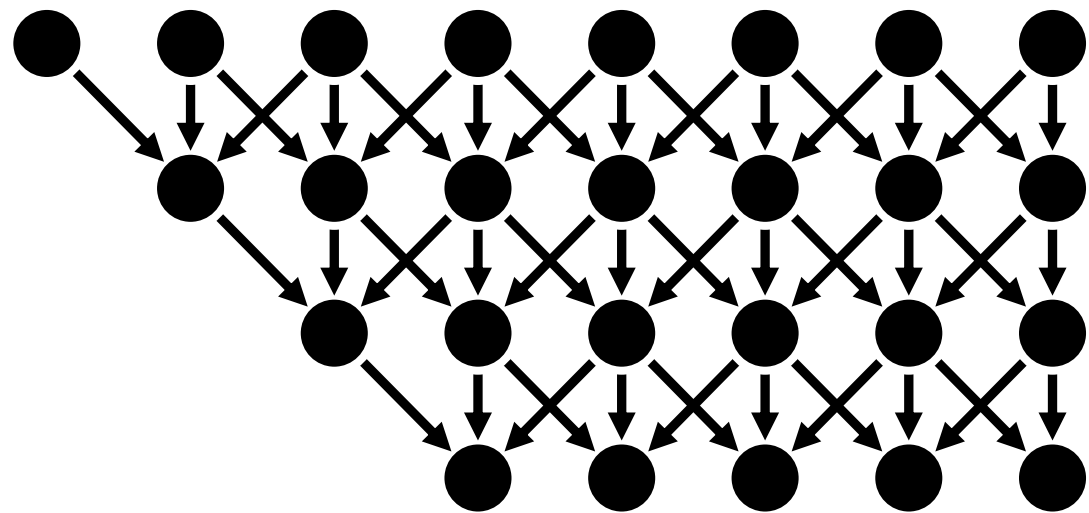
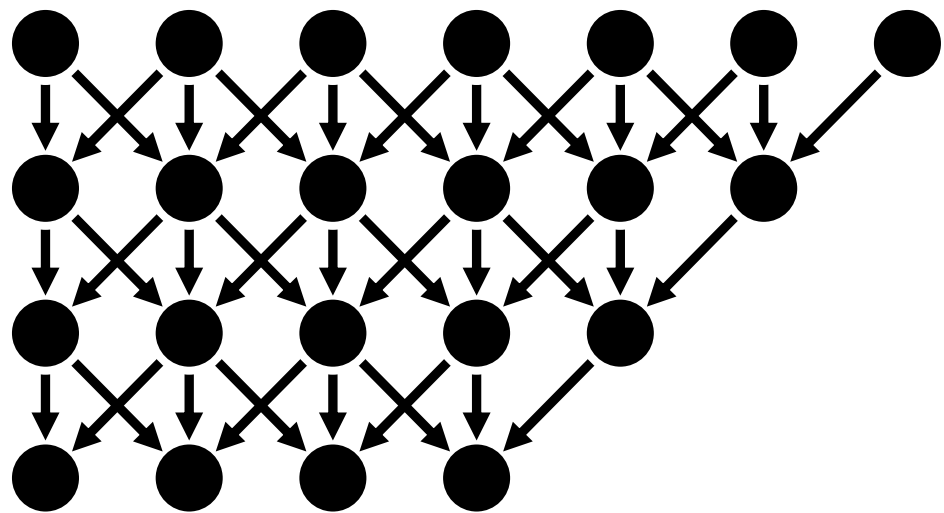


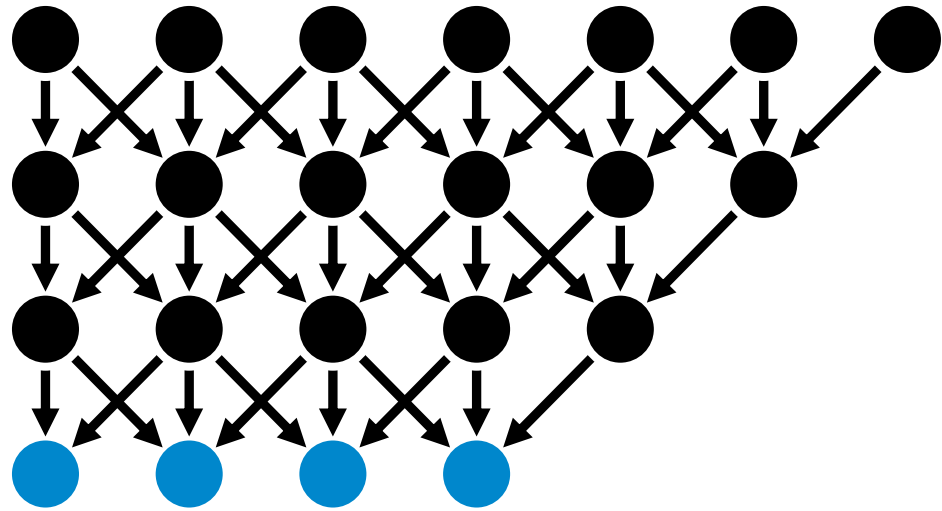
- exchange **qubits** with your neighbors
- update your own **quantum** state
- exchange **qubits** with your neighbors
- update your own **quantum** state
- exchange **qubits** with your neighbors
- update your own **quantum** state
- produce your final output (e.g. color)

$T = 3$
rounds

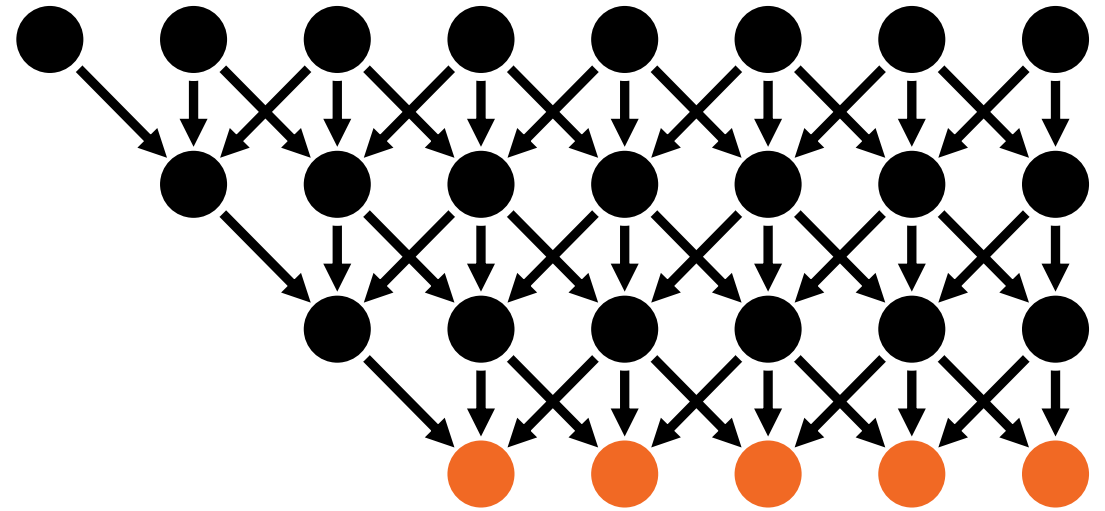






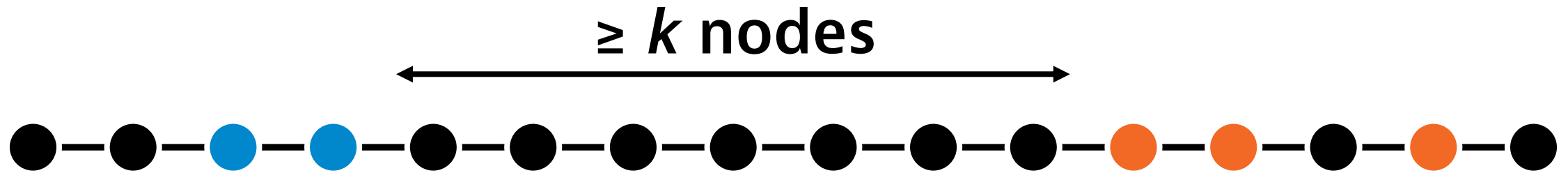


Whatever
these nodes
output ...



... is independent
of whatever
these nodes
output

k -dependent distribution:



These random variables...

... are independent of these random variables

Finely dependent distribution:

k -dependent for some constant k

**Constant-round
distributed algorithm**
(classical or quantum!)



Finely dependent distribution:
 k -dependent for some constant k

**Constant-round
distributed algorithm**
(classical or quantum!)



Show that these
don't exist for
problem X ...

Finely dependent distribution:
 k -dependent for some constant k

**Constant-round
distributed algorithm**
(classical or quantum!)

... and we have
impossibility
here!



Show that these
don't exist for
problem X ...

Finely dependent distribution:
 k -dependent for some constant k

Constant-round distributed algorithm (classical or quantum!)

... and we have impossibility here!

Show that these don't exist for problem X ...

Finely dependent distribution: k -dependent for some constant k

STOC 2024

No Distributed Quantum Advantage for Approximate Graph Coloring

Xavier Coiteux-Roy TU Munich Munich, Germany Munich Center for Quantum Science and Technology Munich, Germany xavier.coiteux-roy@tum.de	Francesco d'Amore Aalto University Espoo, Finland Bocconi University, BIDS Milan, Italy francesco.damore2@unibocconi.it	Rishikesh Gajjala Indian Institute of Science Bangalore, India Aalto University Espoo, Finland rishikeshg@iisc.ac.in
Fabian Kuhn University of Freiburg Freiburg, Germany kuhn@cs.uni-freiburg.de	François Le Gall Nagoya University Nagoya, Japan legall@math.nagoya-u.ac.jp	Henrik Lievenen Aalto University Espoo, Finland henrik.lievenen@aalto.fi
Augusto Modanese Aalto University Espoo, Finland augusto.modanese@aalto.fi	Marc-Olivier Renou Inria Paris, France Université Paris-Saclay Paris, France Institut Polytechnique de Paris Palaiseau, France marc-olivier.renou@inria.fr	Gustav Schmid University of Freiburg Freiburg, Germany schmidg@informatik.uni-freiburg.de
Jukka Suomela Aalto University Espoo, Finland jukka.suomela@aalto.fi		

ABSTRACT
We give an almost complete characterization of the hardness of r -coloring r -chromatic graphs with distributed algorithms, for a wide range of models of distributed computing. In particular, we show that these problems do not admit any distributed quantum advantage. To do that:
(1) We give a new distributed algorithm that finds a coloring in r -chromatic graphs in $O(n^2)$ rounds, with $n = \lfloor \frac{2n}{r} \rfloor$.
(2) We prove that any distributed algorithm for this problem requires $\Omega(n^2)$ rounds.
Our upper bound holds in the classical, deterministic LOCAL model, while the near-matching lower bound holds in the non-signaling model. This model, introduced by Arfaoui and Fraignaud in 2014, captures all models of distributed graph algorithms that obey physical causality; this includes not only classical deterministic LOCAL and randomized LOCAL but also quantum-LOCAL, even with a pre-shared quantum state.
We also show that similar arguments can be used to prove that, e.g., 3-coloring 3-dimensional grids or coloring trees remain hard problems even for the non-signaling model, and in particular do not admit any quantum advantage. Our lower-bound arguments are purely graph-theoretic, at heart, no background on quantum information theory is needed to establish the proofs.

CCS CONCEPTS
• Theory of computation → Distributed computing models; Quantum computation theory; Distributed algorithms.

KEYWORDS
distributed computing, graph coloring, non-signaling model, quantum advantage

ACM Reference Format:
Xavier Coiteux-Roy, Francesco d'Amore, Rishikesh Gajjala, Fabian Kuhn, François Le Gall, Henrik Lievenen, Augusto Modanese, Marc-Olivier Renou, Gustav Schmid, and Jukka Suomela. 2024. No Distributed Quantum Advantage for Approximate Graph Coloring. In *Proceedings of the 58th Annual ACM Symposium on Theory of Computing (STOC '24)*, June 24–28, 2024, Vancouver, BC, Canada. ACM, New York, NY, USA, 10 pages. <https://doi.org/10.1145/3658203.3649679>

This work is licensed under a Creative Commons Attribution 4.0 International License.
STOC '24, June 24–28, 2024, Vancouver, BC, Canada
© 2024 Copyright held by the owner/author(s).
ACM ISBN 978-1-4503-8203-0/24/06.
<https://doi.org/10.1145/3658203.3649679>

1901

**Constant-round
distributed algorithm**
(classical or quantum!)



Finely dependent distribution:
 k -dependent for some constant k

Example: 3-coloring cycles

No constant-round **classical algorithm for 3-colorings in cycles**

[Linial 1992, Naor 1991]

No constant-round **classical algorithm for 3-colorings in cycles**

[Linial 1992, Naor 1991]

Finitely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

No constant-round **classical** algorithm for 3-colorings in cycles

[Linial 1992, Naor 1991]

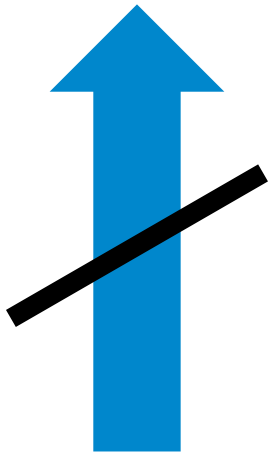


Finely dependent distribution of 3-colorings in cycles

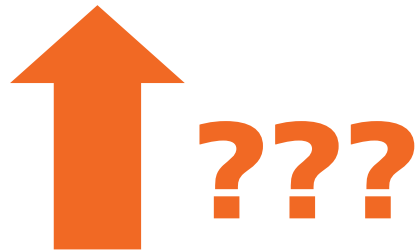
[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

No constant-round **classical** algorithm for 3-colorings in cycles

[Linial 1992, Naor 1991]



Constant-round **quantum** algorithm
for 3-coloring cycles???



Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

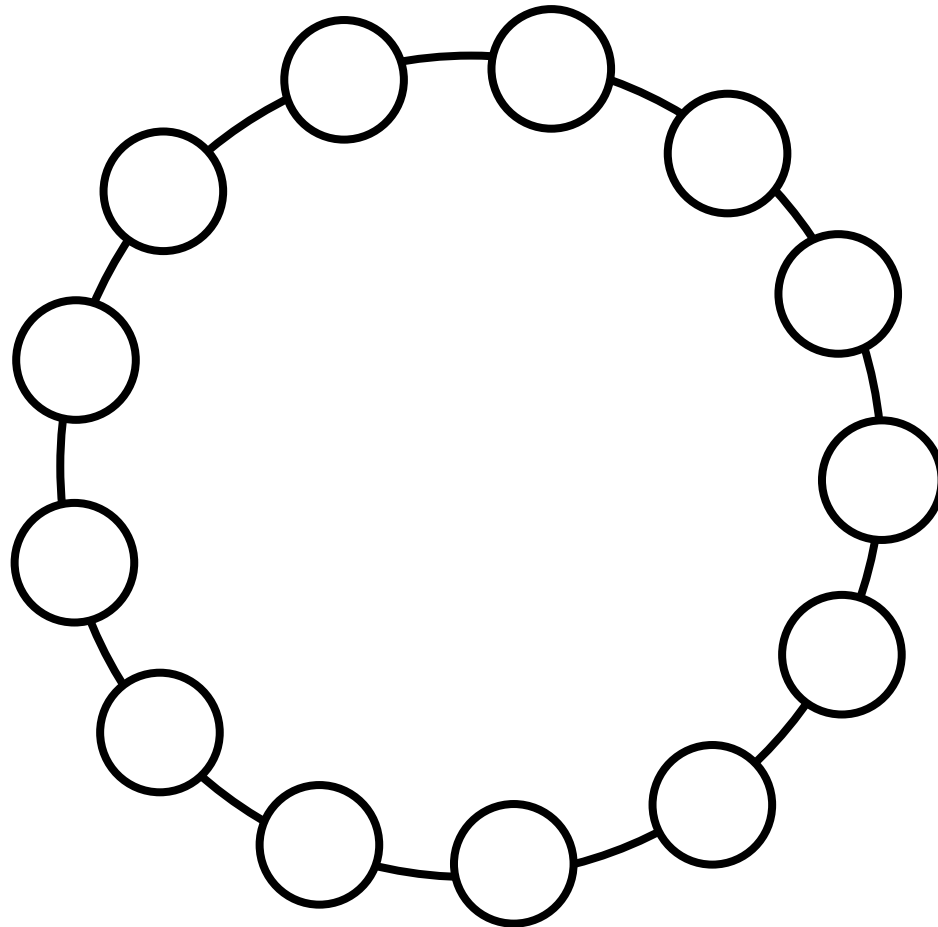
Finitely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Given: n -cycle
(for any $n \geq 3$)

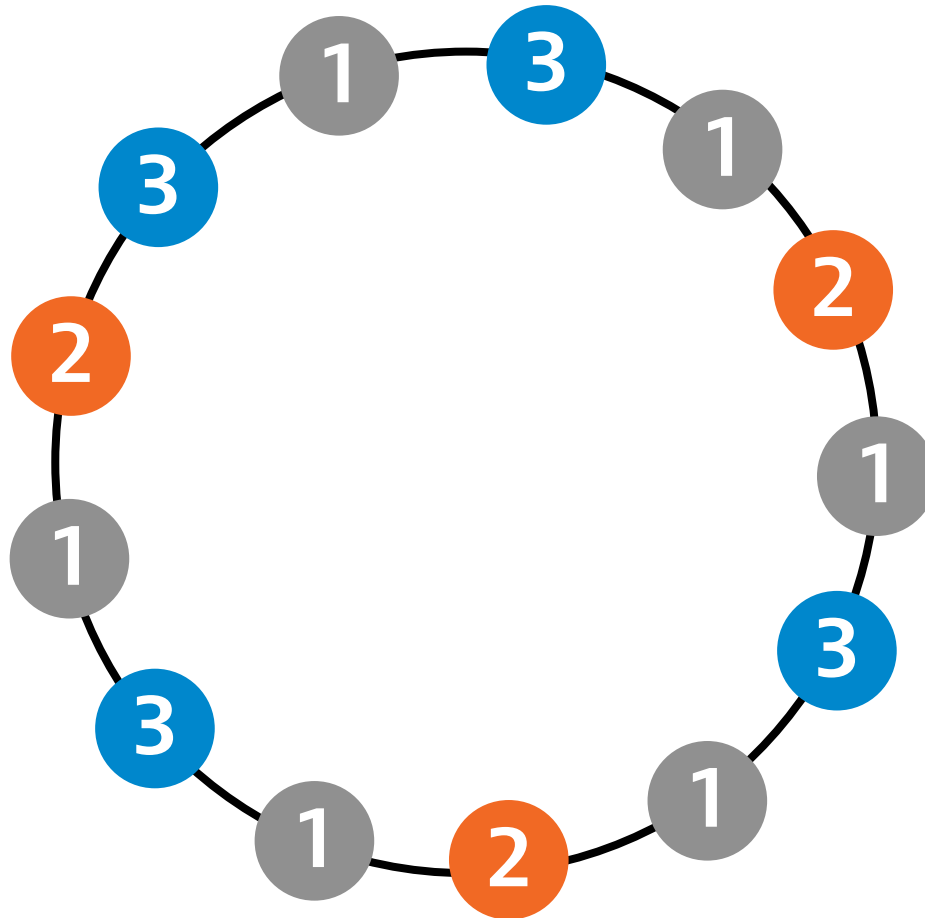


Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Challenge:

design a random process that produces proper 3-colorings that are k -dependent?

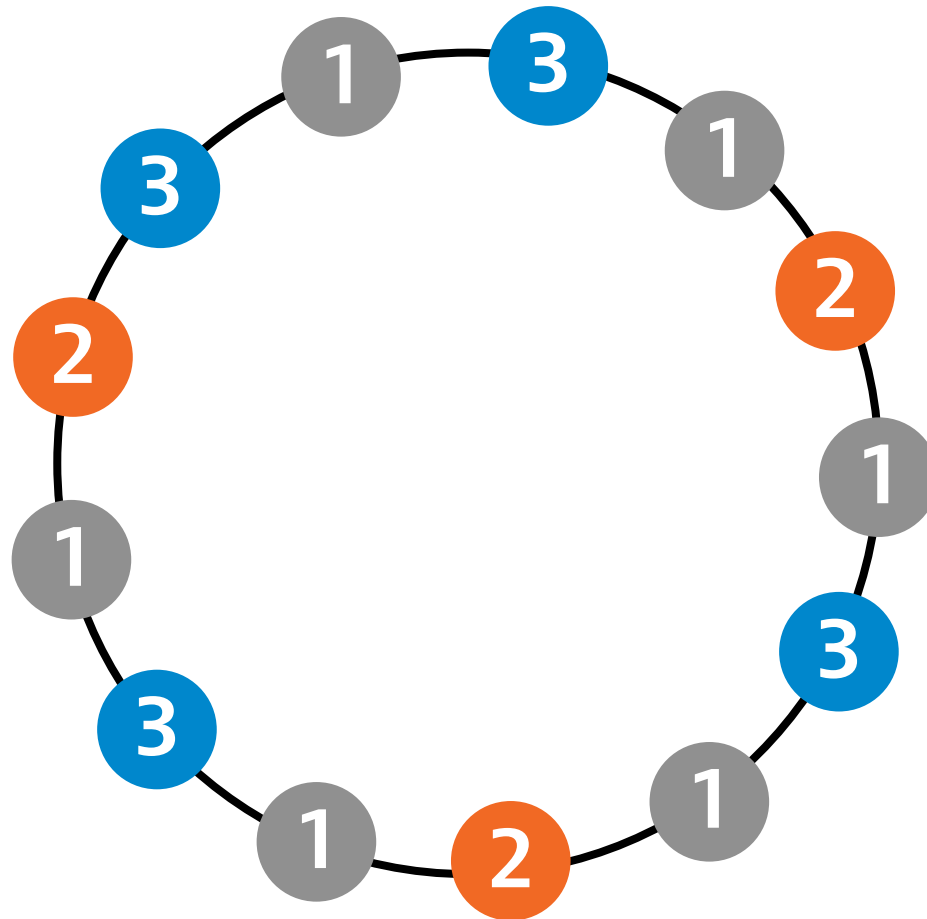


Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Challenge:

design a random process that produces proper 3-colorings that are k -dependent?



Uniformly random colorings do not work!
(exercise: why?)

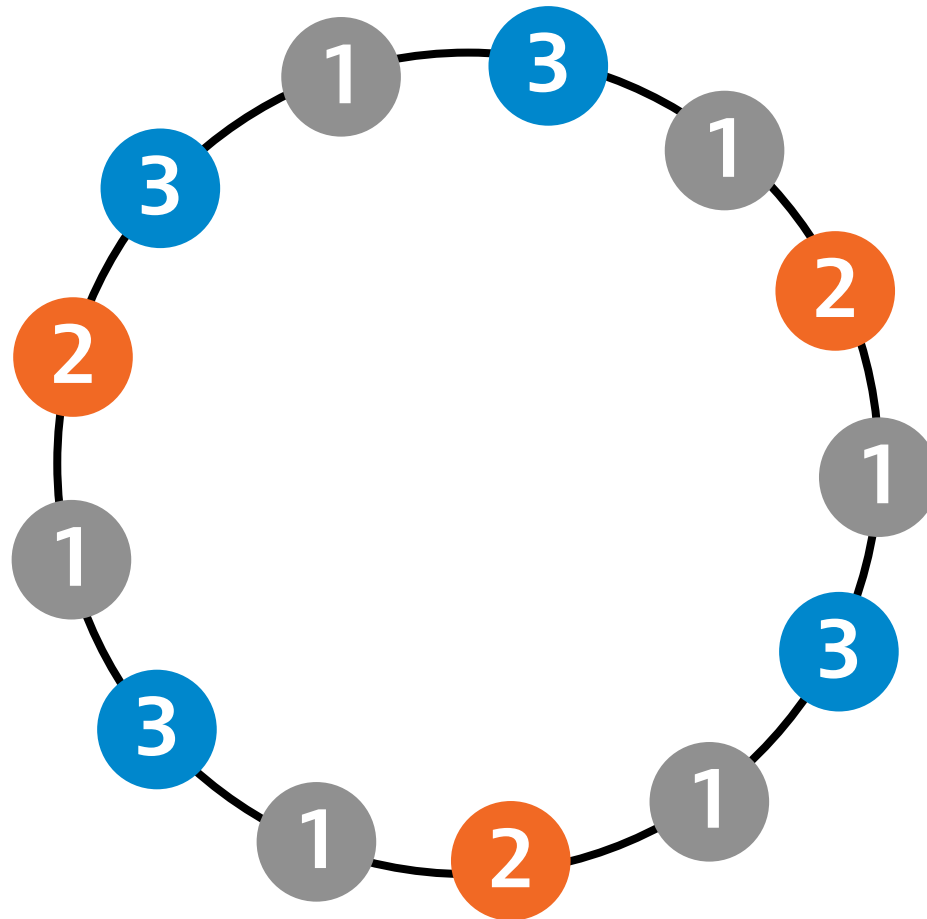
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Challenge:

design a random process that produces proper 3-colorings that are k -dependent?

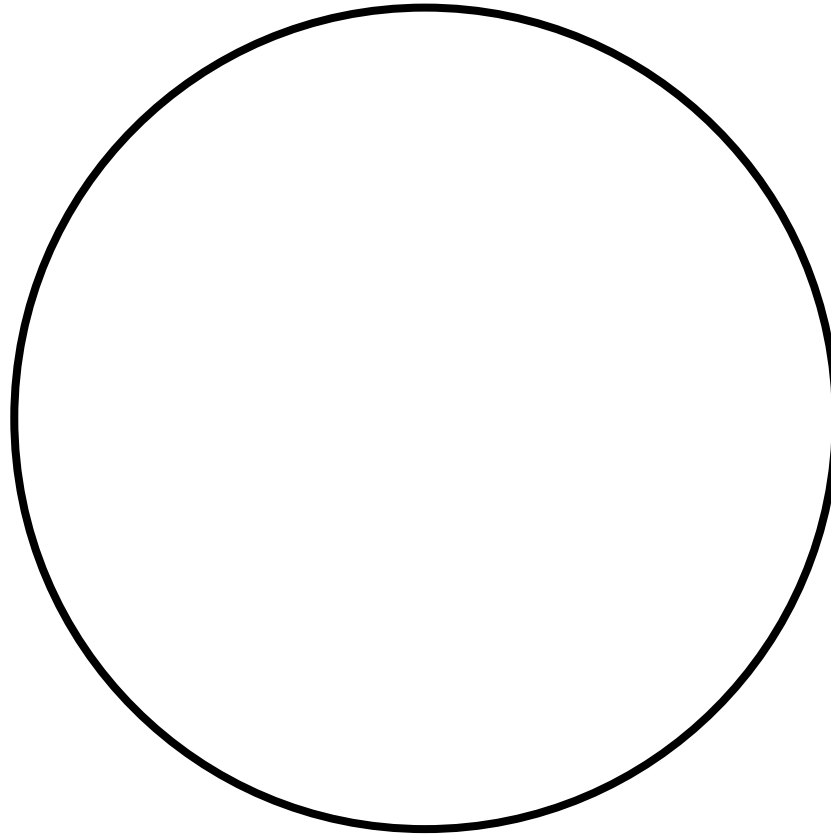
we will have $k = 2$



Uniformly random colorings do not work!
(exercise: why?)

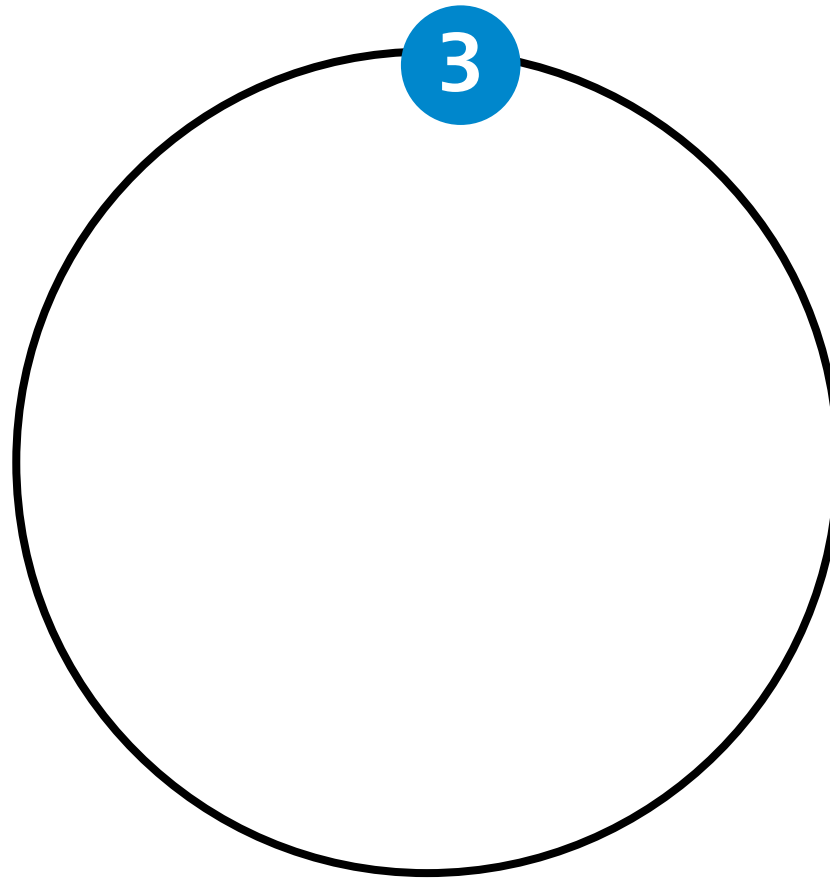
Finitely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



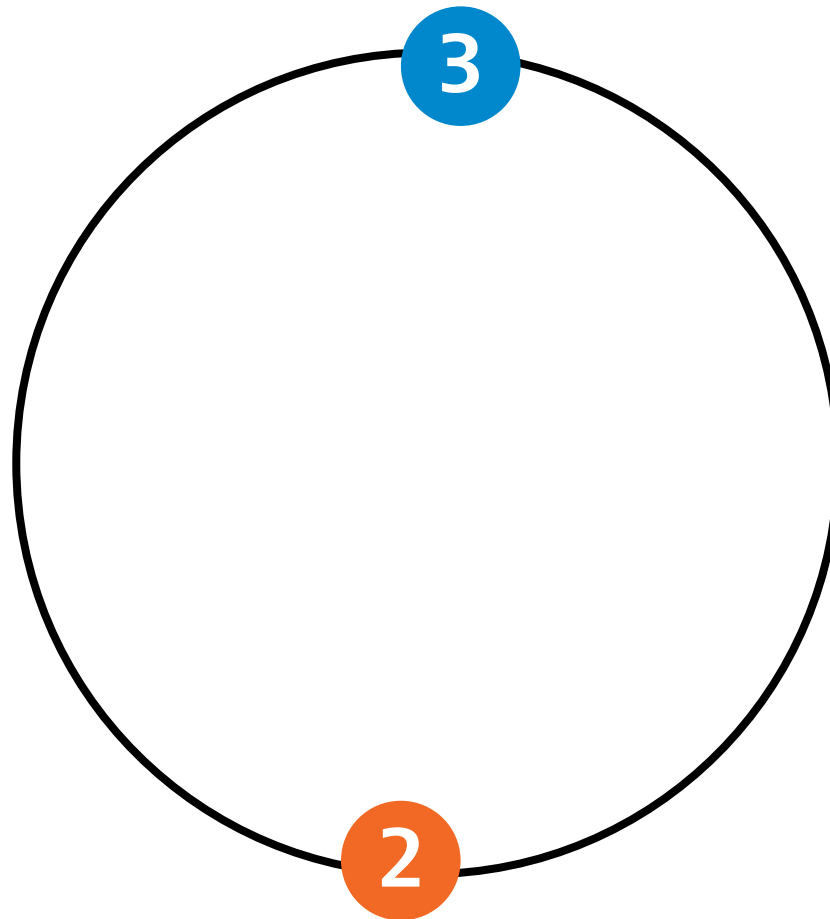
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



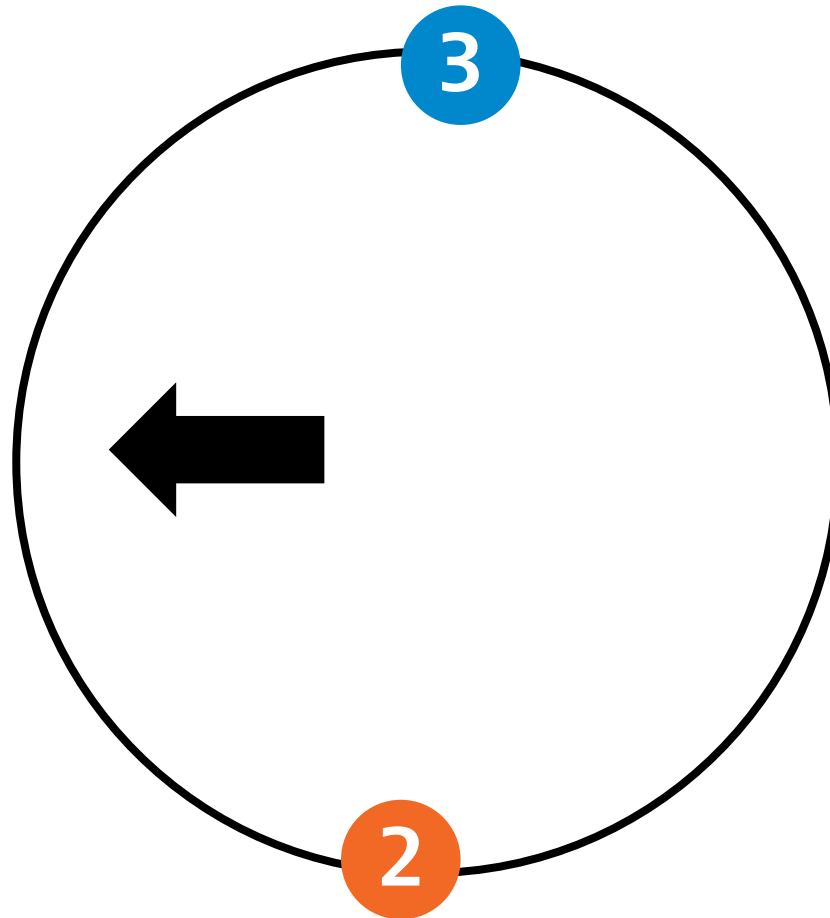
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



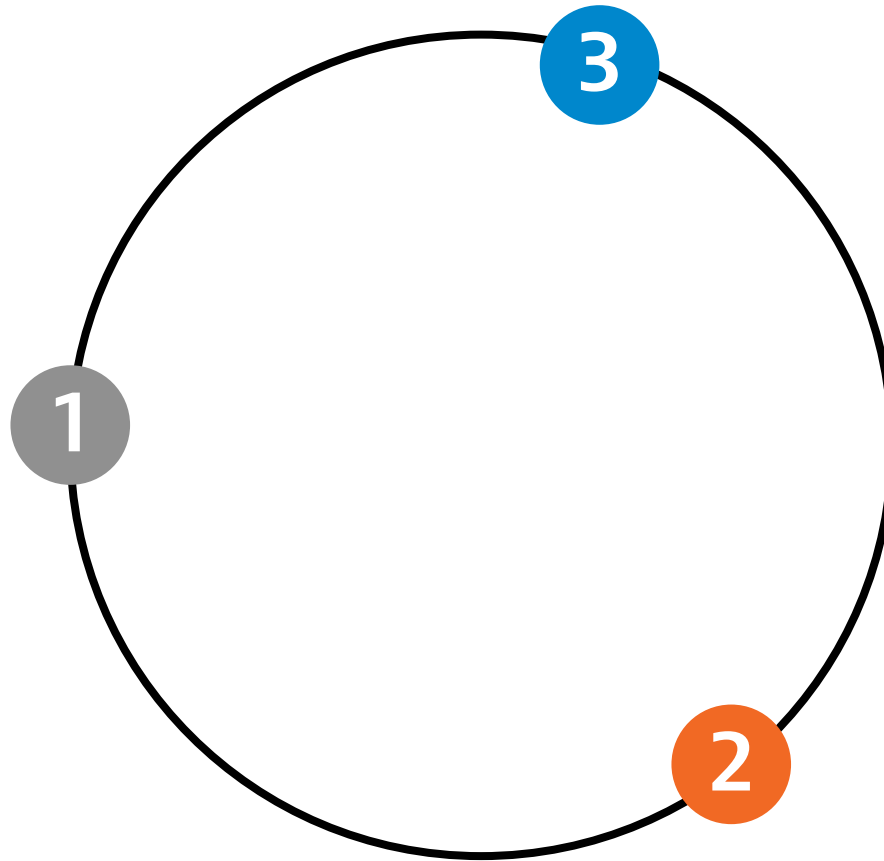
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



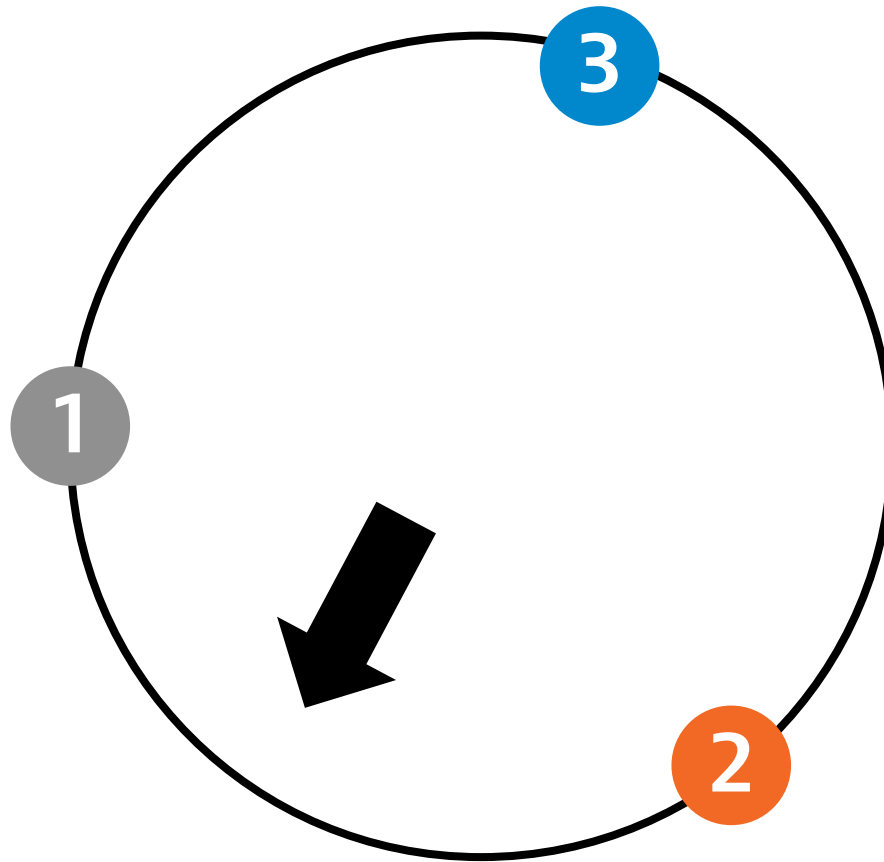
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



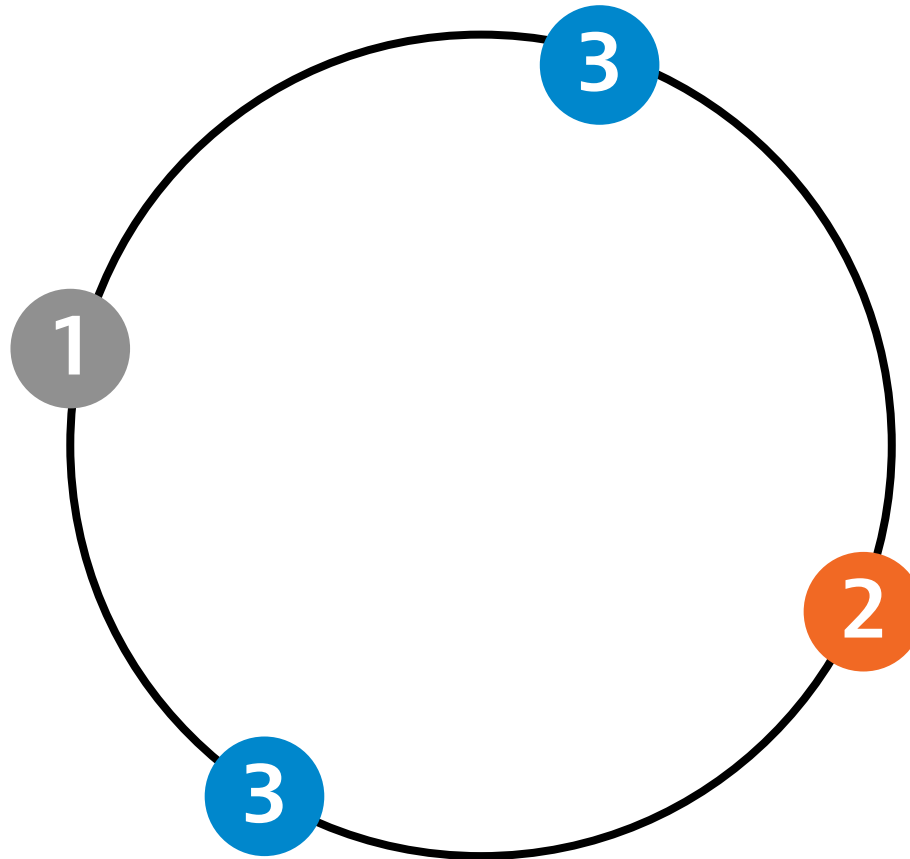
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



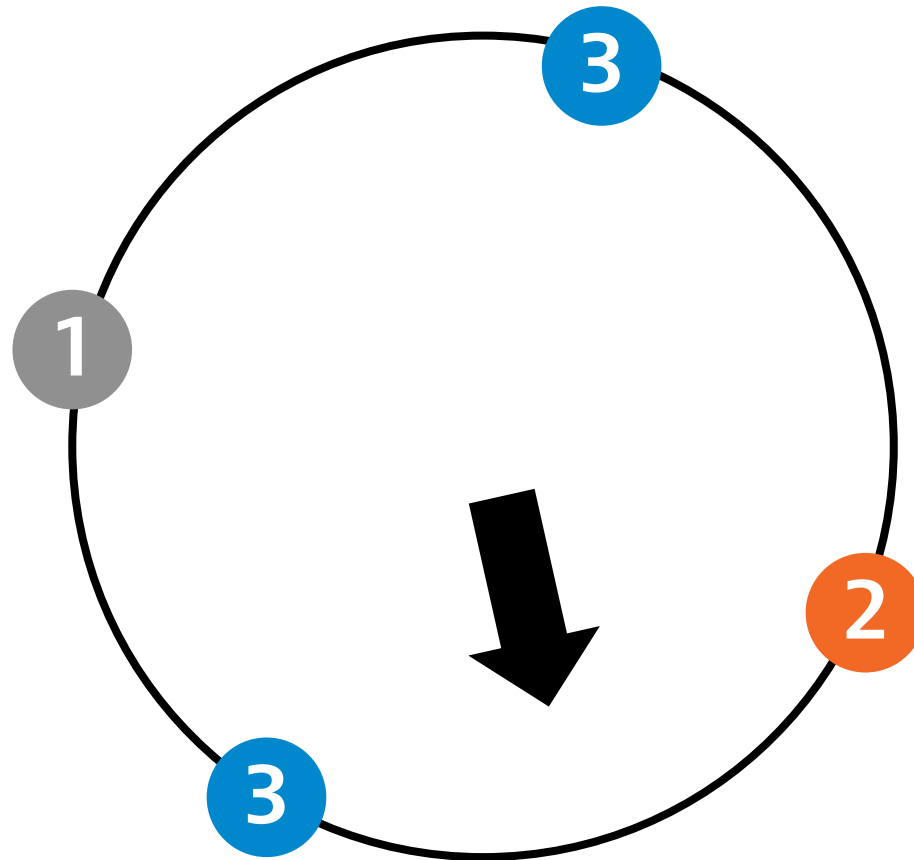
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



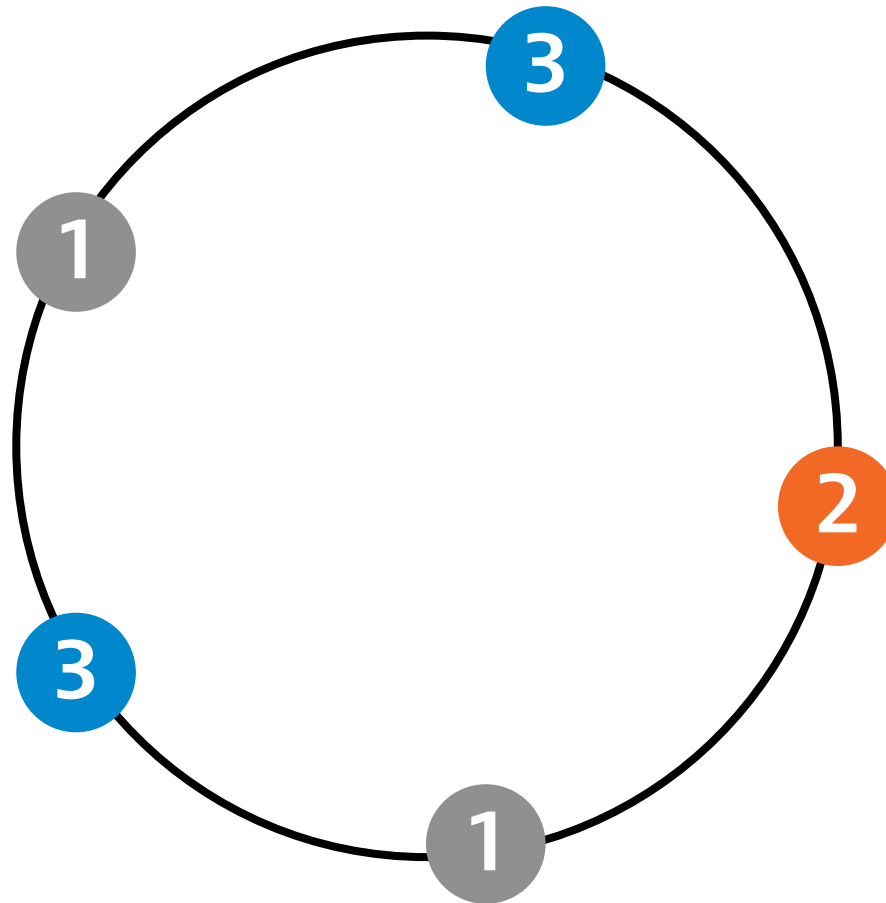
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



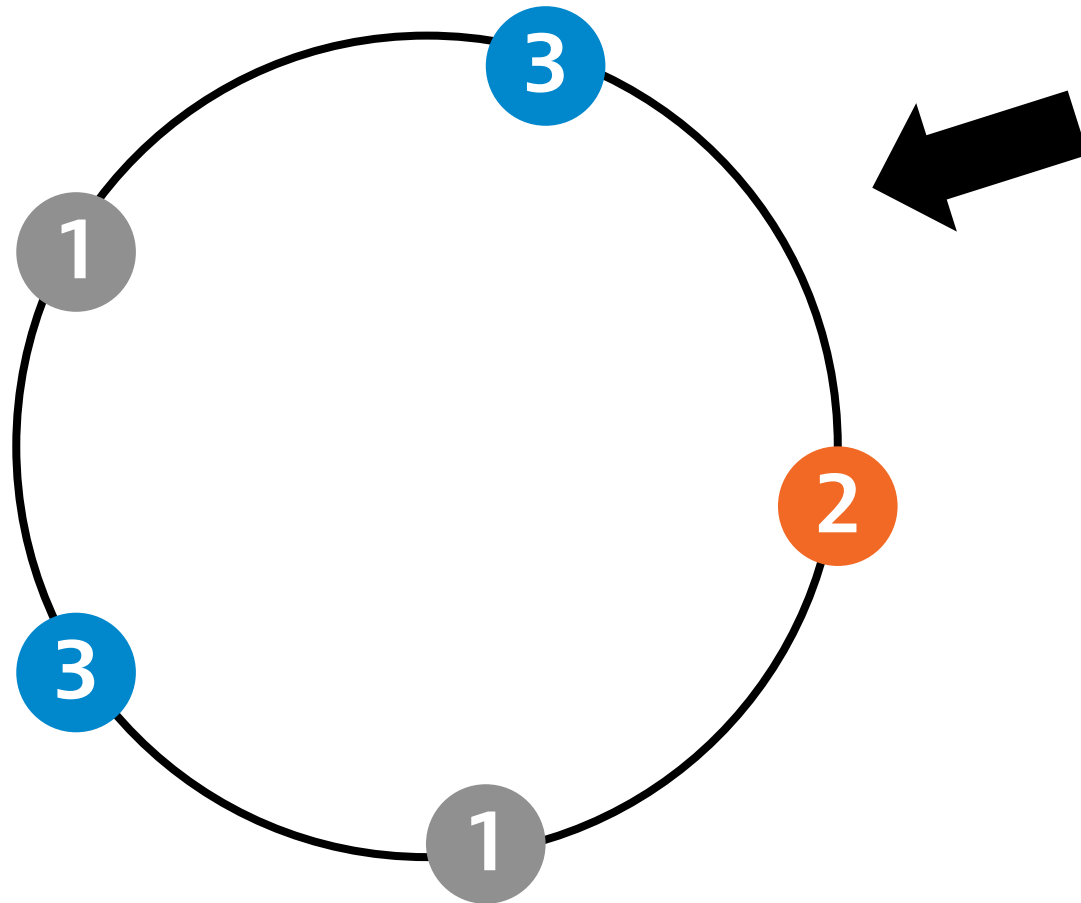
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



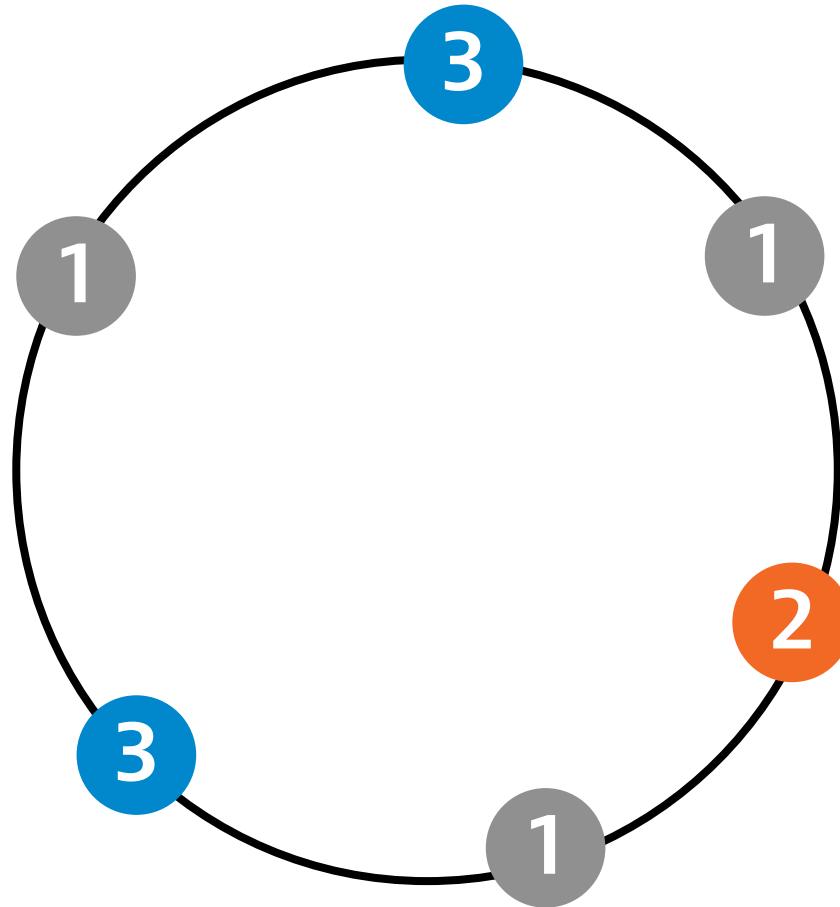
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



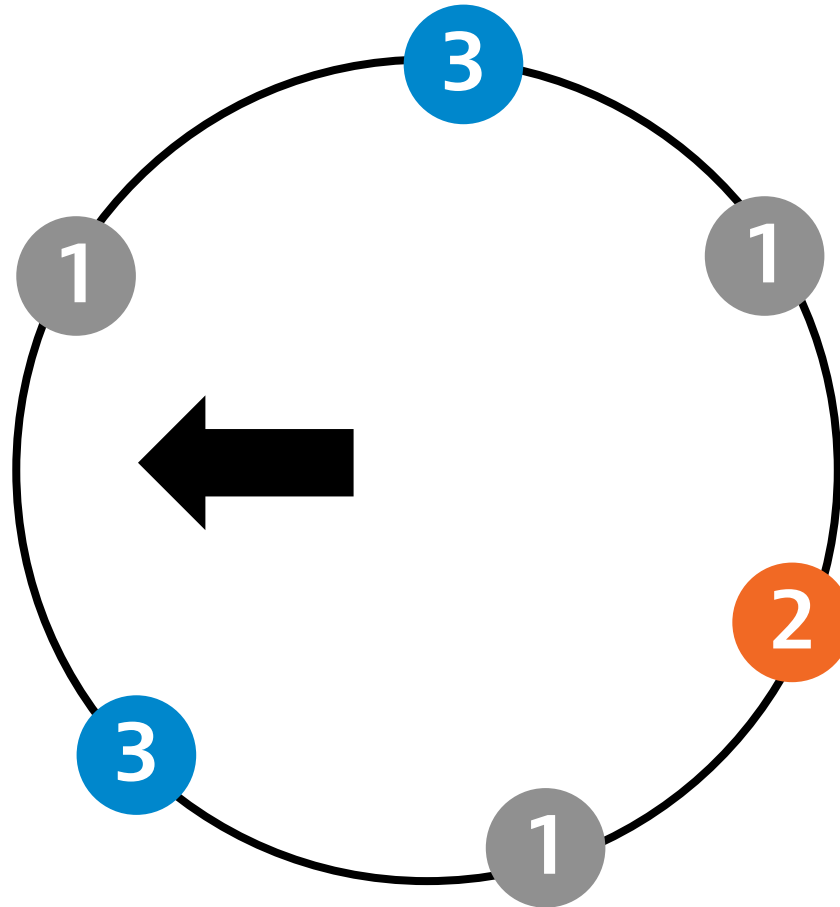
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



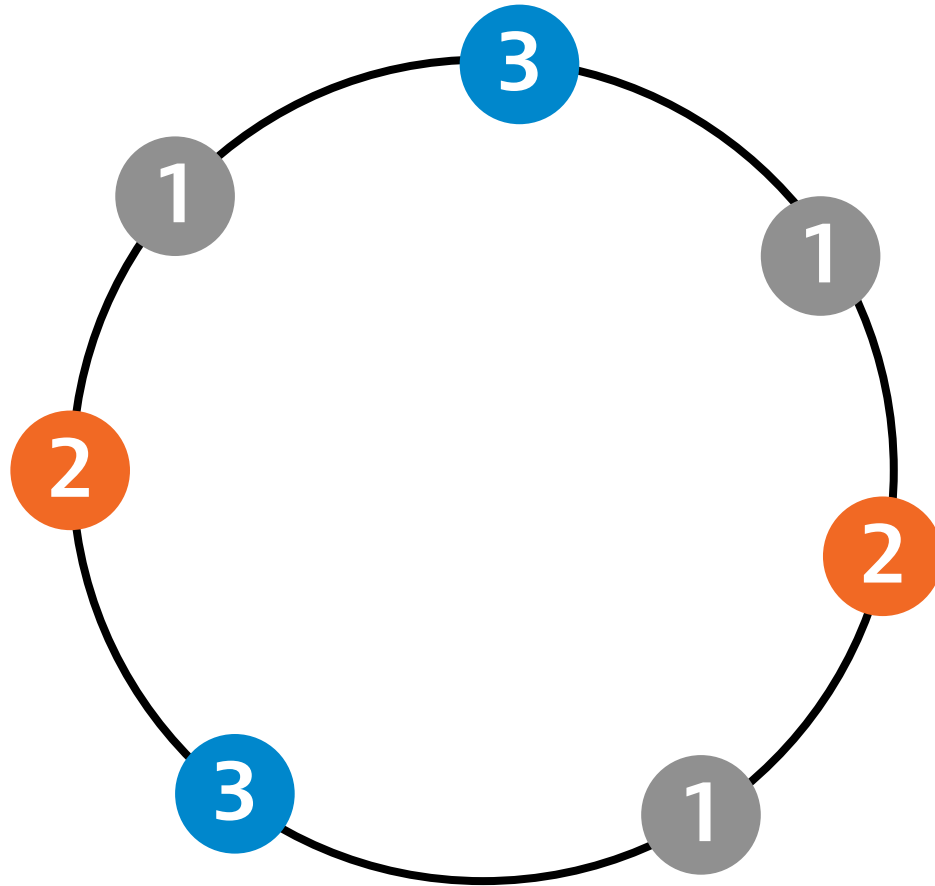
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



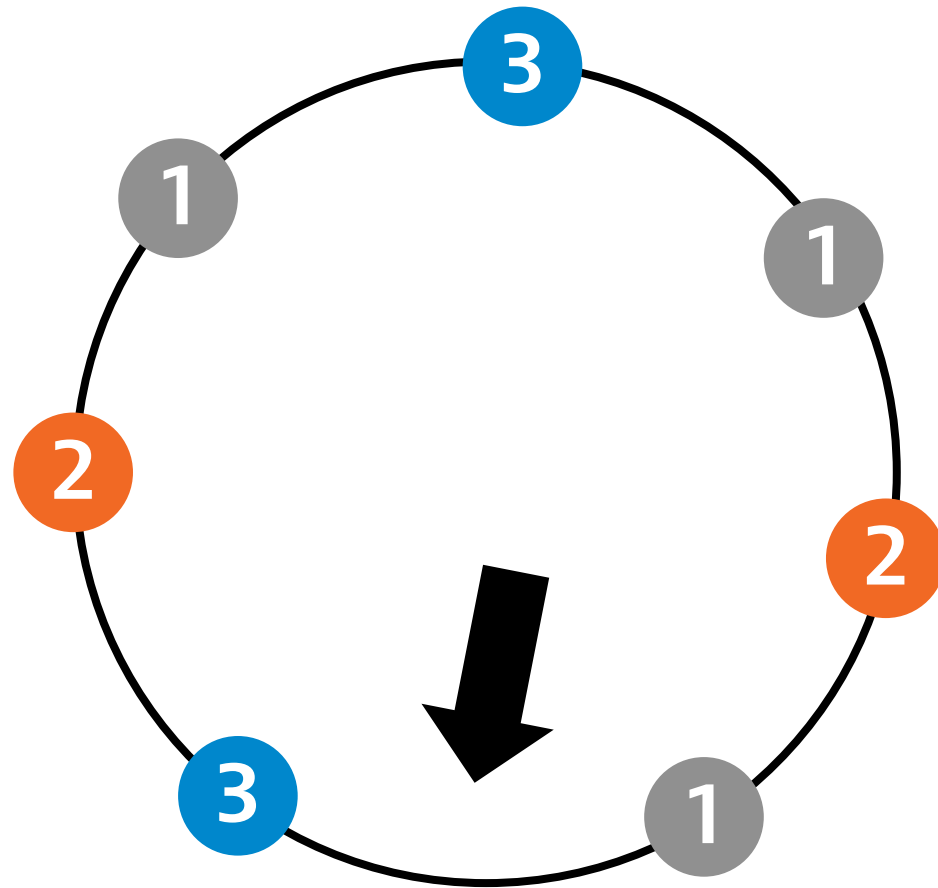
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



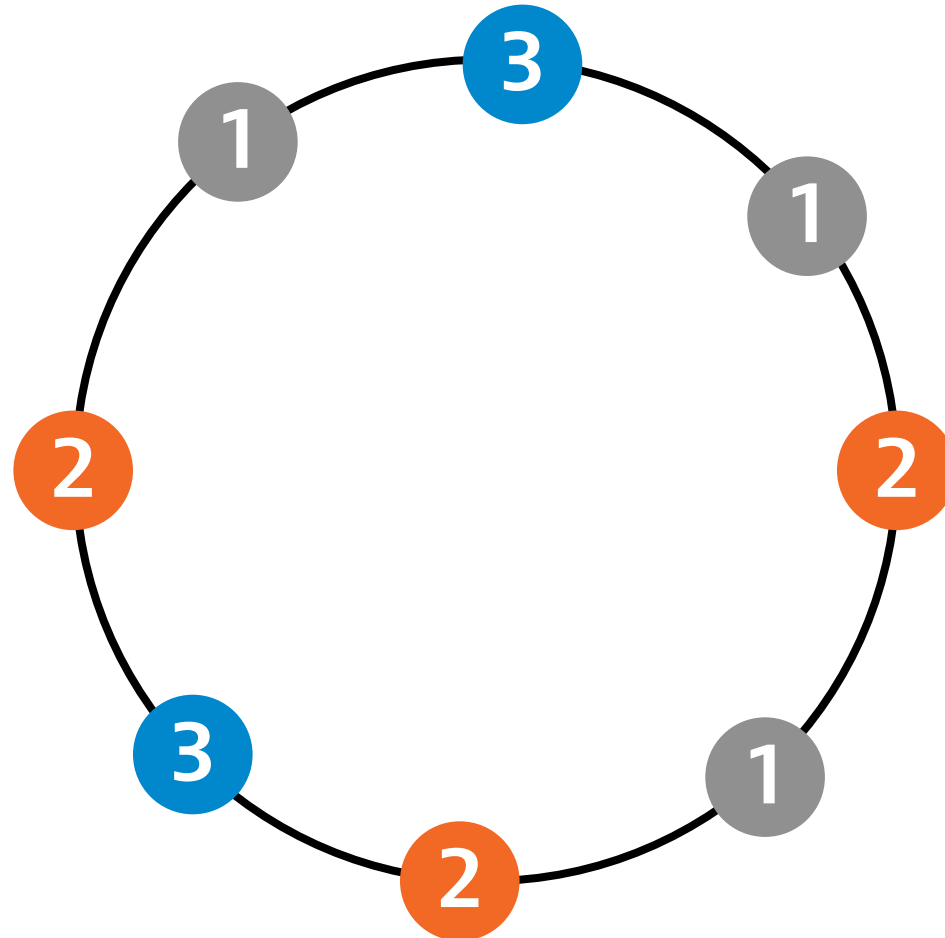
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



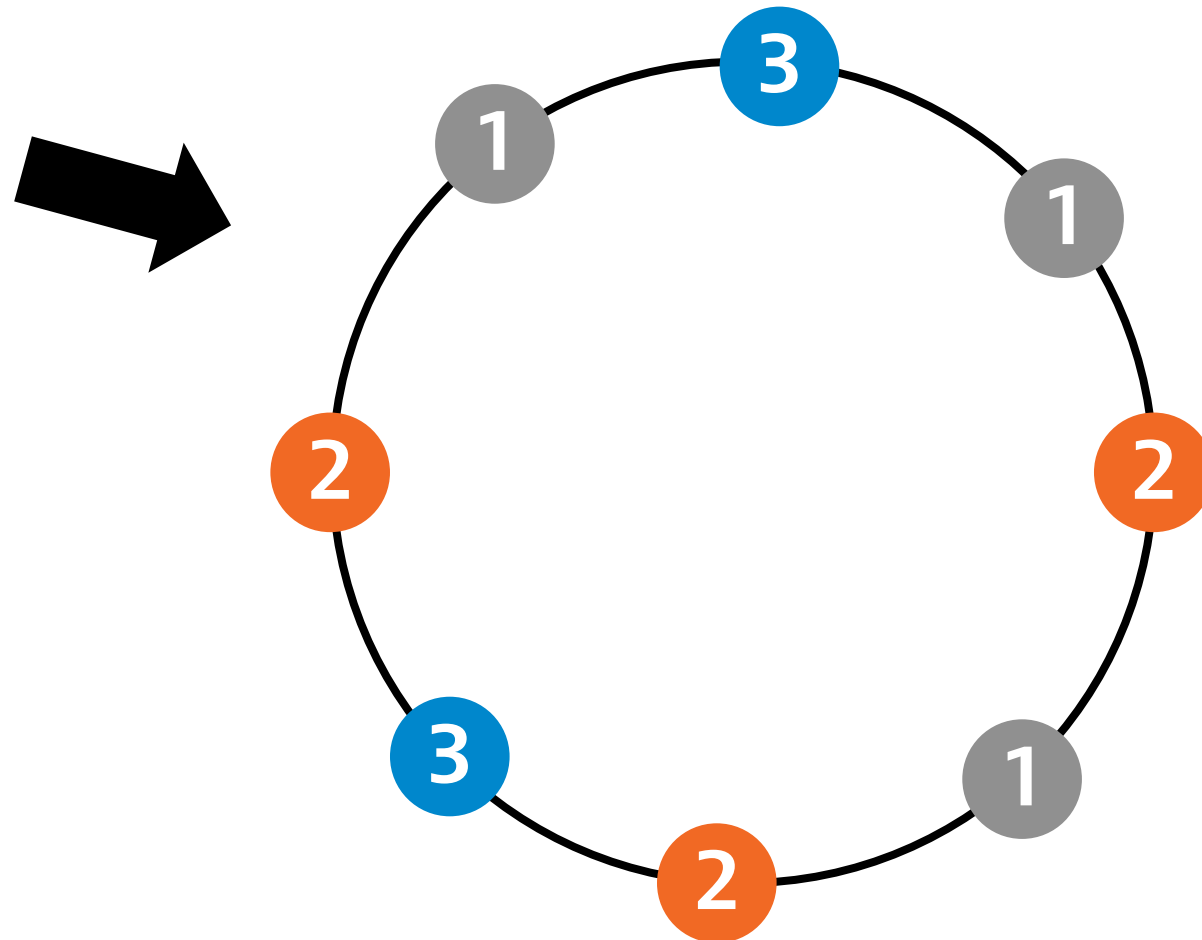
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



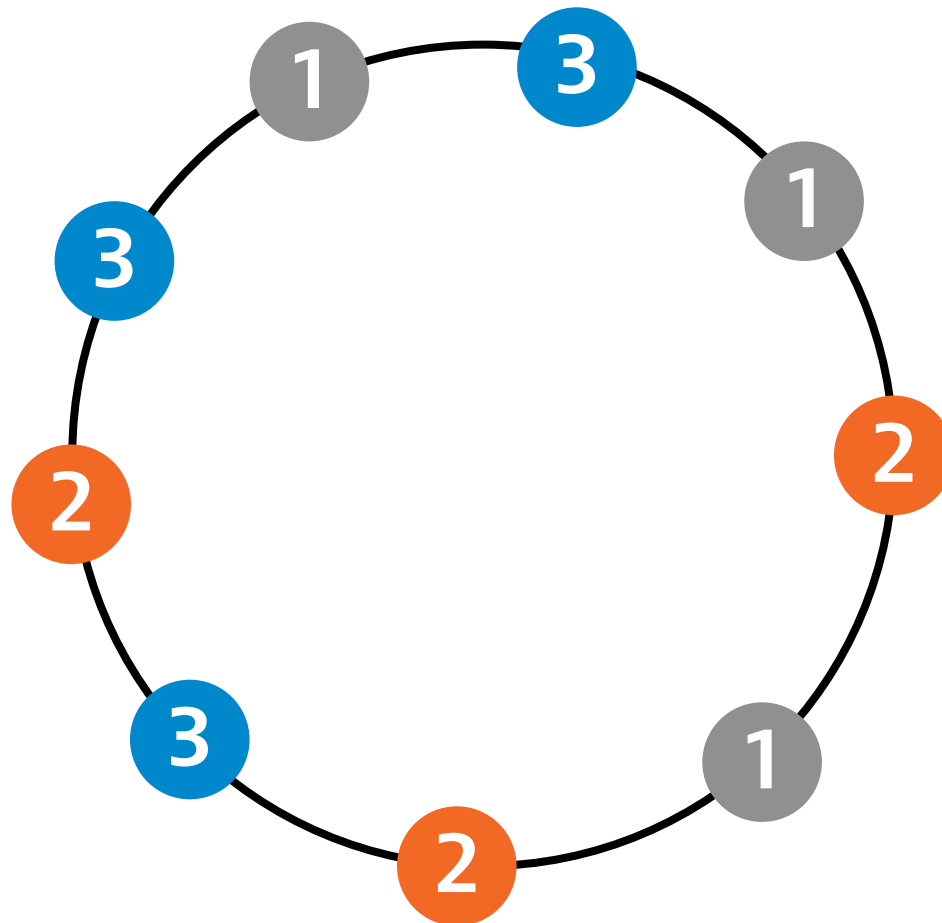
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



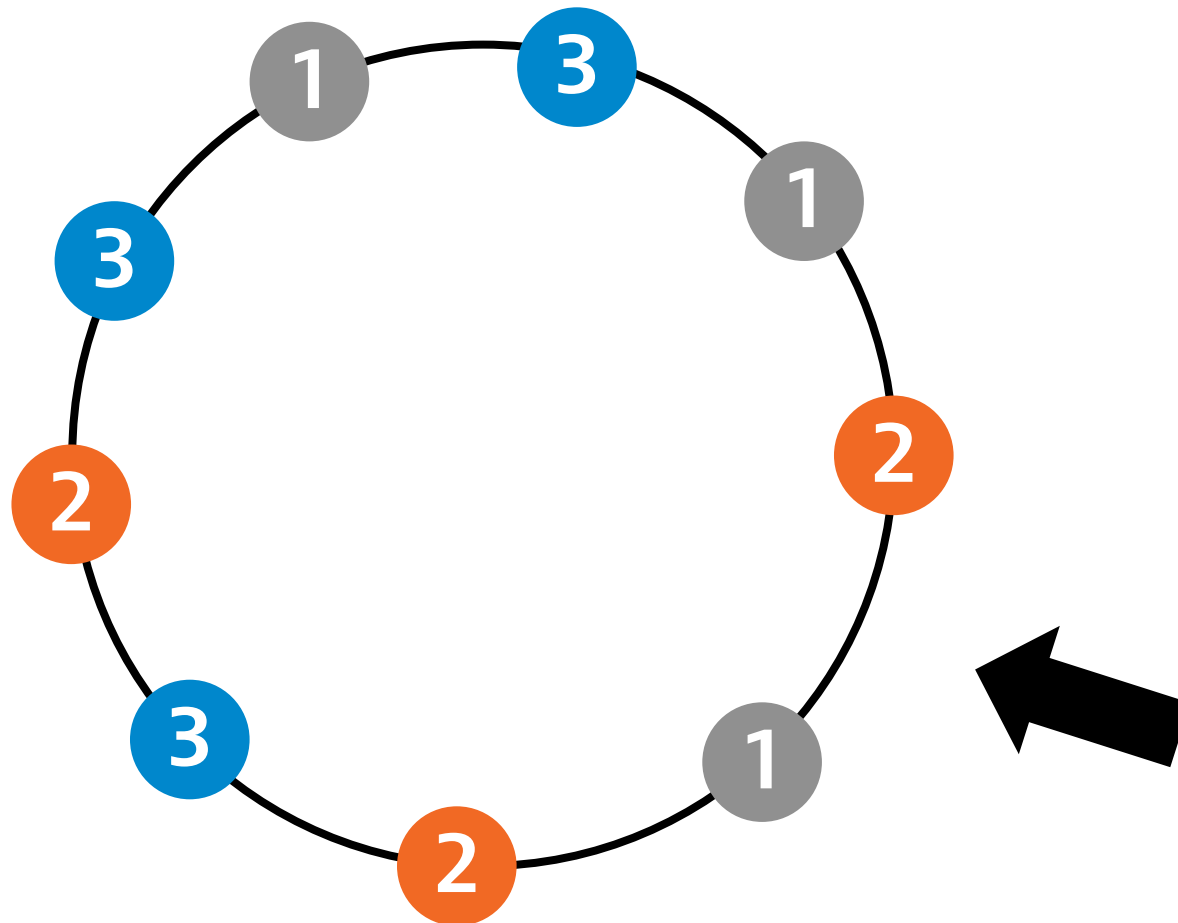
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



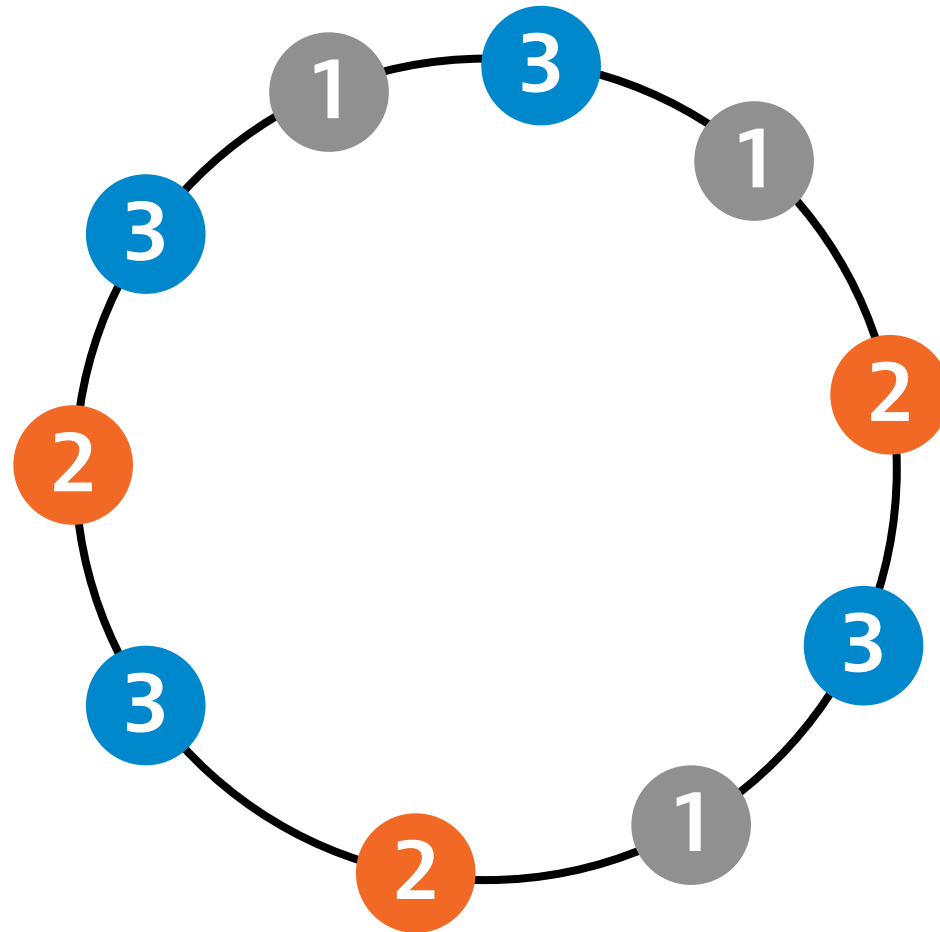
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



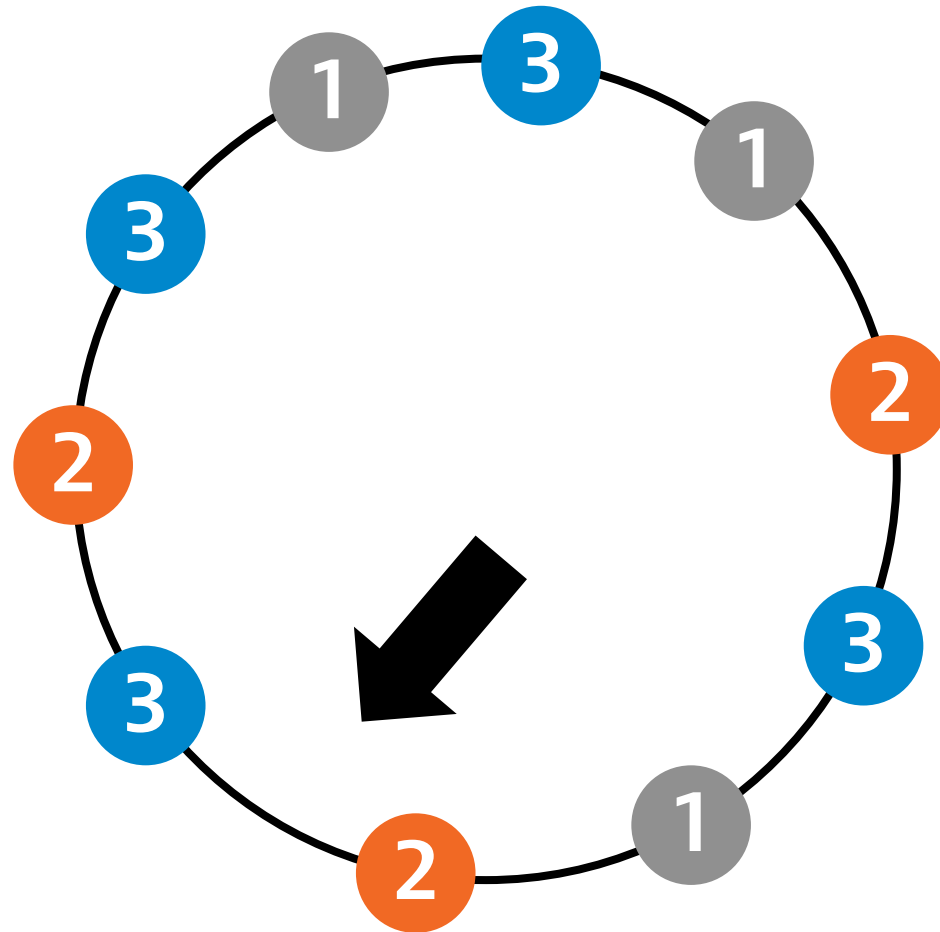
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



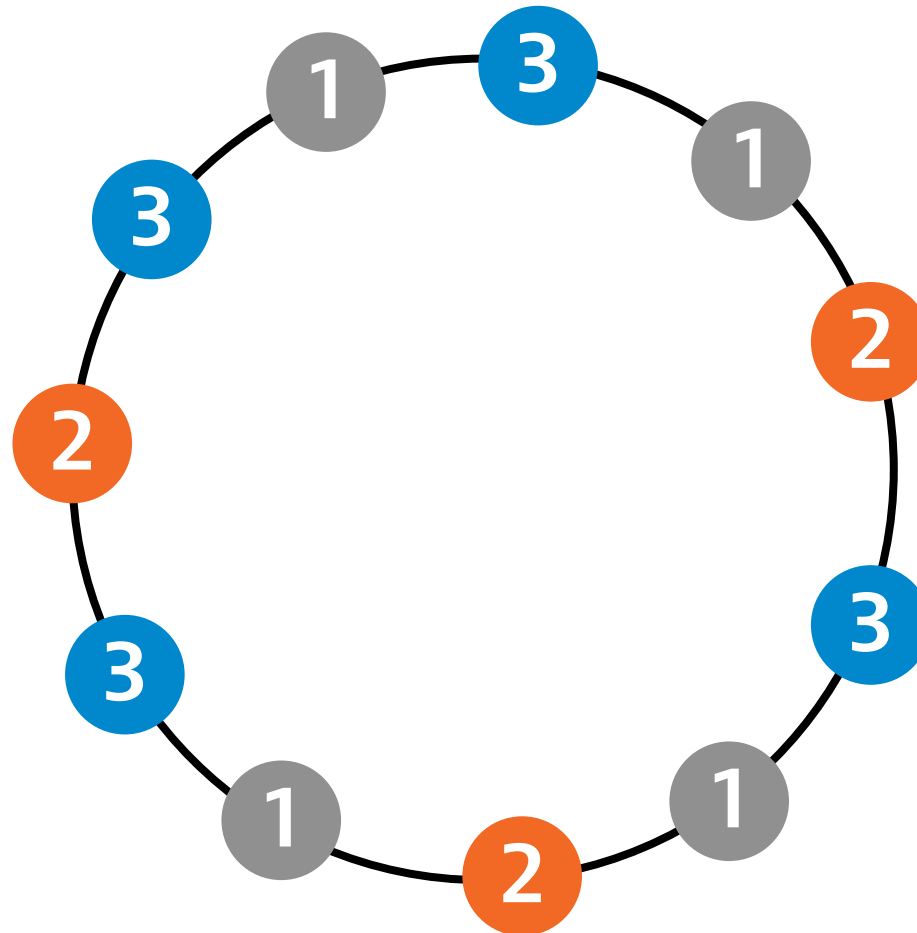
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



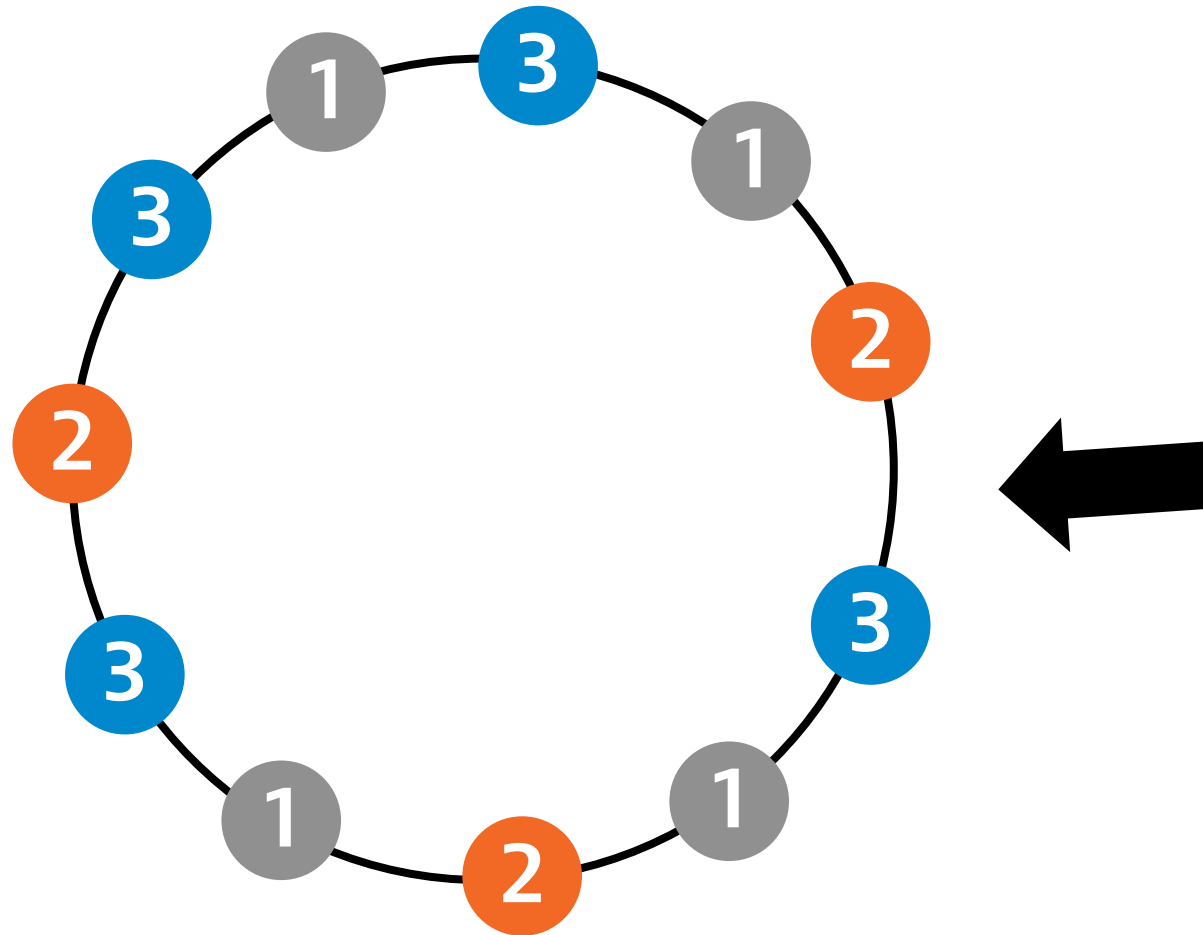
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



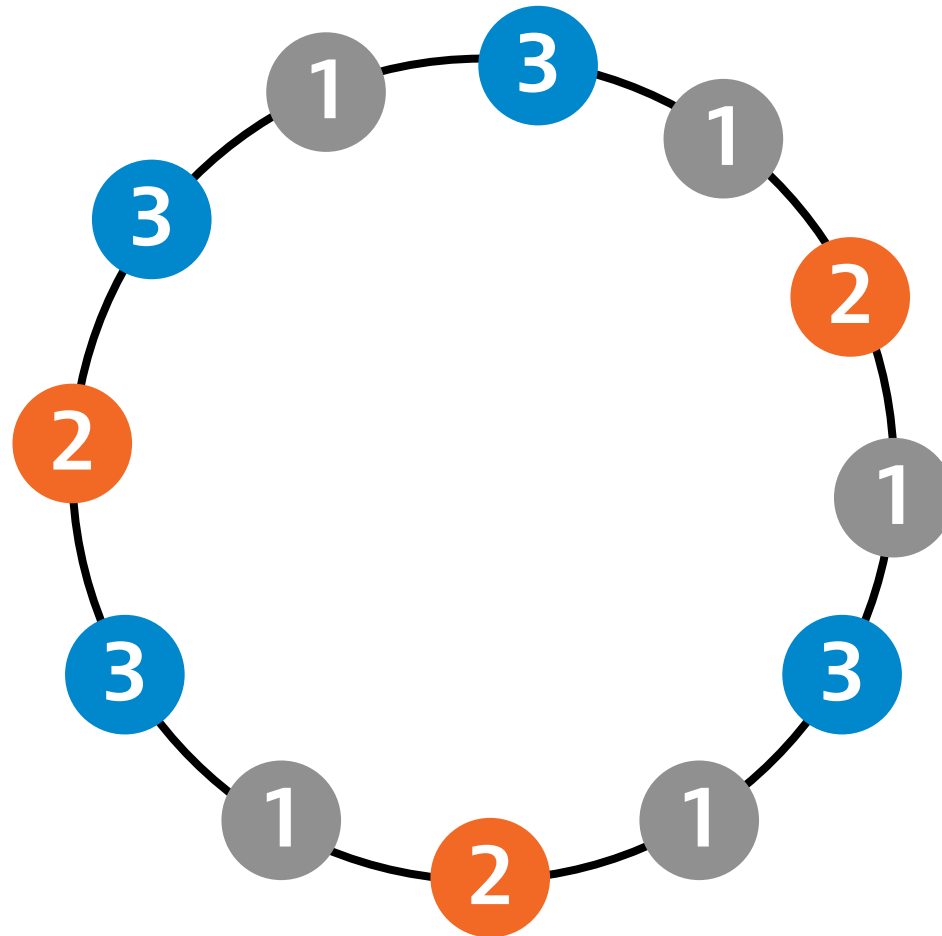
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



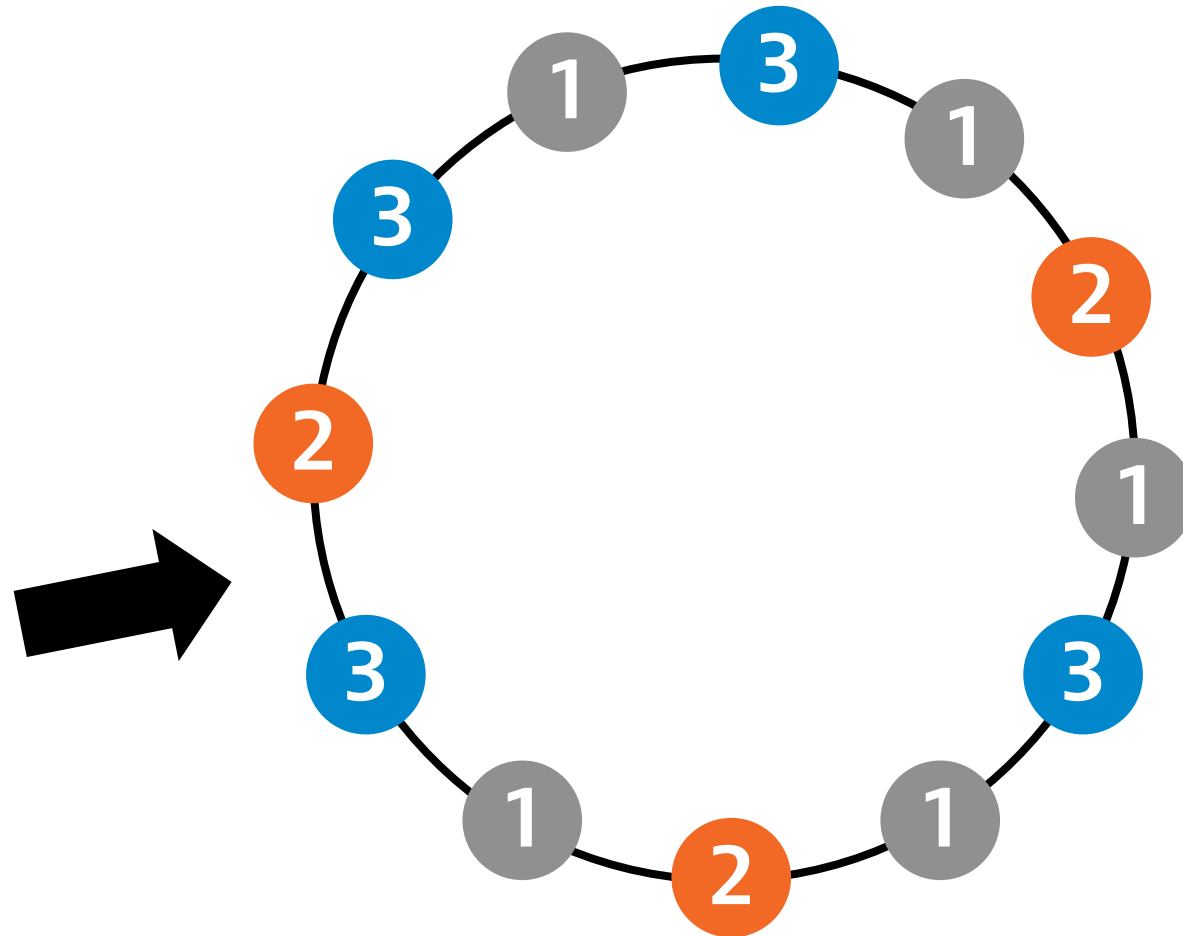
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



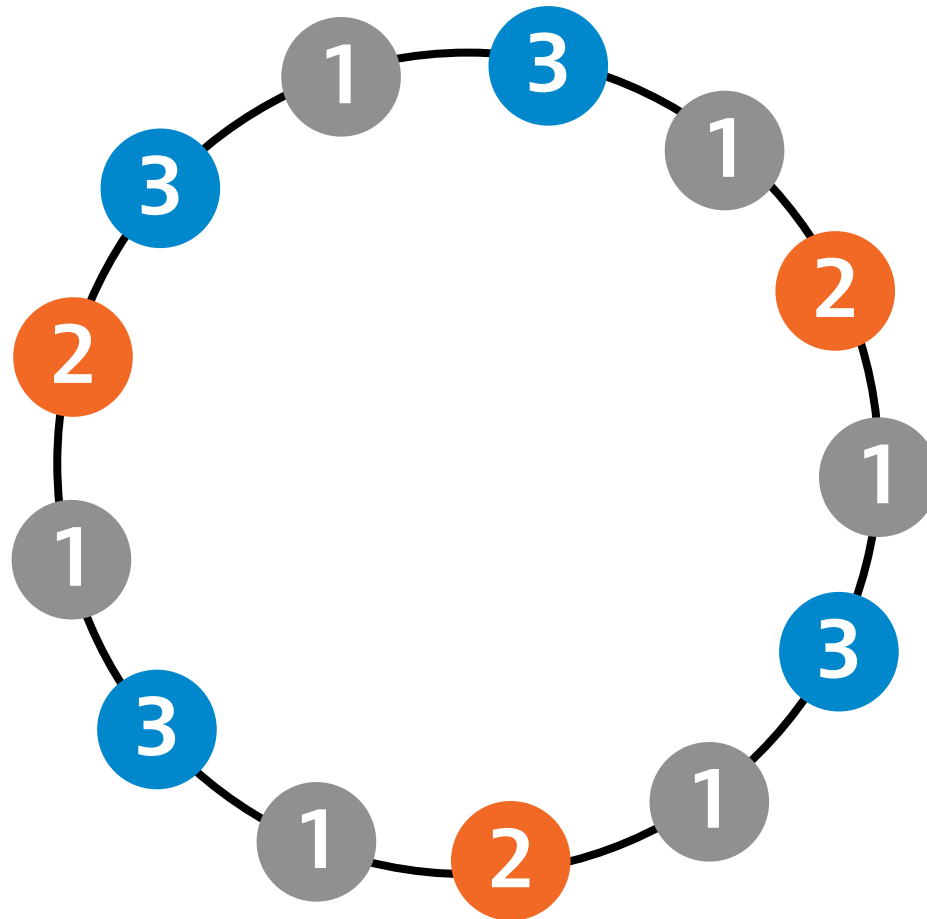
Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



Finely dependent distribution of 3-colorings in cycles

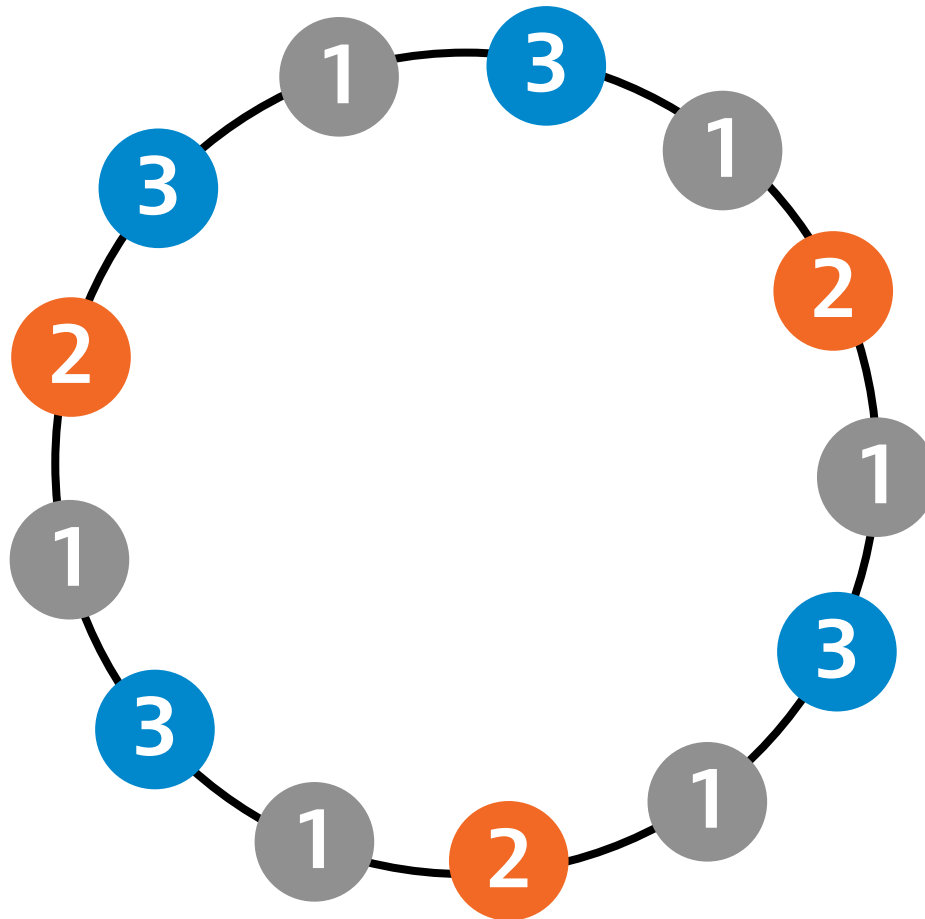
[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]



Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Trivial: stop after
 n steps \rightarrow properly
3-colored n -cycle

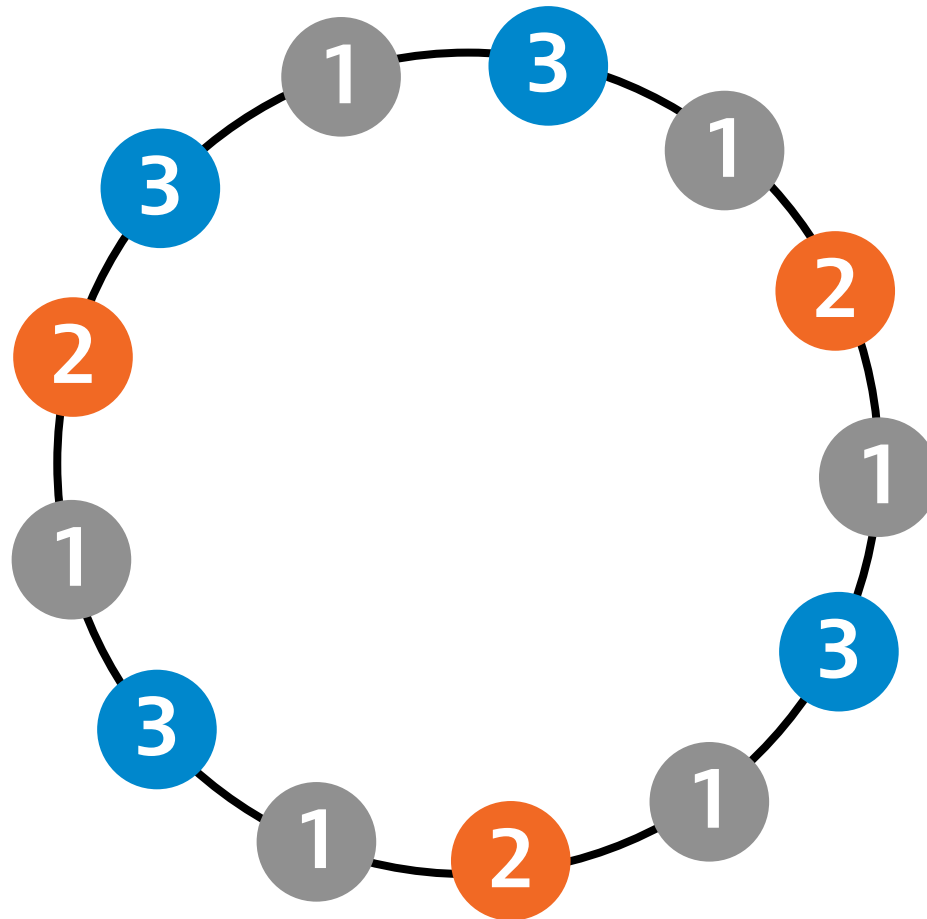


Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Trivial: stop after n steps \rightarrow properly 3-colored n -cycle

Surprise: distribution is 2-dependent!

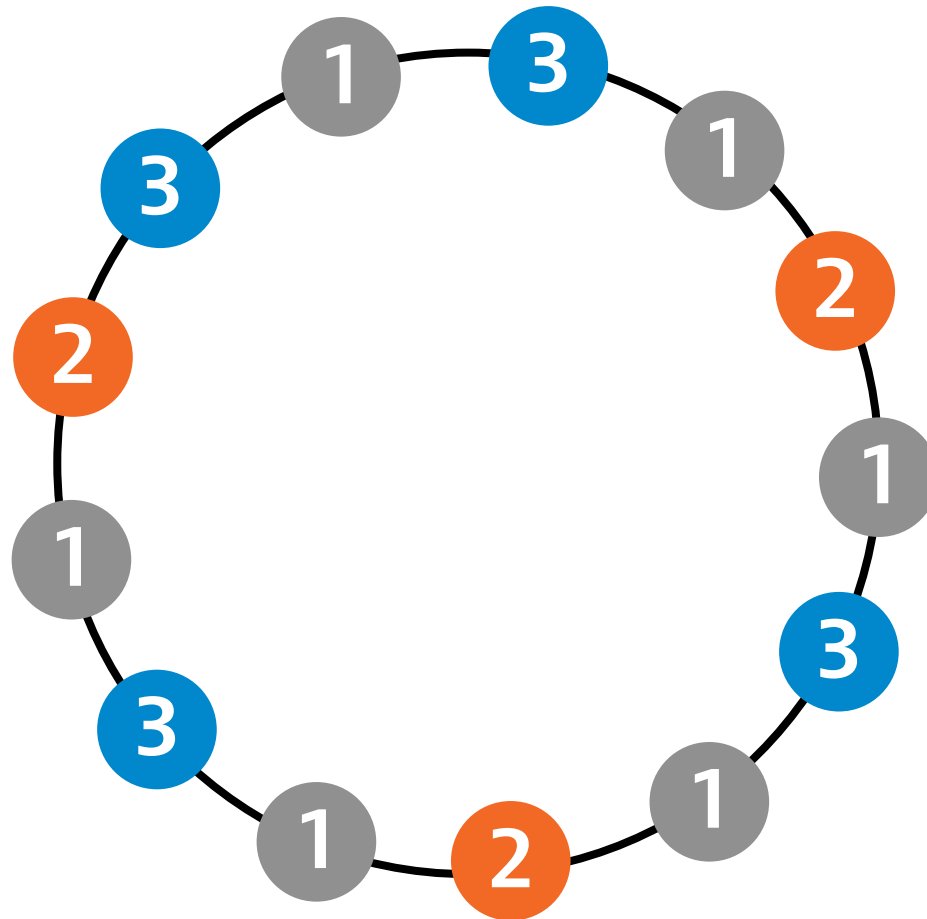


Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Trivial: stop after n steps \rightarrow properly 3-colored n -cycle

Surprise: distribution is 2-dependent!



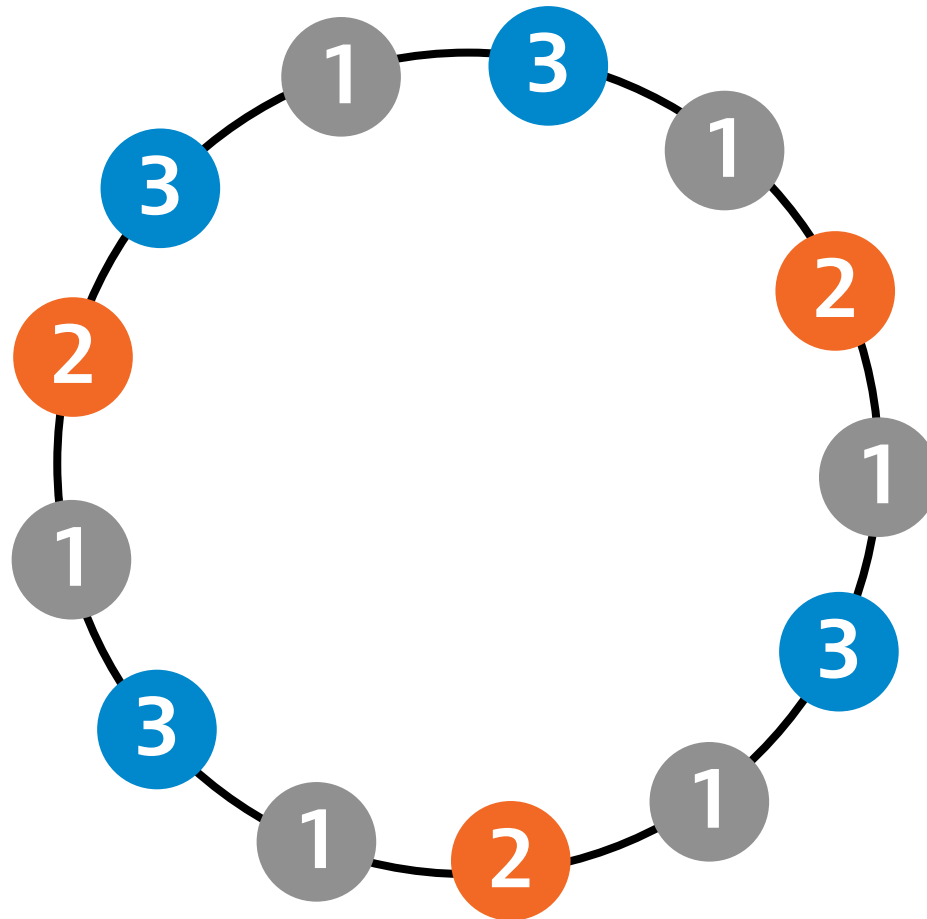
Same for 4 colors: distribution is 1-dependent!

Finely dependent distribution of 3-colorings in cycles

[Holroyd & Liggett 2016, Holroyd, Hutchcroft, Levy 2018]

Trivial: stop after n steps \rightarrow properly 3-colored n -cycle

Surprise: distribution is 2-dependent!



Same for 4 colors: distribution is 1-dependent!

Same for 5 colors: distribution is **not** finitely dependent!!!

- **Known:** constant-round *classical or quantum* algorithm → *finitely dependent distribution*
- **Open:** are there problems that admit finitely dependent distributions but no quantum algorithm?
- **Open:** when does a finitely dependent distribution exist?
- **Key example:** coloring cycles